

Dual First-Order and ADMM Methods for LP

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The Lagrangian Function and Method

We consider

$$f^* := \min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t.} \quad \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{x} \in \mathbf{X}. \quad (1)$$

Recall that the **Lagrangian** function:

$$L(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) - \mathbf{y}^T \mathbf{h}(\mathbf{x}).$$

and the **dual function**:

$$\phi(\mathbf{y}) = \min_{\mathbf{x} \in \mathbf{X}} L(\mathbf{x}, \mathbf{y}); \quad (2)$$

and the **dual problem**

$$(f^* \geq) \phi^* := \max_{\mathbf{y}} \phi(\mathbf{y}). \quad (3)$$

In many cases, one can find \mathbf{y}^* of dual problem (3), a **unconstrained** optimization problem; then go ahead to find \mathbf{x}^* using (2).

The Local Duality Theorem

Suppose \mathbf{x}^* is a local minimizer, and consider the **localized (convex) problem**

$$f(\mathbf{x}^*) := \min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t.} \quad \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{x} \in \mathbf{X}, \|\mathbf{x} - \mathbf{x}^*\|^2 \leq \epsilon. \quad (4)$$

Then, the **localized Lagrangian function**:

$$L_{\mathbf{x}^*}(\mathbf{x}, \mathbf{y}, \mu(\leq 0)) = f(\mathbf{x}) - \mathbf{y}^T \mathbf{h}(\mathbf{x}) - \mu(\|\mathbf{x} - \mathbf{x}^*\|^2 - \epsilon).$$

and the localized dual function:

$$\phi_{\mathbf{x}^*}(\mathbf{y}, \mu) = \min_{\mathbf{x} \in \mathbf{X}, \|\mathbf{x} - \mathbf{x}^*\|^2 \leq \epsilon} L_{\mathbf{x}^*}(\mathbf{x}, \mathbf{y}, \mu); \quad (5)$$

and the **localized dual problem**

$$\max_{\mathbf{y}, \mu \leq 0} \phi(\mathbf{y}, \mu \leq 0). \quad (6)$$

Under certain **constraint qualification and local convexity conditions**, we must have

$f(\mathbf{x}^*) = \phi(\mathbf{y}^*, \mu^* = 0)$ where the localization constraint becomes **inactive**.

The gradient and Hessian of ϕ

Let $\mathbf{x}(\mathbf{y})$ be a minimizer of (2). Then

$$\phi(\mathbf{y}) = f(\mathbf{x}(\mathbf{y})) - \mathbf{y}^T \mathbf{h}(\mathbf{x}(\mathbf{y}))$$

Thus,

$$\begin{aligned} \nabla \phi(\mathbf{y}) &= \nabla f(\mathbf{x}(\mathbf{y}))^T \nabla \mathbf{x}(\mathbf{y}) - \mathbf{y}^T \nabla \mathbf{h}(\mathbf{x}(\mathbf{y})) \nabla \mathbf{x}(\mathbf{y}) - \mathbf{h}(\mathbf{x}(\mathbf{y})) \\ &= (\nabla f(\mathbf{x}(\mathbf{y}))^T - \mathbf{y}^T \nabla \mathbf{h}(\mathbf{x}(\mathbf{y}))) \nabla \mathbf{x}(\mathbf{y}) - \mathbf{h}(\mathbf{x}(\mathbf{y})) \\ &= -\mathbf{h}(\mathbf{x}(\mathbf{y})). \end{aligned}$$

Similarly, we can derive

$$\nabla^2 \phi(\mathbf{y}) = -\nabla \mathbf{h}(\mathbf{x}(\mathbf{y})) (\nabla_{\mathbf{x}}^2 L(\mathbf{x}(\mathbf{y}), \mathbf{y}))^{-1} \nabla \mathbf{h}(\mathbf{x}(\mathbf{y}))^T,$$

where $\nabla_{\mathbf{x}}^2 L(\mathbf{x}(\mathbf{y}), \mathbf{y})$ is the Hessian of the Lagrangian function that is assumed to be positive definite at any (local) minimizer.

