

A Brief Introduction to Lexical Functional Grammar

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LING233B

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* Based on material by
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Architectural Issues

Representation:

Formal encoding of linguistic dependencies

Modularity:

Factoring independent generalizations

Mathematical tractability:

Provable mathematical properties

Implementability:

Transparent, efficient computation

LFG: Formal Concepts

- **Structure**
- **Structural Description**
- **Structural Correspondence (Projection)**

Architectural themes:

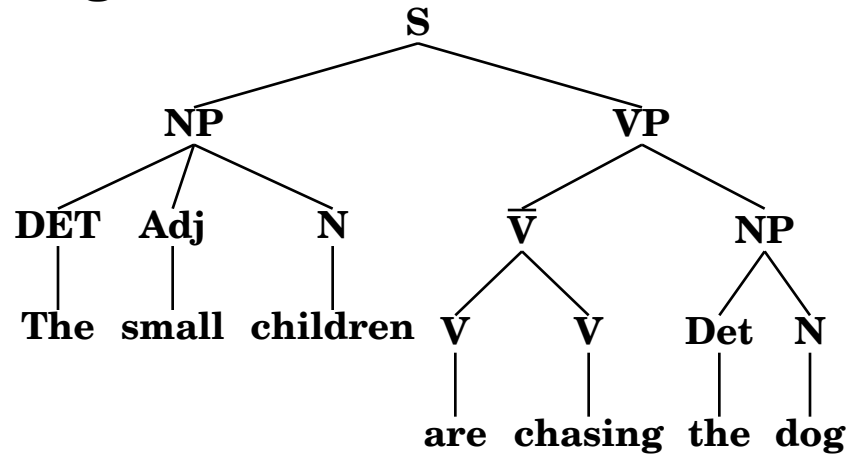
Importance of partial specification

Factor generalizations on related but dissimilar representations

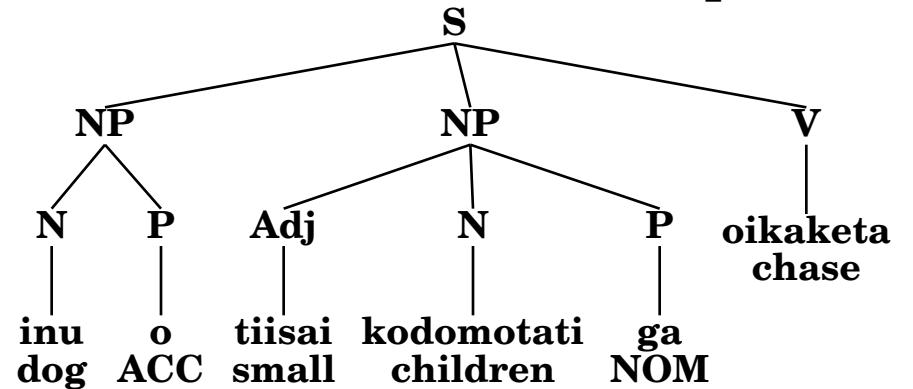
Variation in External Structure

“The small children are chasing the dog.”

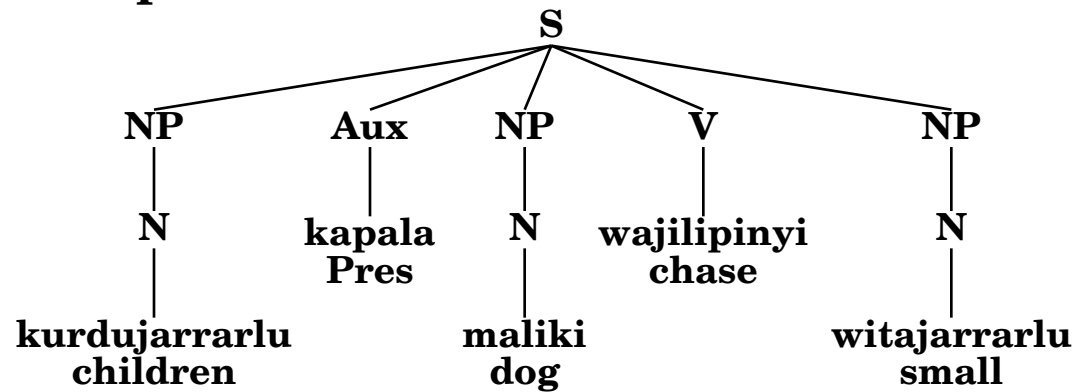
English:



Japanese:

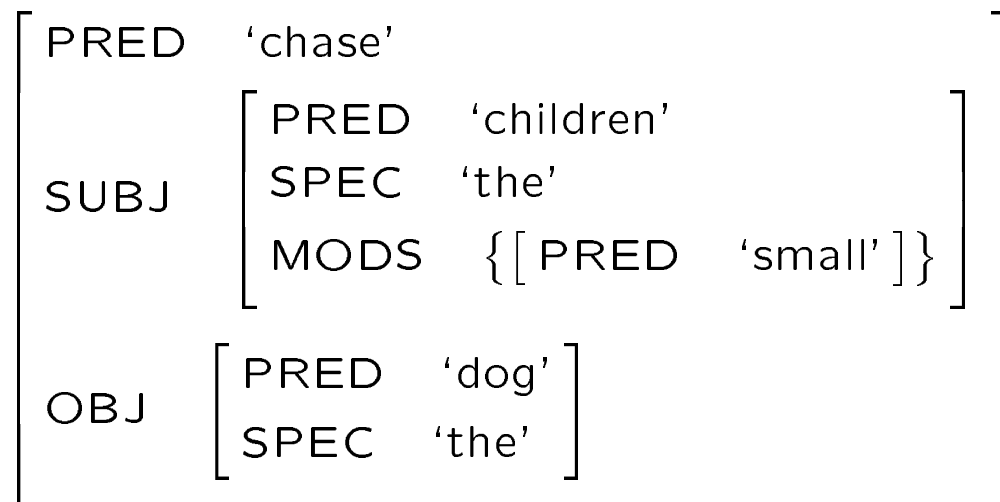


Warlpiri:



