

# MS&E 310 Course Project II: Markov Decision Process

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Markov Decision Processes (MDPs) provide a mathematical framework for modeling sequential decision-making in situations where outcomes are partly random and partly under the control of a decision maker. MDPs are useful for studying a wide range of optimization problems solved via Dynamic Programming (DP), where it was known at least as early as the 1950s (cf. Shapley 1953, Bellman 1957). Modern applications include dynamic planning, reinforcement learning, social networking, and almost all other dynamic/sequential decision game strategy making problems in Mathematical, Physical, Management and Social Sciences.

As talked in class, the MDP problem with  $m$  states and total  $n$  actions can be formulated as a standard form linear program with  $m$  equality constraints and  $n$  variables:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \sum_{j \in \mathcal{A}_1} c_j x_j + \dots + \sum_{j \in \mathcal{A}_m} c_j x_j \\ \text{s.t.} \quad & \sum_{j \in \mathcal{A}_1} (\mathbf{e}_1 - \gamma \mathbf{p}_j) x_j + \dots + \sum_{j \in \mathcal{A}_m} (\mathbf{e}_m - \gamma \mathbf{p}_j) x_j = \mathbf{e}, \\ & \dots \quad x_j \quad \dots \quad \geq 0, \forall j, \end{aligned} \tag{1}$$

where  $\mathcal{A}_i$  represents the set of all actions available in state  $i$ ,  $\mathbf{p}_j$  is the state transition probabilities from state  $i$  to all states and  $c_j$  is the immediate cost when action  $j$  is taken, and  $0 < \gamma < 1$  is the discount factor. Also,  $\mathbf{e} \in R^m$  is the vector of ones, and  $\mathbf{e}_i$  is the unit vector with 1 at the  $i$ -th position and zeros everywhere else. Variable  $x_j$ ,  $j \in \mathcal{A}_i$ , is the state-action frequency or flux, or the expected present value of the number of times in which the process visits state  $i$  and takes state-action  $j \in \mathcal{A}_i$ . Thus, solving the problem entails choosing a state-action frequencies/fluxes that minimize the expected present value sum of total costs. The dual of the LP is

$$\begin{aligned} \text{maximize}_{\mathbf{y}} \quad & \mathbf{e}^T \mathbf{y} = \sum_{i=1}^m y_i \\ \text{subject to} \quad & y_1 - \gamma \mathbf{p}_j^T \mathbf{y} \leq c_j, \quad j \in \mathcal{A}_1 \\ & \vdots \\ & y_i - \gamma \mathbf{p}_j^T \mathbf{y} \leq c_j, \quad j \in \mathcal{A}_i \\ & \vdots \\ & y_m - \gamma \mathbf{p}_j^T \mathbf{y} \leq c_j, \quad j \in \mathcal{A}_m. \end{aligned} \tag{2}$$

where  $y_i$  represents the cost-to-go value in state  $i$ .

**Question 1:** Prove that in (1) every basic feasible solution represent a policy, i.e., the basic variables have exactly one variable from each state  $i$ . Furthermore, prove each basic variable value is no less than 1, and the sum of all basic variable values is  $\frac{m}{1-\gamma}$ .

**Question 2:** Value Iteration Method: Starting with any vector  $\mathbf{y}^0$ , then iteratively update it

$$y_i^{k+1} = \min_{j \in \mathcal{A}_i} \{c_j + \gamma \mathbf{p}_j^T \mathbf{y}^k\}, \forall i. \quad (3)$$

Prove the contraction result:

$$\|\mathbf{y}^{k+1} - \mathbf{y}^*\|_\infty \leq \gamma \|\mathbf{y}^k - \mathbf{y}^*\|_\infty, \forall k.$$

where  $\mathbf{y}^*$  is the fixed-point or optimal value vector, that is,

$$y_i^* = \min_{j \in \mathcal{A}_i} \{c_j + \gamma \mathbf{p}_j^T \mathbf{y}^*\}, \forall i.$$

**Question 3:** In the VI method, if starting with any vector  $\mathbf{y}^0 \geq \mathbf{y}^*$ , then prove the following entry-wise monotone property:

$$\mathbf{y}^* \leq \mathbf{y}^{k+1} \leq \mathbf{y}^k, \forall k.$$

This property has been used in a recent paper (see [SWWY17]) on the VI method using samples.

**Question 4:** Rather than go through all state values in each iteration, we modify the VI method, call it RandomVI: In the  $k$ th iteration, randomly select a subset of states  $B^k$  and do

$$y_i^{k+1} = \min_{j \in \mathcal{A}_i} \{c_j + \gamma \mathbf{p}_j^T \mathbf{y}^k\}, \forall i \in B^k. \quad (4)$$

In RandomVI, we only update a subset of state values at random in each iteration.

What can you tell the convergence of the RandomVI method? Does it make a difference with the classical VI method? How is the sample size affect the performance? Use simulated computational experiments to verify your claims.

Suppose we build an empirical distribution for each action being selected as the winning action in the final policy: the probability of action  $j$  is the past frequency of action  $j$  is being selected as the arg min in the previous iterations, e.g., the the Bayes update where we start with a uniform distribution  $\tilde{p}^0$ . Redo the computational experiments by randomly selecting  $B^k$  using the empirical distribution.

**Question 5:** Here is another modification, called CyclicVI: In the  $k$ th iteration do

- Initialize  $\tilde{\mathbf{y}}^k = \mathbf{y}^k$ .
- For  $i = 1$  to  $m$

$$\tilde{y}_i^k = \min_{j \in \mathcal{A}_i} \{c_j + \gamma \mathbf{p}_j^T \tilde{\mathbf{y}}^k\} \quad (5)$$

- $\mathbf{y}^{k+1} = \tilde{\mathbf{y}}^k$ .

In the CyclicVI method, as soon as a state value is updated, we use it to update the rest of state values.

What can you tell the convergence of the CyclicVI method? Does it make a difference with other VI methods? Use simulated computational experiments to verify your claims.

**Question 6:** In the CyclicVI method, rather than with the fixed cycle order from 1 to  $m$ , we follow a random permutation order, or sample without replacement to update the state values. More precisely, in the  $k$ th iteration do

0. Initialize  $\tilde{\mathbf{y}}^k = \mathbf{y}^k$  and  $B^k = \{1, 2, \dots, m\}$

1. – Randomly select  $i \in B^k$

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$$\tilde{y}_i^k = \min_{j \in \mathcal{A}_i} \{c_j + \gamma \mathbf{p}_j^T \tilde{\mathbf{y}}^k\} \quad (6)$$

– remove  $i$  from  $B^k$  and return to Step 1.

3.  $\mathbf{y}^{k+1} = \tilde{\mathbf{y}}^k$ .

We call it the randomly permuted CyclicVI or RPCyclicVI in short

What can you tell the convergence of the RPCyclicVI method? Does it compare with other VI methods? Use simulated computational experiments to verify your claims.

## References

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