The Toy Example

$$\text{minimize} \quad (x_1 - 1)^2 + (x_2 - 1)^2$$

$$\text{subject to} \quad x_1 + 2x_2 - 1 = 0, \quad 2x_1 + x_2 - 1 = 0.$$

$$L(\mathbf{x}, \mathbf{y}) = (x_1 - 1)^2 + (x_2 - 1)^2 - y_1(x_1 + 2x_2 - 1) - y_2(2x_1 + x_2 - 1).$$

$$x_1 = 0.5y_1 + y_2 + 1, \quad x_2 = y_1 + 0.5y_2 + 1.$$

$$\phi(\mathbf{y}) = -1.25y_1^2 - 1.25y_2^2 - 2y_1y_2 - 2y_1 - 2y_2.$$

$$\nabla\phi(\mathbf{y}) = \begin{pmatrix} 2.5y_1 + 2y_2 + 2 \\ 2y_1 + 2.5y_2 + 2 \end{pmatrix},$$

$$\nabla^2\phi(\mathbf{y}) = - \begin{pmatrix} 1 & 2 \\ 2 & 1 \end{pmatrix} \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}^{-1} \begin{pmatrix} 1 & 2 \\ 2 & 1 \end{pmatrix}^T = - \begin{pmatrix} 2.5 & 2 \\ 2 & 2.5 \end{pmatrix}$$

The Augmented Lagrangian Function

In both theory and practice, we actually consider an **augmented** Lagrangian function (ALF)

$$L_a(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) - \mathbf{y}^T \mathbf{h}(\mathbf{x}) + \frac{\beta}{2} \|\mathbf{h}(\mathbf{x})\|^2,$$

which corresponds to an **equivalent problem** of (1):

$$f^* := \min_{\mathbf{x} \in \mathbf{X}} f(\mathbf{x}) + \frac{\beta}{2} \|\mathbf{h}(\mathbf{x})\|^2 \quad \text{s.t.} \quad \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{x} \in \mathbf{X}.$$

Note that, although at feasibility the additional square term in objective is **redundant**, it helps to improve strict convexity of the Lagrangian function.

The Augmented Lagrangian Dual

Now the dual function:

$$\phi_a(\mathbf{y}) = \min_{\mathbf{x} \in X} L_a(\mathbf{x}, \mathbf{y}); \quad (7)$$

and the dual problem

$$(f^* \geq) \phi_a^* := \max \phi_a(\mathbf{y}). \quad (8)$$

Note that the dual function satisfies $\frac{1}{\beta}$ -Lipschitz condition (see Chapter 14 of L&Y).

For the convex optimization case, say $\mathbf{h}(\mathbf{x}) = \mathbf{A}\mathbf{x} - \mathbf{b}$, we have

$$\nabla^2 L_a(\mathbf{x}, \mathbf{y}) = \nabla^2 f(\mathbf{x}) + \beta(A^T A).$$

The Augmented Lagrangian Method

The **augmented Lagrangian method** (ALM) is:

Start from any $(\mathbf{x}^0 \in X, \mathbf{y}^0)$, we compute a new iterate pair

$$\mathbf{x}^{k+1} = \arg \min_{\mathbf{x} \in X} L_a(\mathbf{x}, \mathbf{y}^k), \text{ and } \mathbf{y}^{k+1} = \mathbf{y}^k - \beta \mathbf{h}(\mathbf{x}^{k+1}).$$

The calculation of \mathbf{x} is used to compute the gradient vector of $\phi_a(\mathbf{y})$, which is a steepest **ascent** direction.

The method converges just like the SDM, because the dual function satisfies $\frac{1}{\beta}$ -**Lipschitz** condition.

Other SDM strategies may be adapted to update \mathbf{y} (the BB, ASDM, Conjugate, Quasi-Newton ...).

Analysis of the Augmented Lagrangian Method

Consider the convex optimization case $\mathbf{h}(\mathbf{x}) = \mathbf{Ax} - \mathbf{b}$. Since \mathbf{x}^{k+1} makes KKT condition:

$$\begin{aligned} \mathbf{0} &= \nabla f(\mathbf{x}^{k+1}) - A^T \mathbf{y}^k + \beta A^T (A\mathbf{x}^{k+1} - \mathbf{b}) \\ &= \nabla f(\mathbf{x}^{k+1}) - A^T (\mathbf{y}^k - \beta(A\mathbf{x}^{k+1} - \mathbf{b})) \\ &= \nabla f(\mathbf{x}^{k+1}) - A^T \mathbf{y}^{k+1}, \end{aligned}$$

we only need to be concerned about whether or not $\|A\mathbf{x}^k - \mathbf{b}\|$ converges to zero and how fast it converges. First, from the convexity of $f(\mathbf{x})$, we have