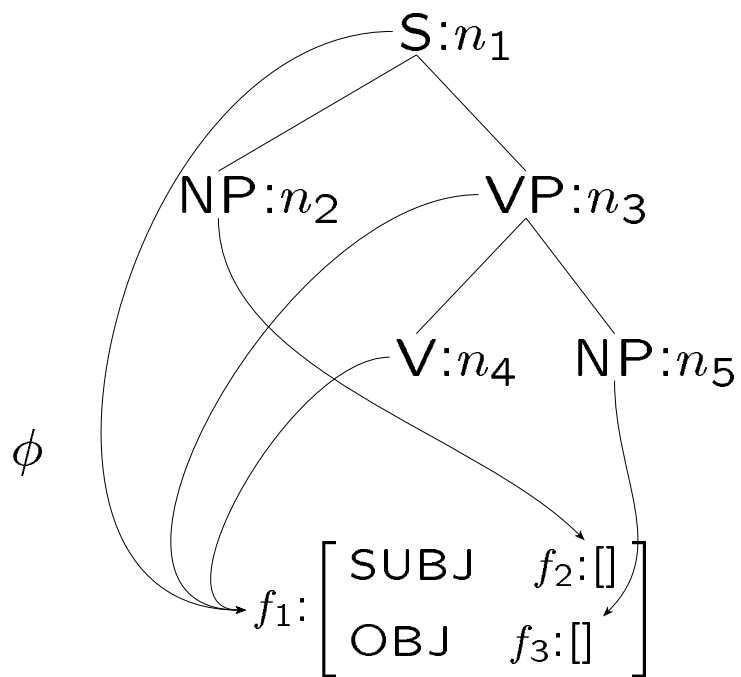
Internal structural invariance:

F-structure for
English, Japanese, and Warlpiri
'The small children are chasing the dog':

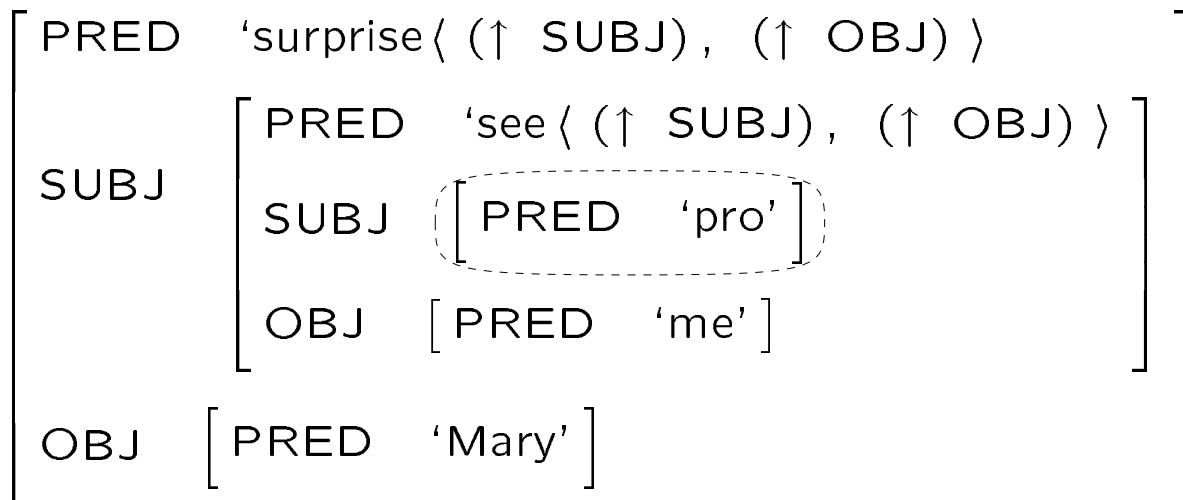
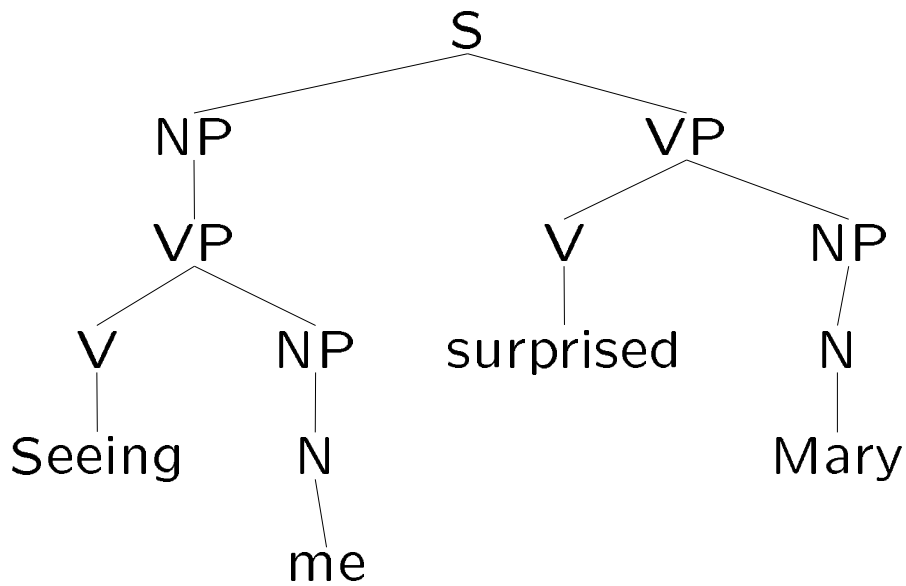


