

Pari-mutuel Betting on Permutations

Shipra Agrawal

Joint work with Zizhuo Wang and Yinyu Ye
Stanford University

Betting as Prediction Markets

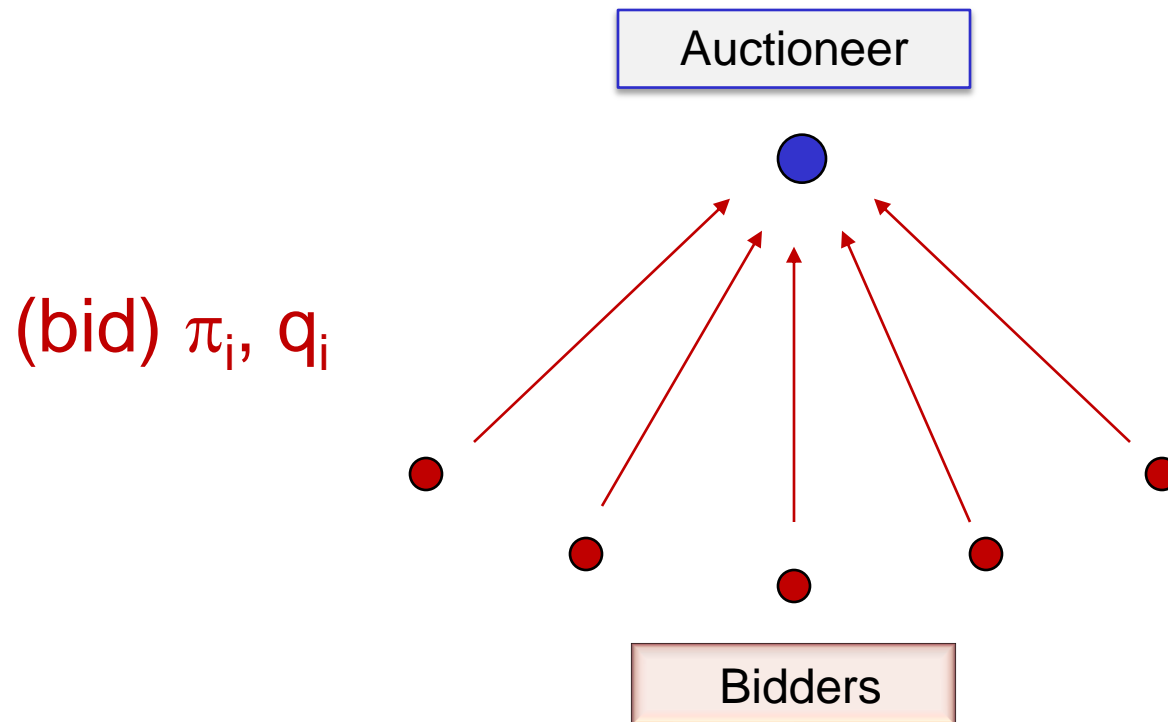
- **Prediction markets**
 - Primary purpose is forecasting events, or the probability that an event will occur
- **Betting markets**
 - Traders speculate the outcome and bet on events they consider probable
- **Prices of bets reflect traders' beliefs**
 - Price of a bet is high if many traders placed similar bets (high price for low risk bets)
 - Price discovery is the end goal of the market organizer

