## Integrating Functions via Distance Sensitive Hashing

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## Problem and Motivation

## Problem

$$
\mathcal{X}=\left\{x_{1}, x_{2}, x_{3}, \ldots, x_{n}\right\} \subset \mathcal{S}^{d-1}, \phi:[-1,1] \rightarrow \mathbb{R}, \text { query } y \in \mathcal{S}^{d-1}
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Partition Function Estimation

## Kernel Density Estimation

$\mathcal{X}=\left\{x_{1}, \ldots, x_{n}\right\} \subset r \mathcal{S}^{d-1}$, distribution $\mathcal{D}$, prob. of $y \in \mathcal{S}^{d-1}$ ?


Problem: Data structure that answers queries in sub-linear time?

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\begin{gathered}
\operatorname{KDE}_{\mathcal{X}}(y)=\frac{1}{n} \sum_{i=1}^{n} K(x, y) \\
K(x, y)=e^{-\frac{\|x-y\|^{2}}{\sigma^{2}}}=e^{\frac{2 r^{2}}{\sigma^{2}}(\langle x, y\rangle-1)} \\
\text { outlier detection, clustering, } \ldots
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## Empirical Gradient Estimation

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\left\{\left(x_{i}, s_{i}\right)\right\}_{i=1}^{n} \subset \mathbb{R}^{d} \times\{ \pm 1\}, \quad \mathcal{L}(\cdot)=\sum_{i=1}^{n} \ell\left(\left\langle s_{i} x_{i}, \cdot\right\rangle\right), \nabla_{y} \mathcal{L}(y) ?
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Discriminative sampling for $S G$

Problem: Data structure that gets lower-variance SG?

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## Applications of "Partition Function Estimation"

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Z(y)=\sum_{i=1}^{n} e^{\phi\left(\left(x x_{i}, y\right)\right.}
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- Kernel Density Estimation
- Robust Optimization: let $\beta=c \log n$

- Variance reduction for Stochastic Gradients


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## Modelling Binary data with Exponential Families

$\mathcal{X}=\{-1,+1\}^{d}$, parameter vector $y \in r \mathcal{S}^{d-1}$, density on $\mathcal{X}$

$$
p_{y}(x)=\frac{1}{Z(y)} e^{\langle x, y\rangle}=\frac{1}{Z(y)} e^{e}
$$

with $Z(y)=\sum_{x \in \mathcal{X}} e^{(x, y)}$ being instrumental for
1 Sampling: given $y$ sample $x \sim p_{y}$
2 Maximum likelihood: estimate gradient
3 Hypothesis testing: $z_{1}, \ldots, z_{m} \sim p_{y}, y=y_{1}$ or $y=y_{2}$ ?
$Z(y)$ requires time $2^{d}$ compute exactly

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Previous Work and Main Result

## Reducing Space through Random Sampling

$L(\phi)$ Lipschitz const, $\phi_{\max }-\phi_{\min } \leq 2 L(\phi)$, equiv. estimate:

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Z(y)=\frac{1}{|\mathcal{X}|} \sum_{x \in \mathcal{X}} e^{\phi(\langle x, y\rangle)-\phi_{\max }} \in\left[e^{-2 L(\phi)}, 1\right]
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Random Sampling and Median-of-means:

mean of $n$ samples

median of $K$ means

For $Z(y)=\mu \in\left[\Omega\left(\frac{1}{n}\right), 1\right]$ we require $O\left(\frac{1}{\epsilon^{2}} \frac{1}{\mu} \log \left(\frac{1}{\delta}\right)\right)$ samples.

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n=\Theta\left(e^{2 L(\phi)}\right) \text { and } \mu \in\left[n^{-1}, 1\right]
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m=\underbrace{\left[\frac{6}{\boldsymbol{\epsilon}^{\mathbf{2}}} \exp (2 L(\phi))\right]}_{\text {mean of } n \text { samples }} \cdot \underbrace{\left[9 \log \left(\frac{1}{\delta}\right)\right]}_{\text {median of } K \text { means }}
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## Previous Work

Mussman and Ermon [ICML'16] employed a series of reductions:

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(1 \pm \epsilon) Z^{-1}(y) \rightarrow \log (1+\epsilon)-\text { MIPS } \rightarrow\left(1+\frac{\epsilon}{r^{2}}\right)-\text { ANN }
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to solve the linear case for $x, y \in r \mathcal{S}^{d-1}$ where $L(\phi)=r^{2}$.
1 Empirically outperforms baseline methods
2 ANN data structure by Andoni-Razenshteyn [STOC'15]