C-structure/f-structure mapping:
Many-to-one

$$\phi: N \rightarrow F$$



C-structure/f-structure mapping: Not 'onto'

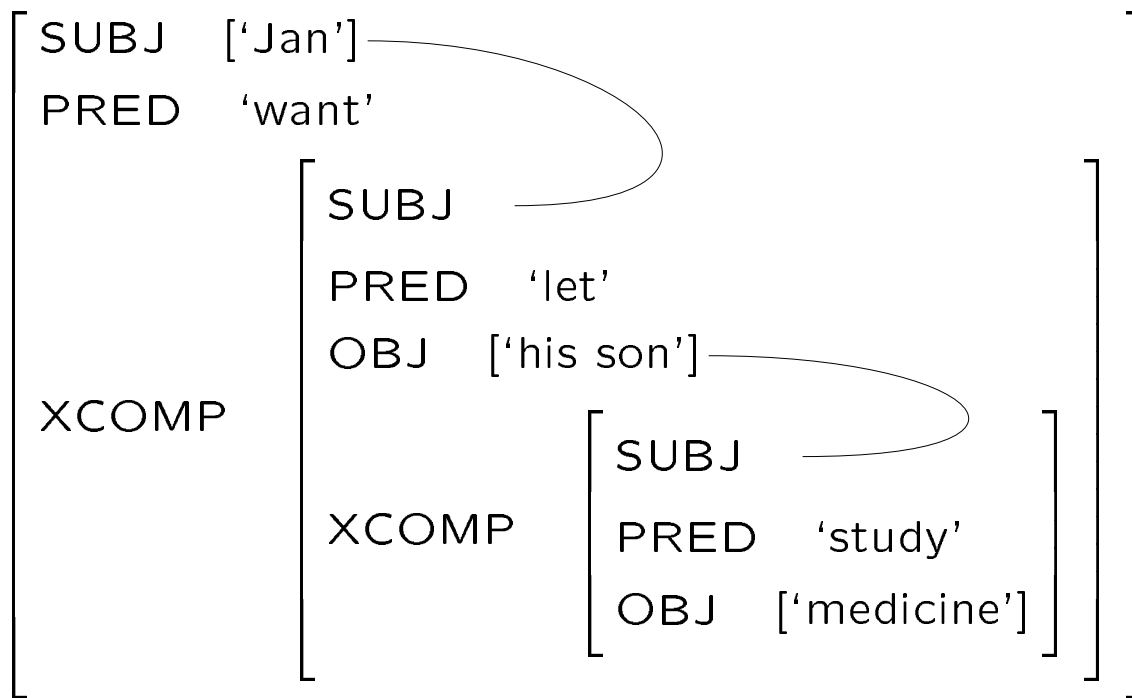
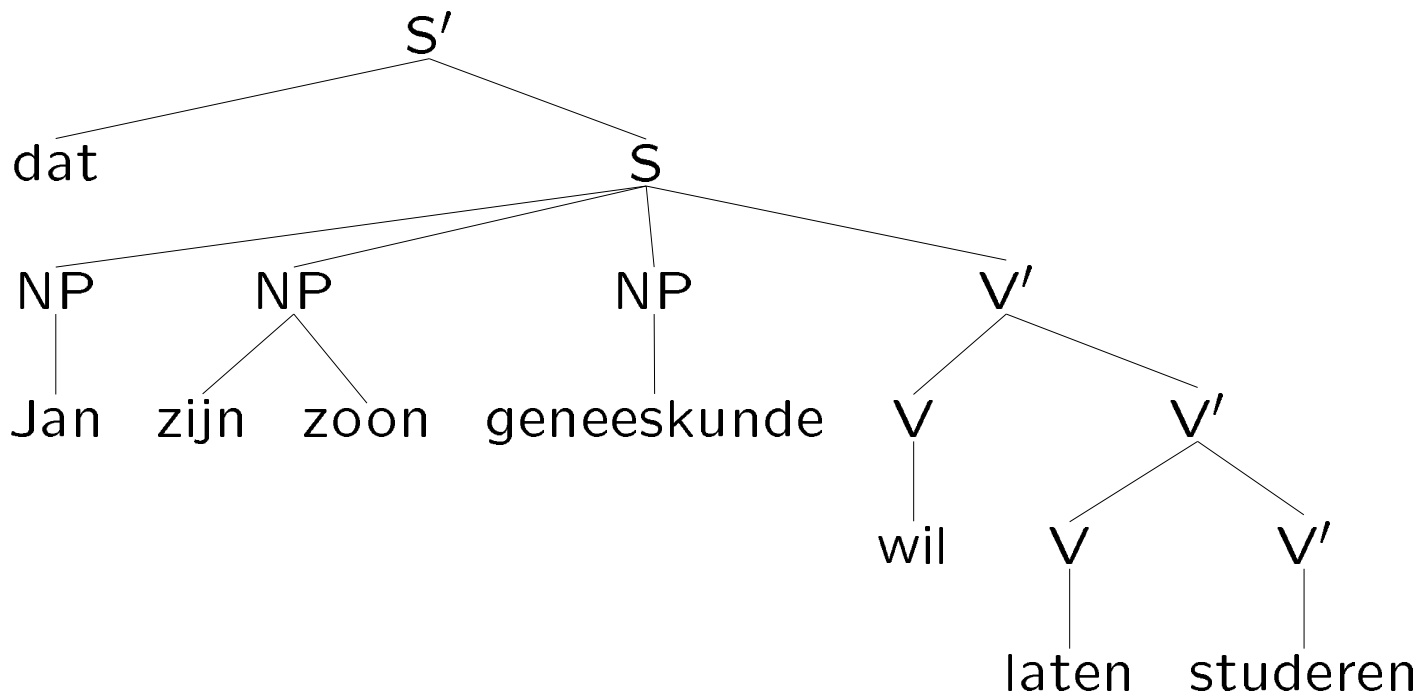