Betting on Permutations

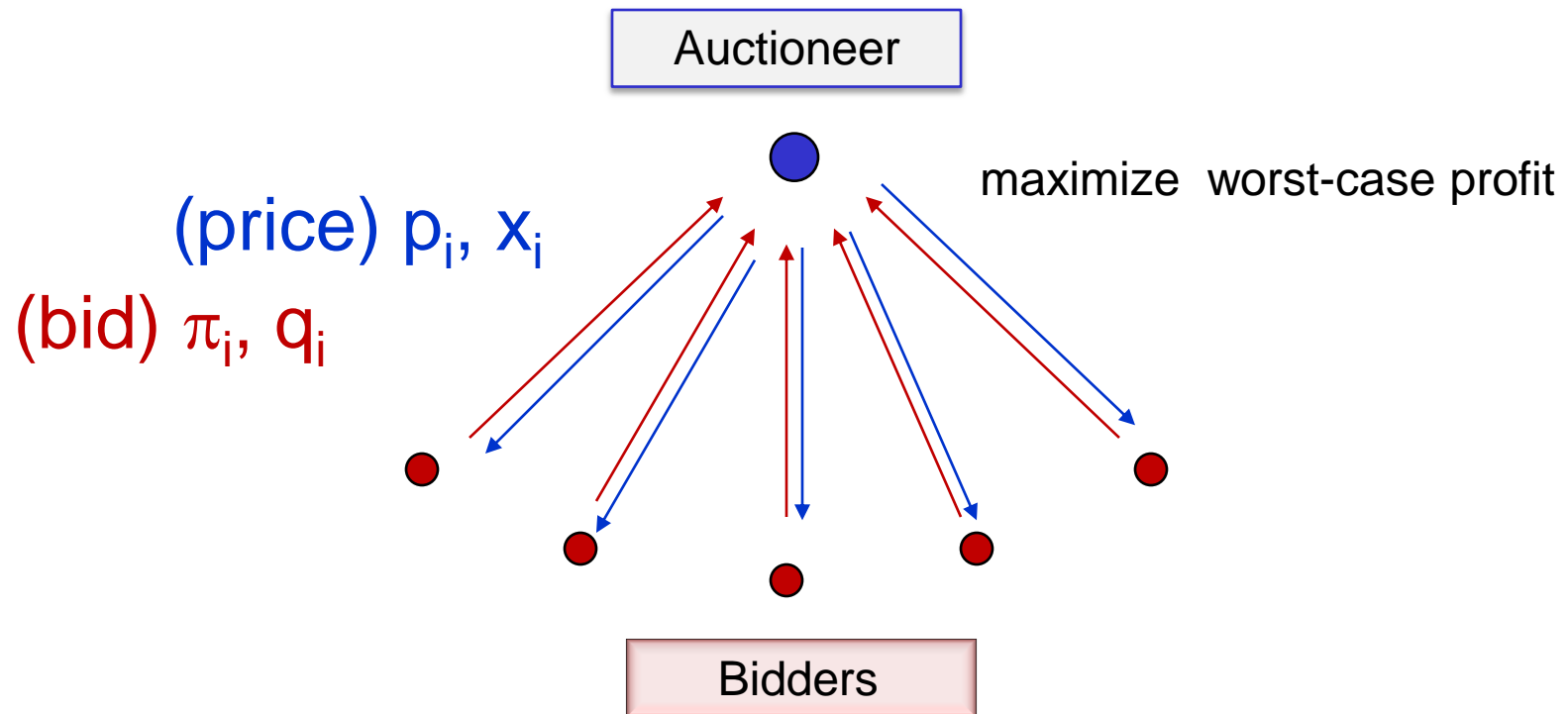
- Horse race
 - $n!$ outcomes, 2^n subset of events
 - Need for restrictive betting languages