$$\begin{aligned} \mathbf{0} &\leq (\nabla f(\mathbf{x}^{k+1}) - \nabla f(\mathbf{x}^k))^T (\mathbf{x}^{k+1} - \mathbf{x}^k) \\ &= (-A^T \mathbf{y}^{k+1} + A^T \mathbf{y}^k)^T (\mathbf{x}^{k+1} - \mathbf{x}^k) \\ &= (\mathbf{y}^{k+1} - \mathbf{y}^k)^T (A\mathbf{x}^{k+1} - A\mathbf{x}^k) \\ &= -\beta(A\mathbf{x}^{k+1} - \mathbf{b})(A\mathbf{x}^{k+1} - \mathbf{b} - (A\mathbf{x}^k - \mathbf{b})), \end{aligned}$$

which implies that $\|A\mathbf{x}^{k+1} - \mathbf{b}\| \leq \|A\mathbf{x}^k - \mathbf{b}\|$, that is, the error is **non-increasing**.

Again, from the convexity, we have

$$\begin{aligned}
 \mathbf{0} &\leq (\nabla f(\mathbf{x}^{k+1}) - \nabla f(\mathbf{x}^*))^T (\mathbf{x}^{k+1} - \mathbf{x}^*) \\
 &= (A^T \mathbf{y}^{k+1} - A^T \mathbf{y}^*)^T (\mathbf{x}^{k+1} - \mathbf{x}^*) \\
 &= (\mathbf{y}^{k+1} - \mathbf{y}^*)^T (A\mathbf{x}^{k+1} - A\mathbf{x}^*) = (\mathbf{y}^{k+1} - \mathbf{y}^*)^T (A\mathbf{x}^{k+1} - \mathbf{b}) \\
 &= \frac{1}{\beta} (\mathbf{y}^{k+1} - \mathbf{y}^*)^T (\mathbf{y}^k - \mathbf{y}^{k+1}).
 \end{aligned}$$

Thus, from the positivity of the cross product, we have

$$\begin{aligned}
 \|\mathbf{y}^k - \mathbf{y}^*\|^2 &= \|\mathbf{y}^k - \mathbf{y}^{k+1} + \mathbf{y}^{k+1} - \mathbf{y}^*\|^2 \\
 &\geq \|\mathbf{y}^k - \mathbf{y}^{k+1}\|^2 + \|\mathbf{y}^{k+1} - \mathbf{y}^*\|^2 \\
 &= \beta \|A\mathbf{x}^{k+1} - \mathbf{b}\|^2 + \|\mathbf{y}^{k+1} - \mathbf{y}^*\|^2.
 \end{aligned}$$

Sum up from 0 to k of the inequality we have

$$\begin{aligned}
 \|\mathbf{y}^0 - \mathbf{y}^*\|^2 &\geq \|\mathbf{y}^{k+1} - \mathbf{y}^*\|^2 + \beta \sum_{l=0}^k \|A\mathbf{x}^{l+1} - \mathbf{b}\|^2 \\
 &\geq \beta \sum_{l=0}^k \|A\mathbf{x}^{l+1} - \mathbf{b}\|^2 \\
 &\geq (k+1)\beta \|A\mathbf{x}^{k+1} - \mathbf{b}\|^2.
 \end{aligned}$$

Two-Block Alternating Direction Method with Multipliers

For the ADMM method, we consider **structured problem**

$$\min f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \quad \text{s.t.} \quad A_1\mathbf{x}_1 + A_2\mathbf{x}_2 = \mathbf{b}, \quad \mathbf{x}_1 \in X_1, \quad \mathbf{x}_2 \in X_2.$$

Consider

$$L(\mathbf{x}_1, \mathbf{x}_2, \mathbf{y}) = f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) - \mathbf{y}^T (A_1\mathbf{x}_1 + A_2\mathbf{x}_2 - \mathbf{b}) + \frac{\beta}{2} \|A_1\mathbf{x}_1 + A_2\mathbf{x}_2 - \mathbf{b}\|^2.$$

Then, for any given $(\mathbf{x}_1^k, \mathbf{x}_2^k, \mathbf{y}^k)$, we compute a new iterate

$$\begin{aligned} \mathbf{x}_1^{k+1} &= \arg \min_{\mathbf{x}_1 \in X_1} L(\mathbf{x}_1, \mathbf{x}_2^k, \mathbf{y}^k), \\ \mathbf{x}_2^{k+1} &= \arg \min_{\mathbf{x}_2 \in X_2} L(\mathbf{x}_1^{k+1}, \mathbf{x}_2, \mathbf{y}^k), \\ \mathbf{y}^{k+1} &= \mathbf{y}^k - \beta (A_1\mathbf{x}_1^{k+1} + A_2\mathbf{x}_2^{k+1} - \mathbf{b}). \end{aligned}$$

Again, we can prove that the iterates converge with the same speed.