Non-native relationships induced by
correspondence:
Functional precedence

$$f_1 <_f f_2 \text{ iff}$$

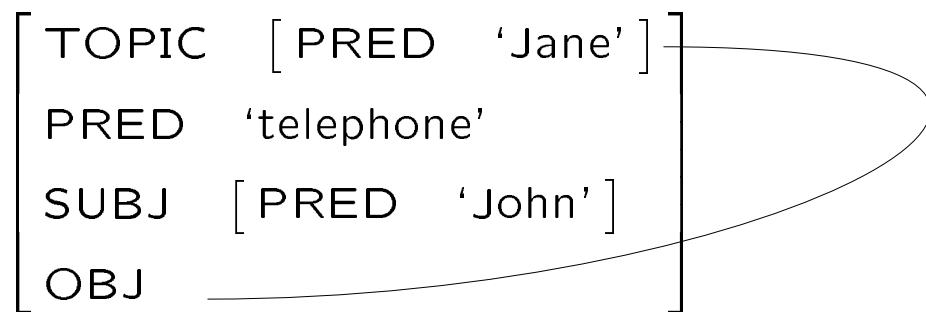
for all $n_1 \in \phi^{-1}(f_1)$ and for all $n_2 \in \phi^{-1}(f_2)$,

$$n_1 <_c n_2$$

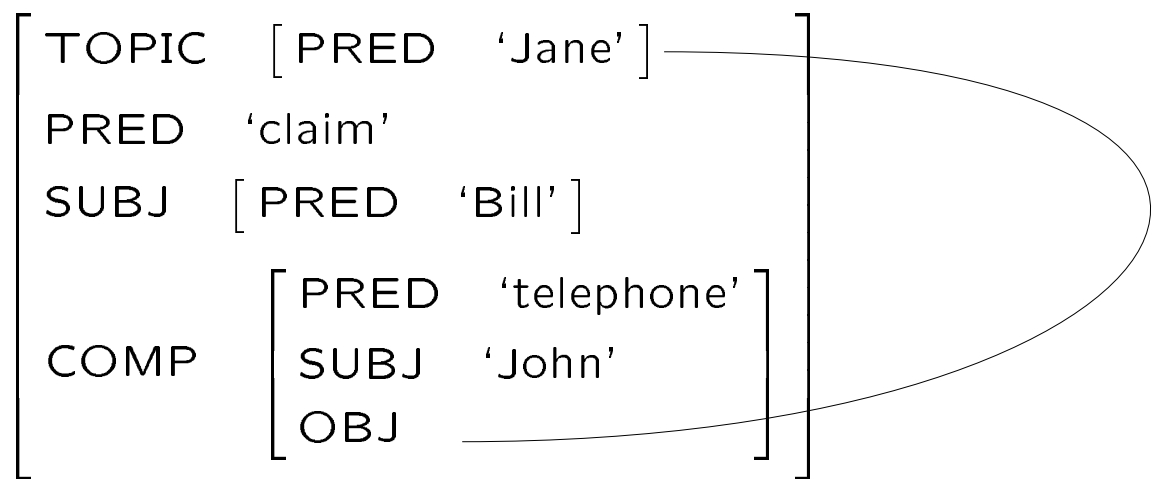


F-Structural Relations: Functional Uncertainty

Jane, John telephoned yesterday.



Jane, Bill claimed that John telephoned
yesterday.



Jane, Bill claimed that Sue said that John
telephoned yesterday.