Pari-mutuel call auction model

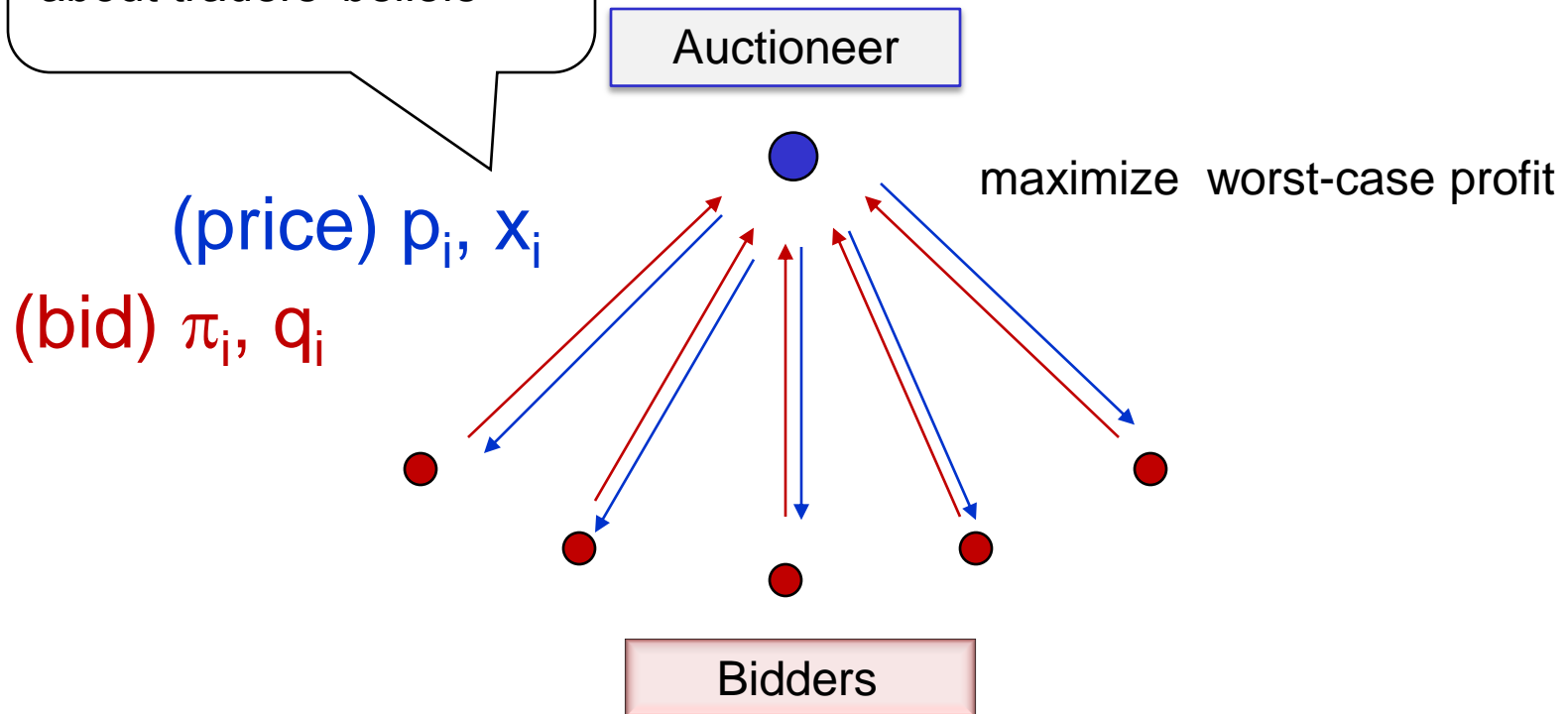


Pari-mutuel call auction model



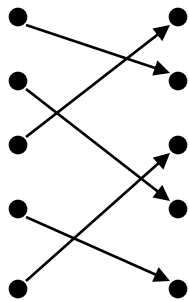
Pari-mutuel call auction model

- Pari-mutuel price