The ADMM method resembles the **Block Coordinate Descent (BCD)** Method ...

Direct Application of ADMM to Linear Programming I

Consider the standard-form LP

$$\begin{array}{ll}
 \text{minimize}_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\
 \text{s.t.} & A\mathbf{x} = \mathbf{b}, \\
 & \mathbf{x} \geq \mathbf{0}.
 \end{array}
 \quad \Rightarrow \quad
 \begin{array}{ll}
 \text{minimize}_{(\mathbf{x}_1, \mathbf{x}_2)} & \mathbf{c}^T \mathbf{x}_1 \\
 \text{s.t.} & A\mathbf{x}_1 = \mathbf{b}, \\
 & \mathbf{x}_1 - \mathbf{x}_2 = \mathbf{0}, \mathbf{x}_2 \geq \mathbf{0}.
 \end{array}$$

$$L(\mathbf{x}_1, \mathbf{x}_2, \mathbf{y}) = \mathbf{c}^T \mathbf{x}_1 - \mathbf{y}^T (A\mathbf{x}_1 - \mathbf{b}) - \mathbf{s}^T (\mathbf{x}_1 - \mathbf{x}_2) + \frac{\beta}{2} (\|A\mathbf{x}_1 - \mathbf{b}\|^2 + \|\mathbf{x}_1 - \mathbf{x}_2\|^2).$$

where \mathbf{y} and \mathbf{s} are the multiplier vectors of first and second equality constraints in the reformulation.

The advantage of such splitting reformulation is that the update of either \mathbf{x}_1 or \mathbf{x}_2 has a simple close form solution.

Direct Application of ADMM to Dual Linear Programming I

Consider the dual LP

$$\begin{aligned} & \text{maximize}_{(\mathbf{y}, \mathbf{s})} && \mathbf{b}^T \mathbf{y} \\ & \text{s.t.} && A^T \mathbf{y} + \mathbf{s} = \mathbf{c}, \mathbf{s} \geq \mathbf{0}. \end{aligned}$$

The augmented Lagrangian function would be

$$L(\mathbf{y}, \mathbf{s}, \mathbf{x}) = -\mathbf{b}^T \mathbf{y} - \mathbf{x}^T (A^T \mathbf{y} + \mathbf{s} - \mathbf{c}) + \frac{\beta}{2} \|A^T \mathbf{y} + \mathbf{s} - \mathbf{c}\|^2,$$

where β is a positive parameter, and \mathbf{x} is the multiplier vector.

Direct Application of ADMM to Dual Linear Programming II

The ADMM for the dual is straightforward: starting from any \mathbf{y}^0 , $\mathbf{s}^0 \geq \mathbf{0}$, and multiplier \mathbf{x}^0 ,

- Update variable \mathbf{y} :

$$\mathbf{y}^{k+1} = \arg \min_{\mathbf{y}} L(\mathbf{y}, \mathbf{s}^k, \mathbf{x}^k);$$

- Update slack variable \mathbf{s} :

$$\mathbf{s}^{k+1} = \arg \min_{\mathbf{s} \geq \mathbf{0}} L(\mathbf{y}^{k+1}, \mathbf{s}, \mathbf{x}^k);$$

- Update multipliers \mathbf{x} :

$$\mathbf{x}^{k+1} = \mathbf{x}^k - \beta(A^T \mathbf{y}^{k+1} + \mathbf{s}^{k+1} - \mathbf{c}).$$

Note that the updates of \mathbf{y} is a **least-squares problem** with constant matrix, and the update of \mathbf{s} has a **simple close form**. (Also note that \mathbf{x} would be non-positive at the end, since we changed maximization to minimization of the dual.)

To split \mathbf{y} into **multi blocks** and update cyclically in random order?