$(\uparrow \text{ TOPIC}) = (\uparrow \text{ COMP}^* \text{ OBJ})$

$(f1 \text{ SUBJ}) = (f1 \text{ GF}^* \text{ OBJ})$

GF* is COMP COMP COMP

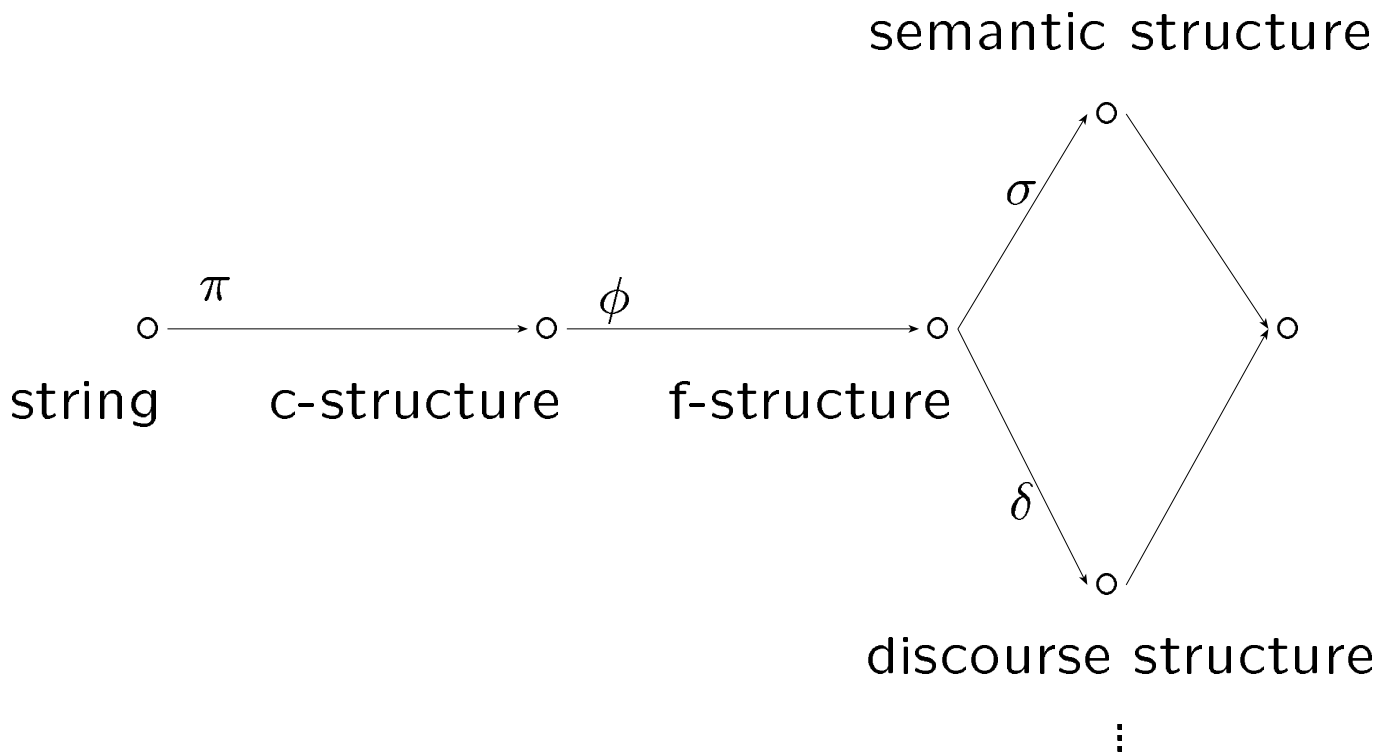
$f6 = f5$

1: $\left[\begin{array}{l} \text{SUBJ } f6 \\ \text{COMP } 2: \left[\begin{array}{l} \text{SUBJ } f7 \\ \text{COMP } 3: \left[\begin{array}{l} \text{SUBJ } f8 \\ \text{COMP } 4: \left[\begin{array}{l} \text{SUBJ } f9 \\ \text{OBJ } f5 \end{array} \right] \end{array} \right] \end{array} \right] \end{array} \right]$

Beyond Syntax: Other Structures

Form

Meaning



Quick introduction to LFG

Annotated phrase structure rules:

$$\text{IP} \longrightarrow \text{NP} \quad \text{I}' \\ (\uparrow \text{ SUBJ}) = \downarrow \quad \uparrow = \downarrow$$

$$\text{I}' \longrightarrow \left(\begin{array}{c} \text{I} \\ \uparrow = \downarrow \end{array} \right) \text{VP} \quad \uparrow = \downarrow$$

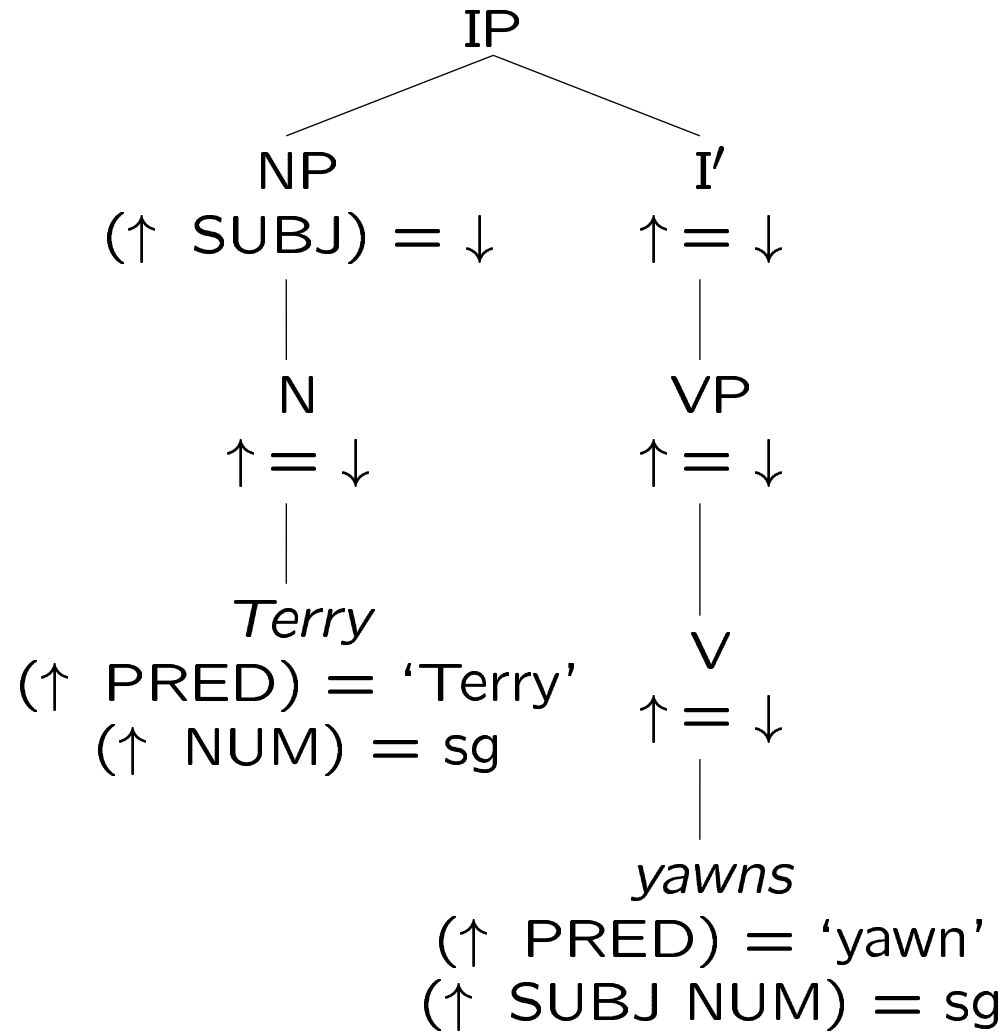
$$\text{VP} \longrightarrow \left(\begin{array}{c} \text{V} \\ \uparrow = \downarrow \end{array} \right) \left(\begin{array}{c} \text{NP} \\ (\uparrow \text{ OBJ}) = \downarrow \end{array} \right)$$