- Reflect information about traders' beliefs



Permutation Betting Mechanism

Horses



Ranks

Outcome

Permutation matrix

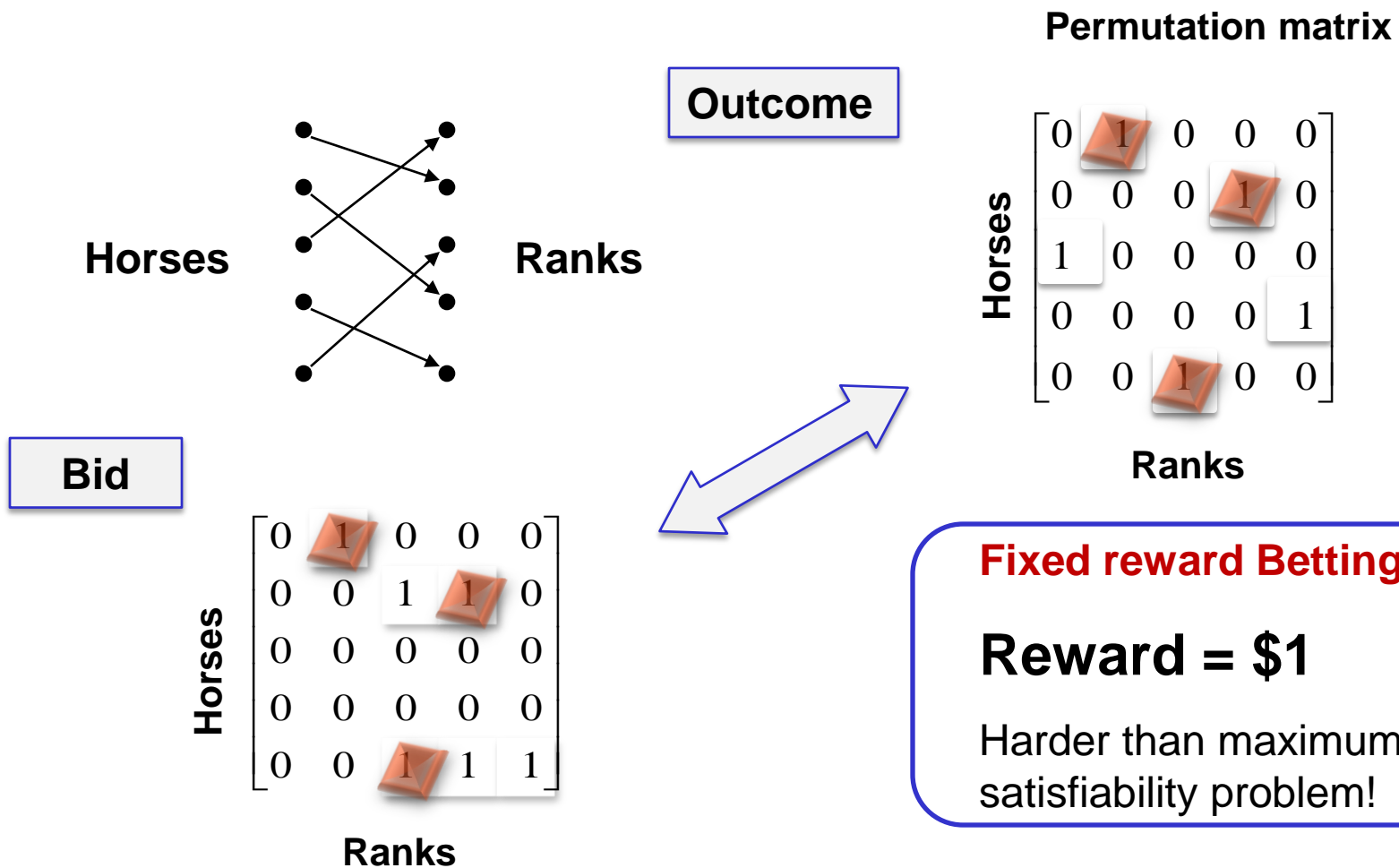
Horses	0	1	0	0	0
	0	0	0	1	0
	1	0	0	0	0
	0	0	0	0	1
	0	0	1	0	0
Ranks					