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## Main Result

## Theorem [S, Charikar'18]

For any convex function $\phi$ there exists a data structure using

- preprocessing time/space $O\left(\frac{1}{\epsilon^{2}}{ }^{L(\phi)} \log \left(\frac{1}{\delta}\right) \cdot d n\right)$
- answers any query $y$ in $O\left(M_{\phi} \frac{\mathbf{1}}{\boldsymbol{\epsilon}^{2}} \frac{1}{\sqrt{\mu}} \log \left(\frac{1}{\delta}\right) d\right)$ time.
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\text { Assuming } n=\Theta\left(e^{2 L(\phi)}\right) \text { and } \mu=\frac{1}{n}
$$

| Method | Space | QUERY |
| :--- | :---: | ---: |
| Random Sampl. | $O(n)$ | $O(n)$ |
| MIPS-ANN | $O\left(n^{2-O(\epsilon)}\right)$ | $O\left(n^{1-O(\epsilon)}\right)$ |
| OURS | $O\left(n^{\frac{3}{2}+o(1)}\right)$ | $O\left(n^{\frac{1}{2}+o(1)}\right)$ |

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| KERNEL | $\phi(\rho)$ | $L(\phi)$ |
| :--- | :---: | :---: |
| $e^{\langle x, y\rangle}$ | $r^{2} \rho$ | $r^{2}$ |
| $e^{-\\|x-y\\|_{2}^{2}}$ | $2 r^{2}(\rho-1)$ | $2 r^{2}$ |
| $\left(\\|x-y\\|_{2}^{2}+1\right)^{-1}$ | $-\log \left(1+(1-\rho) 2 r^{2}\right)$ | $2 r^{2}$ |
| $\overline{1+e^{-\langle x, y\rangle}}$ | $-\log \left(1+e^{-r^{2} \rho}\right)$ | $r^{2}$ |
| $\left(\langle x, y\rangle+c r^{2}\right)^{-k}$ | $-k \log \left(r^{2}(\rho+c)\right)$ | $\frac{k}{c-1}$ |

## Extensions

We show two reductions:
1 Euclidean space $\rightarrow$ Sphere:
partition in thin annuli and round vectors a la Andoni-Razhenshteyn [STOC'15]
2 Vector sums $\rightarrow$ sum of norms:
Estimate vector sums at least as well as estimating the sum of norms

## Data Structure

## Ingredients:

- $T$ hashing schemes $\mathcal{H}_{1}, \ldots, \mathcal{H}_{T}$ with collision probabilities $p_{t}(\langle x, y\rangle)=\mathbb{P}_{h \sim \mathcal{H}_{t}}[h(x)=h(y)]$. [Distance Sensitive]
- $T$ weight functions $w_{1}, \ldots, w_{T}$ such that $e^{\phi(\langle x, y\rangle)}=\sum_{t} w_{t}(\langle x, y\rangle)$ for all $x, y$. [Convex Decomposition]
Preprocessing:
■ Sample $h_{t} \sim \mathcal{H}_{t}$ and create hash table $H_{t}$ for dataset $X$.


## Query Algorithm:

- Let $X_{t}$ be a uniform sample from $H_{t}(q)$ (hash bucket of $q$ )
- Form an unbiased estimator by reweighting:



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$$
Z(y)=\sum_{t=1}^{T} \frac{w_{t}\left(\left\langle X_{t}, y\right\rangle\right)}{p_{t}\left(\left\langle X_{t}, y\right\rangle\right)}\left|H_{t}(y)\right|
$$

## Multi-resolution HBE

Data-structure: median-of-means on unbiased estimator

$$
Z(y)=\sum_{t=1}^{T} \frac{w_{t}\left(\left\langle X_{t}, y\right\rangle\right)}{p_{t}\left(\left\langle X_{t}, y\right\rangle\right)}\left|H_{t}(y)\right|
$$

that we call Multi-resolution Hashing-Based-Estimators.
Challenges:
■ Specify weighting scheme depending on convex fun $\phi$

- Select hashing schemes depending on convex fun $\phi$.

■ Provably bound the variance of the overall estimator.

## Proof Ideas

## Primer on Importance sampling

Setting: weights $w_{1}, \ldots, w_{n}$ e.g. $w_{i}=K\left(x_{i}, y\right)$,
Goal: approximate $\mu=\sum_{i=1}^{n} w_{i}$

## Importance Sampling

Black box $Q$, returns index $i$ with probability $q_{i}$

- Unbiased estimator: let $I \sim Q$ then $Z_{Q}=\frac{w}{q_{I}}$

- Variance: controlled by $\mathbb{E}\left[Z_{Q}^{2}\right]=\sum_{i=1}^{n} \frac{w_{i}^{2}}{q_{i}}$


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■ Variance: controlled by $\mathbb{E}\left[Z_{Q}^{2}\right]=\sum_{i=1}^{n} \frac{w_{i}^{2}}{q_{i}}$

## Locality Sensitive Hashing

Randomized Space Partitions $\mathbb{P}[h(x)=h(y)]=f(\|x-y\|)$


## Algorithmic Framework

Hashing-based-Estimators [Charikar, S.,FOCS'17]:
an estimator $\hat{Z}(y)$ with relative variance $O\left(\frac{1}{\sqrt{\mu}}\right)$

- Scale-free Property is hard to attain.