$$\text{NP} \longrightarrow \text{N} \\ \uparrow = \downarrow$$

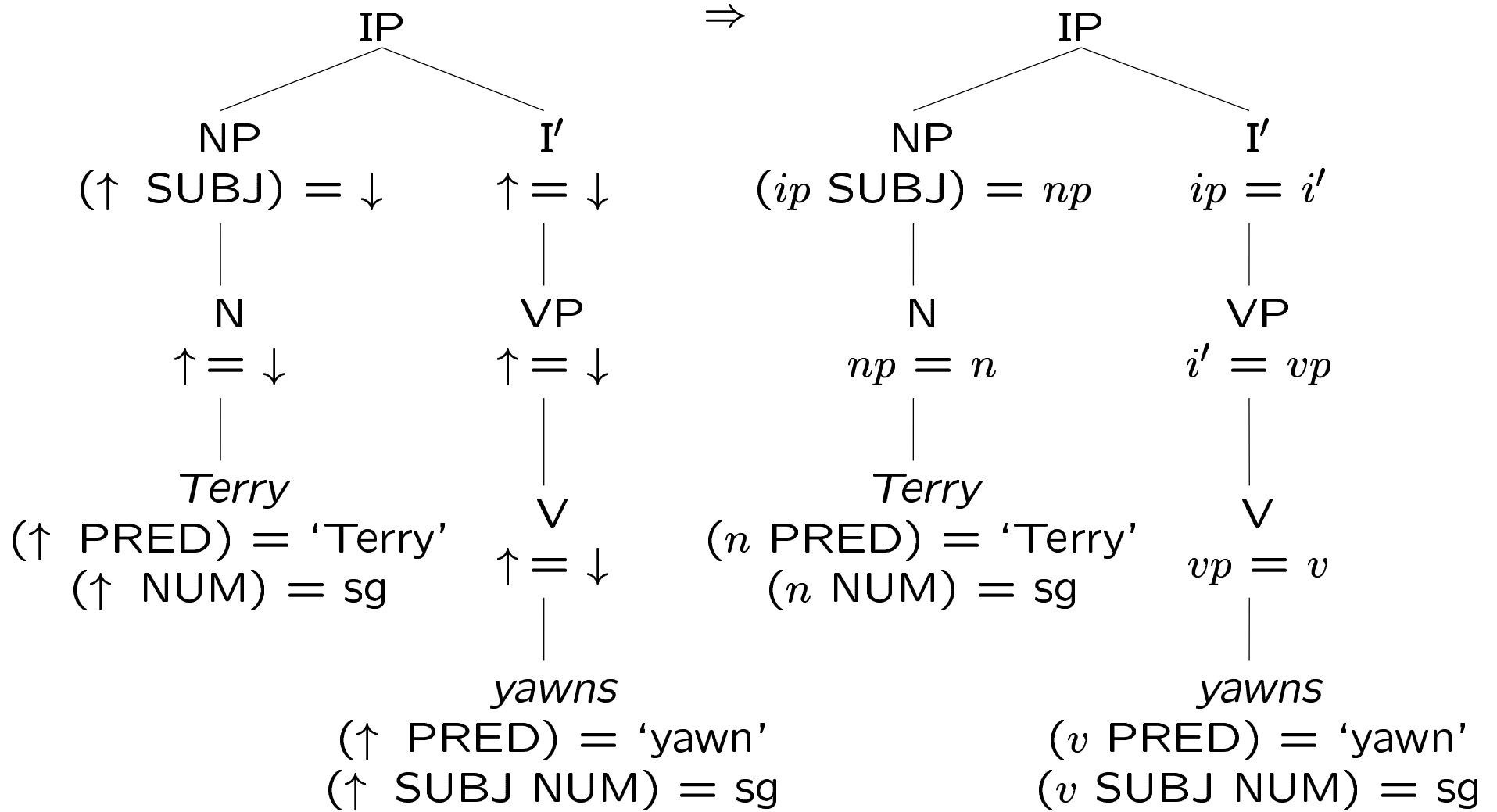
\uparrow is the f-structure of the mother node.

\downarrow is the f-structure of the node the annotation appears on.