Direct Application of ADMM to Dual Linear Programming III

One can also consider to reformulate the dual as

$$\begin{aligned}
 & \text{maximize}_{\mathbf{y}, \mathbf{s}, \mathbf{u}_1, \mathbf{u}_2} && \mathbf{b}^T \mathbf{y} \\
 & \text{s.t.} && A_1^T \mathbf{y}_1 - \mathbf{u}_1 = \mathbf{0}, \quad (\mathbf{v}_1) \\
 & && A_2^T \mathbf{y}_2 - \mathbf{u}_2 = \mathbf{0}, \quad (\mathbf{v}_2) \\
 & && \mathbf{u}_1 + \mathbf{u}_2 + \mathbf{s} = \mathbf{c}, \quad (\mathbf{x}) \\
 & && \mathbf{s} \geq \mathbf{0};
 \end{aligned} \tag{9}$$

with the multiplier \mathbf{v}_1 , \mathbf{v}_2 and \mathbf{x} for the three sets of the equality constraints.

$$\begin{aligned}
 & L^d(\mathbf{y}_1, \mathbf{y}_2, \mathbf{u}_1, \mathbf{u}_2, \mathbf{s}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{x}) \\
 & = -\mathbf{b}_1^T \mathbf{y}_1 - \mathbf{b}_2^T \mathbf{y}_2 - \mathbf{v}_1^T (A_1^T \mathbf{y}_1 - \mathbf{u}_1) - \mathbf{v}_2^T (A_2^T \mathbf{y}_2 - \mathbf{u}_2) - \mathbf{x}^T (\mathbf{u}_1 + \mathbf{u}_2 + \mathbf{s} - \mathbf{c}) \\
 & \quad + \frac{\beta}{2} (\|A_1^T \mathbf{y}_1 - \mathbf{u}_1\|^2 + \|A_2^T \mathbf{y}_2 - \mathbf{u}_2\|^2 + \|\mathbf{u}_1 + \mathbf{u}_2 + \mathbf{s} - \mathbf{c}\|^2).
 \end{aligned} \tag{10}$$

Note that \mathbf{y}_i , $i = 1, 2$, and $\mathbf{s} \geq \mathbf{0}$ can be updated independently and in parallel, and \mathbf{u}_i , $i = 1, 2$, can be updated jointly with a close form(?). This is essentially a two-block ADMM!

Direct Application of ADMM to Dual LP IV: Barrier Regularization

Now consider dual linear program with the logarithmic barrier function

$$\begin{aligned} & \text{maximize}_{\mathbf{y}, \mathbf{s}} && \mathbf{b}^T \mathbf{y} + \mu \sum_j \ln(s_j) \\ & \text{s.t.} && A^T \mathbf{y} + \mathbf{s} = \mathbf{c}, \end{aligned} \tag{11}$$

where μ is a fixed small positive constant.

The augmented Lagrangian function would be

$$L_\mu(\mathbf{y}, \mathbf{s}, \mathbf{x}) = -\mathbf{b}^T \mathbf{y} - \mu \sum_j \ln(s_j) - \mathbf{x}^T (A^T \mathbf{y} + \mathbf{s} - \mathbf{c}) + \frac{\beta}{2} \|A^T \mathbf{y} + \mathbf{s} - \mathbf{c}\|^2,$$

Apply the path-following idea to the Dual ADMM with barrier.

“An ADMM-Based Interior-Point Method for Large-Scale Linear Programming,”

(<https://arxiv.org/abs/1805.12344>).

ADMM for Multi-block Convex Minimization

Why not consider convex minimization problems with *three blocks*:

$$\begin{aligned}
 \min \quad & f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) + f_3(\mathbf{x}_3) \\
 \text{s.t.} \quad & A_1\mathbf{x}_1 + A_2\mathbf{x}_2 + A_3\mathbf{x}_3 = \mathbf{b} \\
 & \mathbf{x}_1 \in \mathcal{X}_1, \mathbf{x}_2 \in \mathcal{X}_2, \mathbf{x}_3 \in \mathcal{X}_3
 \end{aligned}$$