Bid

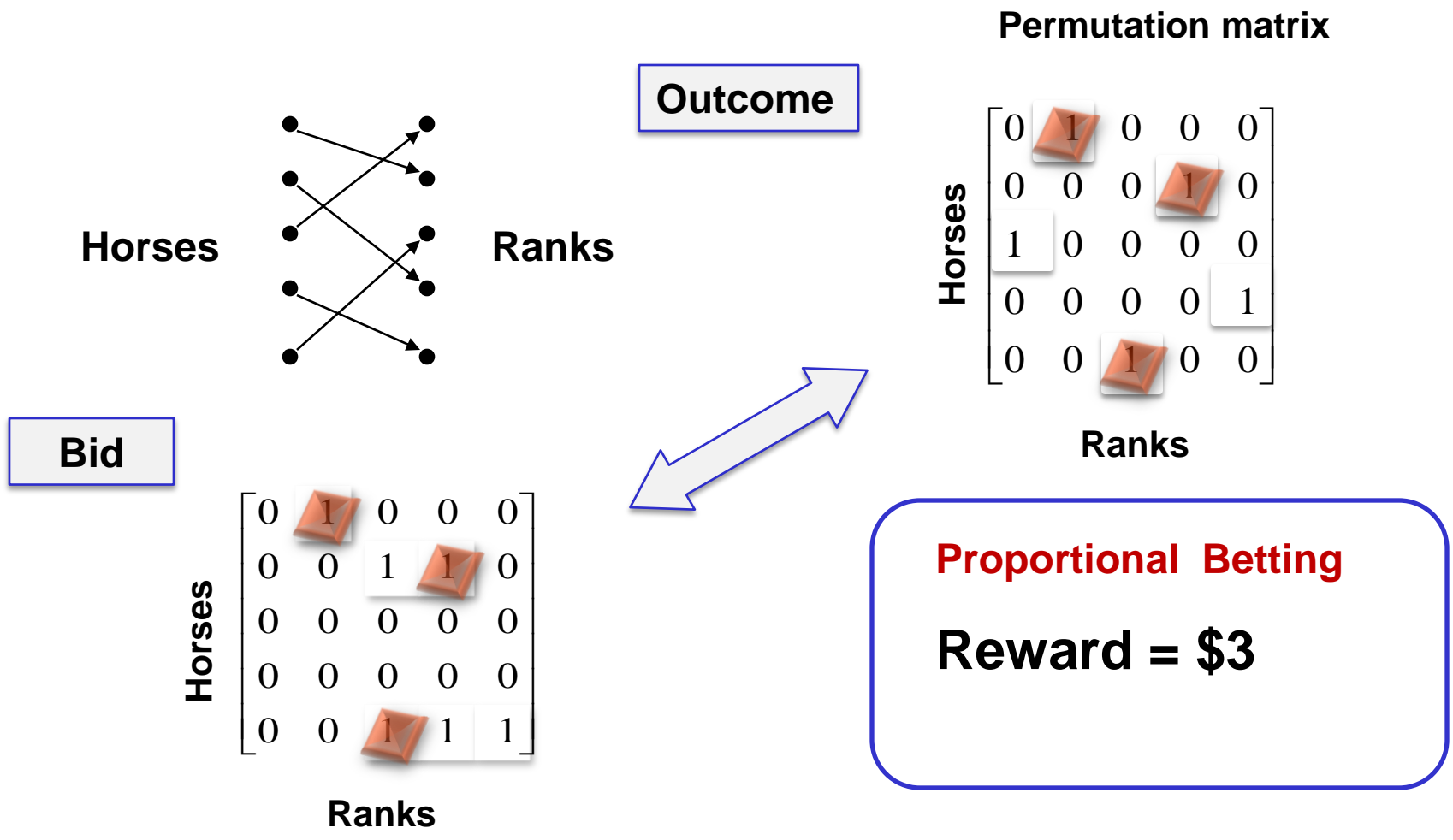
Horses	0	1	0	0	0
	0	0	1	1	0
	0	0	0	0	0
	0	0	0	0	0
	0	0	1	1	1
Ranks					

Ranks

Permutation Betting Mechanism



Permutation Betting Mechanism



Mathematical Formulation

n! prices



$$\begin{array}{ll} \max_{\mathbf{x}} & \pi^T \mathbf{x} - r \\ \text{s.t.} & r \geq (\sum_k x_k A_k) \bullet M_{\sigma} \quad \text{for all permutations } \sigma \\ & 0 \leq \mathbf{x} \leq \mathbf{q} \end{array}$$