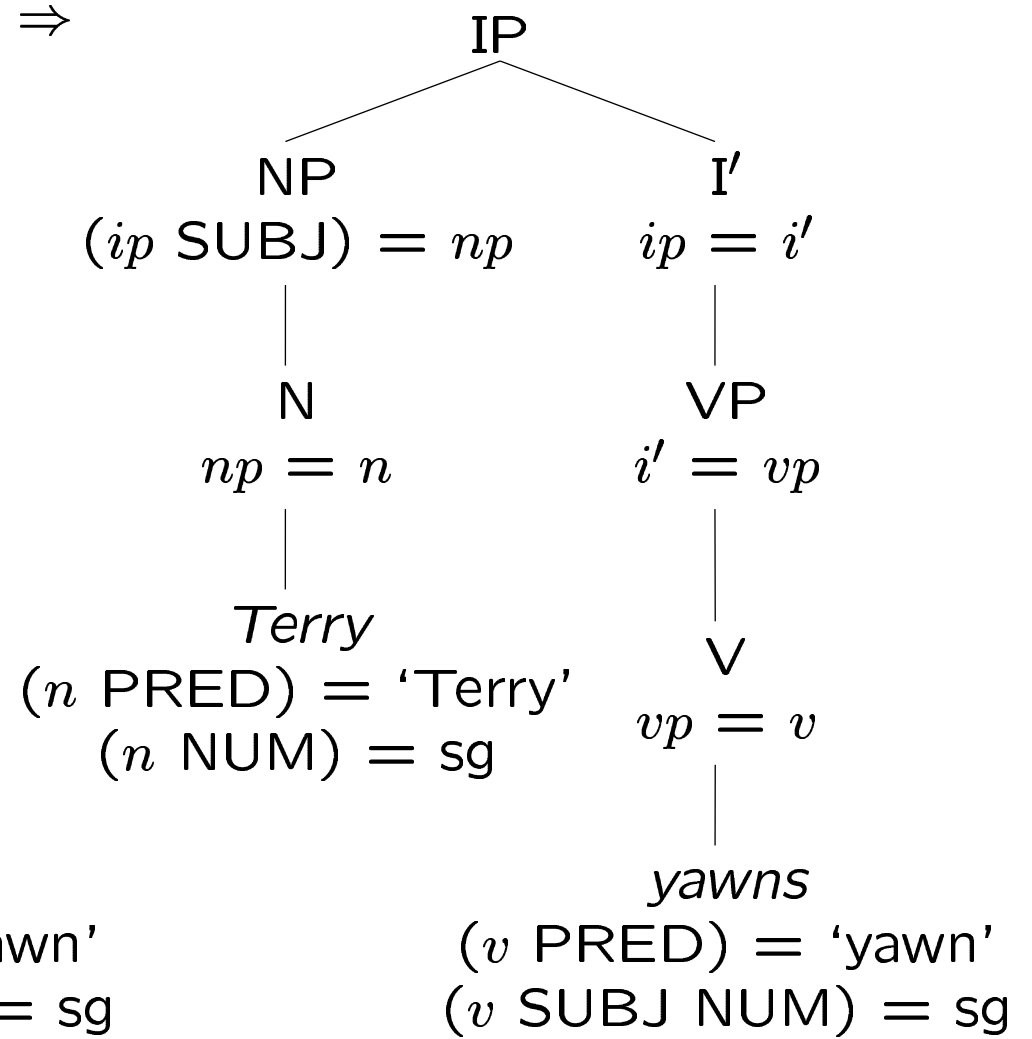
Terry yawns.

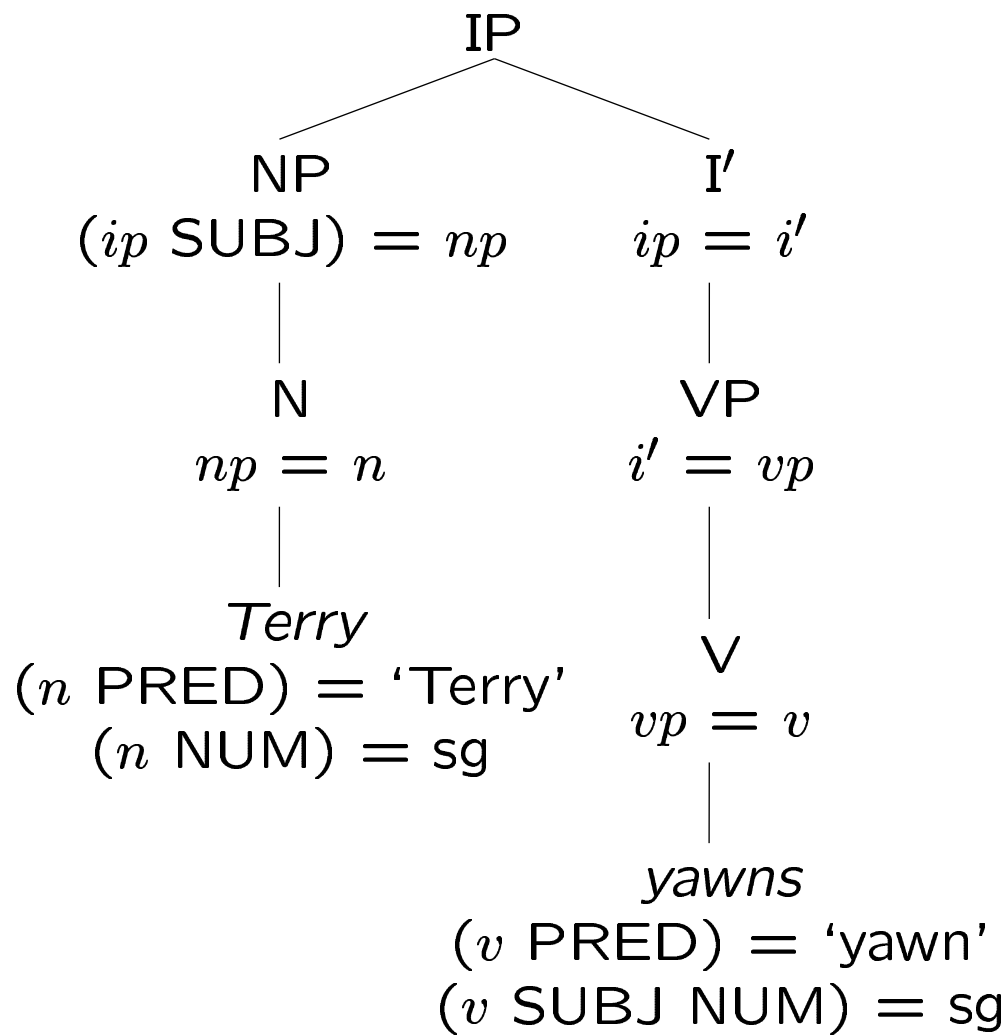


Terry yawns.