The direct and natural extension of ADMM with null objectives:

$$\left\{ \begin{aligned}
 \mathbf{x}_1^{k+1} &= \arg \min \{ L(\mathbf{x}_1, \mathbf{x}_2^k, \mathbf{x}_3^k, \mathbf{y}^k) \mid \mathbf{x}_1 \in \mathcal{X}_1 \} \\
 \mathbf{x}_2^{k+1} &= \arg \min \{ L(\mathbf{x}_1^{k+1}, \mathbf{x}_2, \mathbf{x}_3^k, \mathbf{y}^k) \mid \mathbf{x}_2 \in \mathcal{X}_2 \} \\
 \mathbf{x}_3^{k+1} &= \arg \min \{ L(\mathbf{x}_1^{k+1}, \mathbf{x}_2^{k+1}, \mathbf{x}_3, \mathbf{y}^k) \mid \mathbf{x}_3 \in \mathcal{X}_3 \} \\
 \mathbf{y}^{k+1} &= \mathbf{y}^k - \beta(A_1\mathbf{x}_1^{k+1} + A_2\mathbf{x}_2^{k+1} + A_3\mathbf{x}_3^{k+1} - \mathbf{b})
 \end{aligned} \right.$$

$$L(\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3, \mathbf{y}) = \sum_{i=1}^3 f_i(\mathbf{x}_i) - \mathbf{y}^T \left(\sum_{i=1}^3 A_i \mathbf{x}_i - \mathbf{b} \right) + \frac{\beta}{2} \left\| \sum_{i=1}^3 A_i \mathbf{x}_i - \mathbf{b} \right\|^2$$

Divergent Example of the Extended ADMM I

Should it converge? (Not easy to analyze the convergence of ADMM with more than two blocks; or the proving operator theory of two-block cannot be directly extended to the ADMM with three blocks.)

Consider the system of homogeneous linear equations with three variables and null objective functions:

$$A_1x_1 + A_2x_2 + A_3x_3 = \mathbf{0}, \text{ where } A = (A_1, A_2, A_3) = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 2 \\ 1 & 2 & 2 \end{pmatrix}.$$

Then the extended ADMM with $\beta = 1$ can be specified as a linear map

$$\begin{pmatrix} 3 & 0 & 0 & 0 & 0 & 0 \\ 4 & 6 & 0 & 0 & 0 & 0 \\ 5 & 7 & 9 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 2 & 0 & 1 & 0 \\ 1 & 2 & 2 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} x_1^{k+1} \\ x_2^{k+1} \\ x_3^{k+1} \\ \mathbf{y}^{k+1} \end{pmatrix} = \begin{pmatrix} 0 & -4 & -5 & 1 & 1 & 1 \\ 0 & 0 & -7 & 1 & 1 & 2 \\ 0 & 0 & 0 & 1 & 2 & 2 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} x_1^k \\ x_2^k \\ x_3^k \\ \mathbf{y}^k \end{pmatrix}.$$

Divergent Example of the Extended ADMM II

Or equivalently,

$$\begin{pmatrix} x_2^{k+1} \\ x_3^{k+1} \\ \mathbf{y}^{k+1} \end{pmatrix} = M \begin{pmatrix} x_2^k \\ x_3^k \\ \mathbf{y}^k \end{pmatrix},$$

where

$$M = \frac{1}{162} \begin{pmatrix} 144 & -9 & -9 & -9 & 18 \\ 8 & 157 & -5 & 13 & -8 \\ 64 & 122 & 122 & -58 & -64 \\ 56 & -35 & -35 & 91 & -56 \\ -88 & -26 & -26 & -62 & 88 \end{pmatrix}.$$

Divergent Example of the Extended ADMM III

The matrix $M = V \text{Diag}(d) V^{-1}$, where $d = \begin{pmatrix} 0.9836 + 0.2984i \\ 0.9836 - 0.2984i \\ 0.8744 + 0.2310i \\ 0.8744 - 0.2310i \\ 0 \end{pmatrix}$. Note that

$$\rho(M) = |d_1| = |d_2| > 1.$$

Theorem 1 *There existing an example where the direct extension of ADMM of three blocks with a real number initial point is not necessarily convergent for any choice of β . Moreover, for any randomly generated initial point, ADMM diverges with probability 1.*

Multi-block problems and ADMM

In general, consider a convex optimization problem

$$\begin{aligned} \min_{\mathbf{x} \in R^N} \quad & f_1(\mathbf{x}_1) + \cdots + f_n(\mathbf{x}_n), \\ \text{subject to} \quad & A\mathbf{x} \triangleq A_1\mathbf{x}_1 + \cdots + A_n\mathbf{x}_n = \mathbf{b}, \\ & \mathbf{x}_i \in \mathcal{X}_i \subset R^{d_i}, \quad i = 1, \dots, n. \end{aligned} \tag{12}$$

$$L(\mathbf{x}_1, \dots, \mathbf{x}_n; \mathbf{y}) = \sum_i f_i(x_i) - \mathbf{y}^T \left(\sum_i A_i \mathbf{x}_i - \mathbf{b} \right) + \frac{\beta}{2} \left\| \sum_i A_i \mathbf{x}_i - \mathbf{b} \right\|^2$$