\mathbf{p}_{σ}

Mathematical Formulation

$n!$ prices

$$\begin{array}{ll} \max_{\mathbf{x}} & \pi^T \mathbf{x} - r \\ \text{s.t.} & r \geq (\sum_k x_k A_k) \bullet M_{\sigma} \quad \text{for all permutations } \sigma \\ & 0 \leq \mathbf{x} \leq \mathbf{q} \end{array}$$

\mathbf{p}_{σ}

$$\begin{array}{ll} r = \max_M & (\sum_k x_k A_k) \bullet M \\ \text{s.t.} & M^T \mathbf{e} = \mathbf{e} \\ & M \mathbf{e} = \mathbf{e} \\ & M_{ij} \geq 0 \end{array}$$

Mathematical Formulation

$$\begin{aligned} \max_x \quad & \pi^T x - r \\ \text{s.t.} \quad & r \geq (\sum_k x_k A_k) \bullet M_\sigma \quad \text{for all permutations } \sigma \\ & 0 \leq x \leq q \end{aligned}$$

$n!$ prices

p_σ

$$\begin{aligned} r = \max_M \quad & (\sum_k x_k A_k) \bullet M = \min_{v,w} \quad e^T v + e^T w \\ \text{s.t.} \quad & M^T e = e \quad \text{s.t.} \quad v_j + w_j \geq (\sum_k x_k A_k)_{ij} \\ & M e = e \\ & M_{ij} \geq 0 \end{aligned}$$

n^2 prices

$$\begin{aligned} \max_x \quad & \pi^T x - e^T v - e^T w \\ \text{s.t.} \quad & v_j + w_j \geq (\sum_k x_k A_k)_{ij} \quad 1 \leq i, j \leq n \\ & 0 \leq x \leq q \end{aligned}$$

Q_{ij}

Marginal Prices

$$\begin{array}{c} \mathbf{Q} \\ \text{Horses} \end{array} \begin{array}{c} \text{Ranks} \\ \left[\begin{array}{ccccc} 0.2 & 0.1 & 0.3 & 0.3 & 0.1 \\ 0.1 & 0.3 & 0.1 & 0.2 & 0.3 \\ 0.2 & 0.2 & 0.1 & 0.3 & 0.2 \\ 0.3 & 0.3 & 0.2 & 0.1 & 0.1 \\ 0.2 & 0.1 & 0.3 & 0.1 & 0.3 \end{array} \right] \end{array} = \mathbf{1} \quad \begin{array}{c} \text{Marginal} \\ \text{Distributions} \end{array}$$

$= \mathbf{1}$

Theorem:

- One can compute in polynomial-time, an $n \times n$ marginal price matrix Q which is sufficient to price the bets in the Proportional Betting mechanism.
- Further, the price matrix is unique, pari-mutuel, and satisfies the desired price-consistency constraints.

Pricing the permutations

		Ranks					
Horses		0.2	0.1	0.3	0.3	0.1	= 1
		0.1	0.3	0.1	0.2	0.3	
		0.2	0.2	0.1	0.3	0.2	
		0.3	0.3	0.2	0.1	0.1	
		0.2	0.1	0.3	0.1	0.3	
		= 1					

$$p_1 \begin{bmatrix} 1 & & & & \\ & 1 & & & \\ & & 1 & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix} + p_2 \begin{bmatrix} & 1 & & & \\ 1 & & & & \\ & & 1 & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix} + p_3 \begin{bmatrix} & & 1 & & \\ & 1 & & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix} + \dots + p_{n!} \begin{bmatrix} & & & & 1 \\ & & & & & 1 \\ & & & & & & 1 \\ & & & & & & & 1 \\ & & & & & & & & 1 \end{bmatrix}$$

Joint Distribution p over permutations

Maximum entropy criteria

$$\min_p \sum_{\sigma} p_{\sigma} \log p_{\sigma}$$

$$\text{s.t. } \sum_{\sigma} p_{\sigma} M_{\sigma} = Q \quad \Rightarrow \quad p_{\sigma} = e^{Y \cdot M_{\sigma}}$$
$$p_{\sigma} \geq 0$$

- Closest distribution to uniform prior
- Completely specified by n^2 parameters
- Maximum likelihood estimator