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■ Collision probability $p(x, y)=\Theta(\sqrt{K(x, y)})$ then one can get an estimator $\hat{Z}(y)$ with relative variance $O\left(\frac{1}{\sqrt{\mu}}\right)$.

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## Limitations of HBE

Scale-free Property is hard to attain:

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p(x, y)=\Theta(\sqrt{K(x, y)})
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## ■ Gaussian, Exponential and "polynomial" using LSH

- Collision prob. that near 0 or $\gg 1$ exhibited the desirec (exponential, gaussian or polynomial) decay with distance.
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## Main contributions

■ Generalize results on HBE to Multi-resolution HBE .
■ Distance Sensitive Hashing on the Sphere instead of LSH.
■ Approximation Theory for Log-convex functions on Sphere.

## Intuition

Given a function $w_{0}:[-1,1] \rightarrow \mathbb{R}$ want to approximate

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\int_{-1}^{1} w_{0}(\rho) d \rho
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- Find appropriate hashing probabilities $\left\{p_{t}\right\}_{t \in[T]}$.
- Design a HBE for each $w_{t}$ (Multi-resolution HBE)
- Bound the variance of resulting estimators.


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## Distance Sensitive Hashing [Aumuller et al. 2017]

$$
\begin{gathered}
g_{+}, g_{-} \sim \mathcal{N}\left(0, I_{d}\right),\left\{\left\langle x, g_{+}\right\rangle \geq \tau \wedge\left\langle x, g_{-}\right\rangle \leq-\gamma \tau\right\}, e^{O\left(\tau^{2}\right)} \text { times } \\
\log \left(p_{\gamma, \tau}(\rho)\right)=\Theta\left(-\left(\frac{1-\rho}{1+\rho}+\gamma^{2} \frac{1+\rho}{1-\rho}\right) \frac{\tau^{2}}{2}\right)
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Effect of $\gamma$


Effect of $\tau$


## Multi-resolution HBE


hashing schemes $\left\{\mathcal{H}_{t}\right\}$, coll. prob. $\left\{p_{t}\right\}$, and weight func. $\left\{w_{t}\right\}$.

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Z_{T}(y)=\frac{1}{n} \sum_{t \in[T]} \frac{w_{t}\left(X_{t}, y\right)}{p_{t}\left(X_{t}, y\right)}\left|H_{t}(y)\right|, \quad X_{t} \sim H_{t}(y) \text { for } t \in[T]
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Technique to bound variance from [Charikar, S., FOCS'17],

Key design principle $w_{t}(x, y)=\frac{p_{t}^{2}(x, y)}{\sum_{t^{\prime}} p_{t}^{2}(x, y)} \cdot w_{0}(x, y)$ results in "Variance of Multi-resolution HBE is bounded by Variance of HBE with collision probability $p_{*}(x, y)=\max _{t \in[T]\{ }\left\{p_{t}(x, y)\right\}$ " Goal: hashing scheme $p_{*}(x, y)=\Theta\left(\sqrt{w_{0}(x, y)}\right)=\Theta\left(e^{\frac{1}{2} \phi(x, y)}\right)$

Fortunately, $\frac{1}{2} \phi(x, y)$ remains convex and lipschitz.

## $p^{2}$-weighting scheme

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## Approximation of Convex Functions I

Goal, pick a set of parameters $\left\{\left(\gamma_{t}, \tau_{t}\right)\right\}_{t \in T}$ such that:

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\left|\sup _{t \in T}\left\{\log \left(p_{\gamma_{t}, \tau_{t}}(\rho)\right)\right\}-\frac{1}{2} \phi(\rho)\right|=O(1)
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Approximation Theory of Convex Functions:

- Approximate Convex Func. by $O(\sqrt{L}(\phi))$ Piecewise Linear (Sandwich Algorithm [Rote'92])
- Approximate Linear func. using $O(\log (L(\phi)))$ hash functions.
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## Approximation of Convex Functions II

Goal, pick a set of parameters $\left\{\left(\gamma_{t}, \tau_{t}\right)\right\}_{t \in T}$ such that:

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$$




## Recap




■ Partition Function Estimation via Distance Sensitive Hashing.

- Improve upon state of the art by $\sqrt{n}$ factor.

■ Multi-resolution HBE and Log-Convex Functions.

## Future Work

■ Design and implement more practical Hashing Schemes.
■ Applications in Optimization and Learning.

## Thank You!

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