The direct Cyclic Extension Multi-block ADMM:

$$\begin{cases} \mathbf{x}_1 \longleftarrow \arg \min_{\mathbf{x}_1 \in \mathcal{X}_1} L(\mathbf{x}_1, \dots, \mathbf{x}_n; \mathbf{y}), \\ \quad \vdots \\ \mathbf{x}_n \longleftarrow \arg \min_{\mathbf{x}_n \in \mathcal{X}_n} L(\mathbf{x}_1, \dots, \mathbf{x}_n; \mathbf{y}), \\ \mathbf{y} \longleftarrow \mathbf{y} - \beta(A\mathbf{x} - \mathbf{b}), \end{cases}$$

How to Make it Work?

There are many “correction” methods to deal with the problem, but ...

Is there a “simple way” to make the ADMM with the multi-block work?

Permute the updating order of \mathbf{x}_i randomly, and it works for the example – the expected $\rho(M)$ equals **0.9723!**

Randomly Permuted ADMM

Random-Permuted ADMM (RP-ADMM): each round, draw a random permutation $\sigma = (\sigma(1), \dots, \sigma(n))$ of $\{1, \dots, n\}$, and

$$\text{Update } \mathbf{x}_{\sigma(1)} \rightarrow \mathbf{x}_{\sigma(2)} \rightarrow \dots \rightarrow \mathbf{x}_{\sigma(n)} \rightarrow \mathbf{y}.$$

(This is sample without replacement.)

Interpretation: Force “absolute fairness” among blocks.

Simulation Test Result on solving linear equations: always converges!

Any theory behind the success?

We produced a positive result for ADMM on solving the system of linear equations.

Random Permuted ADMM for Linear Systems

- Consider solving any square system of linear equation ($f_i = 0, \forall i$).

$$\begin{aligned} \text{(S)} \quad & \min_{\mathbf{x} \in R^N} \quad 0, \\ & \text{s.t.} \quad A_1 \mathbf{x}_1 + \cdots + A_n \mathbf{x}_n = \mathbf{b}, \end{aligned}$$

where $A = [A_1, \dots, A_n] \in R^{N \times N}$ is full-rank, $\mathbf{x}_i \in R^{d_i}$ and $\sum_i d_i = N$.

- RP-ADMM: Pick a permutation σ of $\{1, \dots, n\}$ uniformly at random, then compute (with $\beta = 1$) $\mathbf{x}_{\sigma(1)}^{k+1}, \dots, \mathbf{x}_{\sigma(n)}^{k+1}, \mathbf{y}^{k+1}$ by

$$\begin{cases} -A_{\sigma(1)}^T \mathbf{y}^k + A_{\sigma(1)}^T (A_{\sigma(1)} \mathbf{x}_{\sigma(1)}^{k+1} + \sum_{l=2}^n A_{\sigma(l)} \mathbf{x}_{\sigma(l)}^k - \mathbf{b}) = \mathbf{0}, \\ \dots \\ -A_{\sigma(n)}^T \mathbf{y}^k + A_{\sigma(n)}^T (\sum_{j=1}^{n-1} A_{\sigma(j)} \mathbf{x}_{\sigma(j)}^{k+1} + A_{\sigma(n)} \mathbf{x}_{\sigma(n)}^{k+1} - \mathbf{b}) = \mathbf{0}, \\ \mathbf{y}^{k+1} = \mathbf{y}^k - (\sum_{i=1}^n A_i \mathbf{x}_i^{k+1} - \mathbf{b}) \end{cases}$$

- WOLG, assume $\mathbf{b} = \mathbf{0}$.

Main Result: Convergence in Expectation

- After k rounds, RP-ADMM generates \mathbf{z}^k , an r.v. depending on

$$\boldsymbol{\xi}_k = (\sigma_1, \dots, \sigma_k), \quad \mathbf{z}^i = M_{\sigma_i} \mathbf{z}^{i-1}, \quad i = 1, \dots, k,$$

where σ_i is the picked permutation at i -th round.