Complexity of computing

$$\begin{aligned} \max_Y \quad & Q \bullet Y - 1 \\ \text{s.t.} \quad & \sum_{\sigma} e^{Y \bullet M_{\sigma}} M_{\sigma} \geq Q \\ & Y \leq 0 \end{aligned}$$

- #P-hard to compute parameter Y
 - Reduction from the problem of computing permanent of a non-negative matrix
 - Permanent of matrix $[e^{Y_{ij}}]$ is $\sum_{\sigma} e^{Y \bullet M_{\sigma}}$

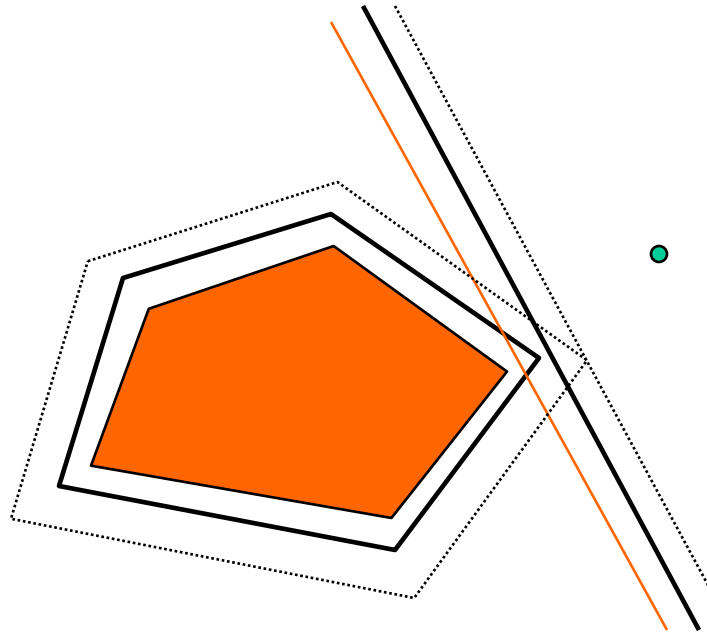
Polynomial-time approximation

$$\begin{aligned} \max_Y \quad & Q \bullet Y - 1 \\ \text{s.t.} \quad & \sum_{\sigma} e^{Y \bullet M_{\sigma}} M_{\sigma} \geq Q \\ & Y \leq 0 \end{aligned}$$

- Use ellipsoid method for convex optimization
 - In conjunction with sampling to estimate the separating hyperplane at each step.
 - MCMC method for estimating permanent of a matrix used for sampling

Illustration

- Example: problem of finding feasible point in convex set



- We are guaranteed to find a feasible point in the outermost convex set

Ongoing and future work

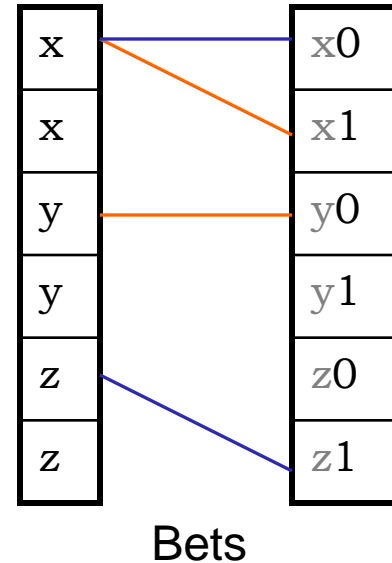
- Sequential auction model
 - Traders come one by one
 - Sequential decisions
- Truthfulness of the price vector
 - How well price captures the beliefs?
 - Empirical studies exploring this aspect
- Betting mechanisms
 - What are the gains of fixed reward betting from information collection perspective?
 - What other betting mechanisms are solvable/useful?

Questions?

Reduction

$$(x \vee \neg y) \wedge (\neg x \vee z) \wedge \dots$$

Clauses



- A bet is won if and only if a clause is satisfied
- Assignment satisfying K clauses satisfies K or more bets and vice-versa