- Denote the expected output $\boldsymbol{\phi}^k \triangleq E_{\boldsymbol{\xi}_k}(\mathbf{z}^k)$

Theorem 2 *The expected output converges to the unique solution, i.e.*

$$\{\boldsymbol{\phi}^k\}_{k \rightarrow \infty} \longrightarrow \begin{bmatrix} \mathbf{0} \\ \mathbf{0} \end{bmatrix}.$$

- **Remark:** Expected convergence \neq convergence, but is a strong evidence for convergence for solving most problems, e.g., when iterates are bounded.

The Average Mapping is a Contraction

- The update equation of RP-ADMM for (S) is

$$\mathbf{z}^{k+1} = M_{\sigma} \mathbf{z}^k,$$

where $M_{\sigma} \in \mathbb{R}^{2N \times 2N}$ depend on σ .

- Define the expected update matrix as

$$M = E_{\sigma}(M_{\sigma}) = \frac{1}{n!} \sum_{\sigma} M_{\sigma}.$$

Theorem 3 *The spectral radius of M , $\rho(M)$, is strictly less than 1.*

- Remark: For A in the divergence example, $\rho(M_{\sigma}) > 1$ for any σ
 - Averaging Helps, a lot.

Math Problem of Theorem 3

- Define

$$Q \triangleq E(L_\sigma^{-1}) = \frac{1}{n!} \sum_{\sigma} L_\sigma^{-1}. \quad (13)$$

- Example:

$$L_{(231)} = \begin{bmatrix} 1 & A_1^T A_2 & A_1^T A_3 \\ 0 & 1 & 0 \\ 0 & A_3^T A_2 & 1 \end{bmatrix}.$$

- Need to prove that, for all A , $\rho(M) < 1$ where

$$M = \begin{bmatrix} I - QA^T A & QA^T \\ -A + AQA^T A & I - AQA^T \end{bmatrix}.$$

Difficulties of Proving Theorem 3

- **Difficulty 1:** Few tools deal with spectral radius of non-symmetric matrices.
 - E.g. $\rho(X + Y) \leq \rho(X) + \rho(Y)$ and $\rho(XY) \leq \rho(X)\rho(Y)$ don't hold.
 - Though $\rho(M) < \|M\|$, it turns out $\|M\| > 2.3$ for the counterexample.
- **Difficulty 2:** M is a complicated function of A .
 - $n = 3$, let $(A^T A)_{k,l} = b_{kl}$, then $Q_{12} = -\frac{1}{2}b_{12} + \frac{1}{6}b_{13}b_{23}$.
 - $n = 4$, $Q_{12} = -\frac{1}{2!}b_{12} + \frac{1}{3!}(b_{13}b_{32} + b_{14}b_{42}) - \frac{1}{4!}(b_{13}b_{34}b_{42} + b_{14}b_{43}b_{32})$.
- **Solution:** Symmetrization and Mathematical Induction.

Two Main Lemmas to Prove Theorem 3

- **Step 1:** Relate M to a symmetric matrix AQA^T .

Lemma 1

$$\mathbf{y} \in \text{eig}(M) \iff \frac{(1 - \mathbf{y})^2}{1 - 2\mathbf{y}} \in \text{eig}(AQA^T).$$

Since Q defined by (13) is *symmetric*, we have

$$\rho(M) < 1 \iff \text{eig}(AQA^T) \subseteq (0, \frac{4}{3}).$$

- **Step 2:** Bound eigenvalues of AQA^T - prove by math induction.

Lemma 2

$$\text{eig}(AQA^T) \subseteq (0, \frac{4}{3}).$$

- Remark: $4/3$ is “almost” tight; for $n = 3$, maximum ≈ 1.18 . Increase to $4/3$ as n increases.

More Randomization: Randomly Assembled ADMM

Randomly Assembled ADMM (RA-ADMM) – Random Variable Sampling without Replacement in Each ADMM round:

1. Set the initial set N_x as all (primal) decision variables.
2. Randomly select a subset of variables from N_x to optimize.
3. Remove this set of variables from N_x and return to Step 1 till N_x is empty.
4. Update the dual multipliers as usual.

Provide much better results for non-convex and discrete/combinatorial optimization (Mihic et al. 2020).

Extensions and Research Directions

- Non-square system of linear equations: resolved
- Non-separable convex quadratic minimization: resolved
- Theory: Convergence w.h.p.?
- Theory: Generalize to inequality systems or convex optimization at large?
- Theory: Overall complexity of Interior-Point ADMM for LP?
- Theory: More analyses on RA-ADMM?
- Implementation and computation development!