\Rightarrow





F-description:

(ip SUBJ) = np
 np = n
 (n PRED) = 'Terry'
 (n NUM) = sg
 ip = i'
 i' = vp
 vp = v
 (v PRED) = 'yawn'
 (v SUBJ NUM) = sg

$(ip \text{ SUBJ}) = np$
 $np = n$
 $(n \text{ PRED}) = \text{'Terry'}$
 $(n \text{ NUM}) = \text{sg}$
 $ip = i'$
 $i' = vp$
 $vp = v$
 $(v \text{ PRED}) = \text{'yawn'}$
 $(v \text{ SUBJ NUM}) = \text{sg}$

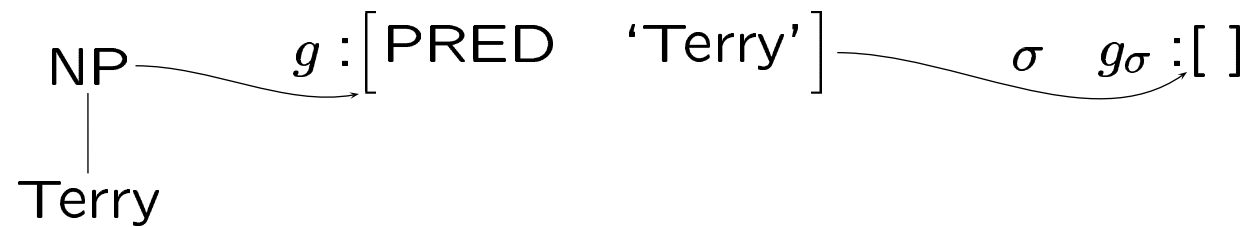
Minimal solution to f-description:

$$ip, i', vp, v : \left[\begin{array}{cc} \text{PRED} & \text{'yawn'} \\ \text{SUBJ} & np, n : \left[\begin{array}{cc} \text{PRED} & \text{'Terry'} \\ \text{NUM} & \text{sg} \end{array} \right] \end{array} \right]$$

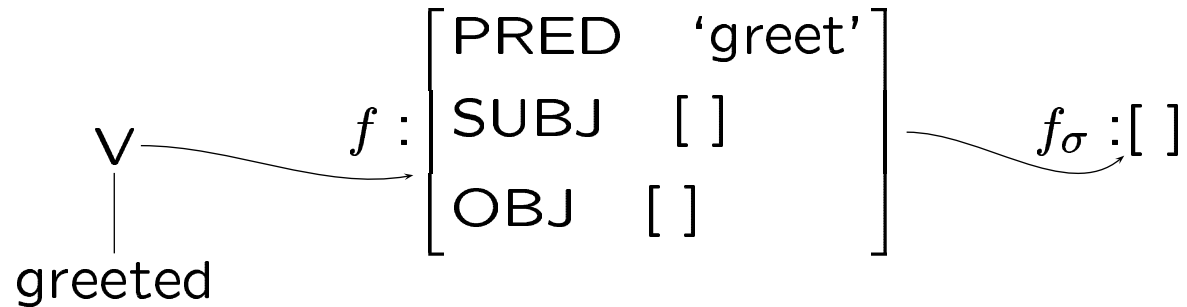
F-structures and semantic projections: Dalrymple et al. (1999)

$(g \text{ PRED}) = \text{'Terry'}$

$Terry : g_\sigma$



“Glue language” deductions:



$$\lambda X.\lambda Y.greet(X, Y) : (f \text{ SUBJ})_\sigma \multimap ((f \text{ OBJ})_\sigma \multimap f_\sigma)$$

- if we have a resource for f 's subject ($f \text{ SUBJ}$)
- then (\multimap) if we have a resource for f 's object ($f \text{ OBJ}$)
- then we have a resource for f .

Augmenting the lexicon with semantic information:

Sam N (\uparrow PRED) = 'Sam'
Sam : \uparrow_σ

Terry N (\uparrow PRED) = 'Terry'
Terry : \uparrow_σ

greeted V (\uparrow PRED) = 'greet'
 $\lambda X.\lambda Y.greet(X, Y) : (\uparrow \text{ SUBJ})_\sigma \multimap ((\uparrow \text{ OBJ})_\sigma \multimap \uparrow_\sigma)$

Sam greeted Terry.

$$f : \left[\begin{array}{l} \text{PRED} \quad \text{'greet'} \\ \text{SUBJ} \quad g : [\text{PRED} \quad \text{'Sam'}] \\ \text{OBJ} \quad h : [\text{PRED} \quad \text{'Terry'}] \end{array} \right]$$

sam : *Sam* : g_σ

terry : *Terry* : h_σ

greeted : $\lambda X.\lambda Y.greet(X, Y) : g_\sigma \multimap (h_\sigma \multimap f_\sigma)$

$$f : \left[\begin{array}{l} \text{PRED} \quad \text{'greet'} \\ \text{SUBJ} \quad g : [\text{PRED} \quad \text{'Sam'}] \\ \text{OBJ} \quad h : [\text{PRED} \quad \text{'Terry'}] \end{array} \right]$$

sam : $Sam : g_\sigma$

terry : $Terry : h_\sigma$

greeted : $\lambda X. \lambda Y. \text{greet}(X, Y) : g_\sigma \multimap (h_\sigma \multimap f_\sigma)$

sam-greeted : $\lambda Y. \text{greet}(Sam, Y) : h_\sigma \multimap f_\sigma$

sam \otimes **terry** \otimes **greeted** (*Premises.*)

\vdash **sam-greeted** \otimes **terry**

\vdash $\text{greet}(Sam, Terry) : f_\sigma$

$$f : \left[\begin{array}{l} \text{PRED} \quad \text{'greet'} \\ \text{SUBJ} \quad g : [\text{PRED} \quad \text{'Sam'}] \\ \text{OBJ} \quad h : [\text{PRED} \quad \text{'Terry'}] \end{array} \right]$$

sam : $Sam : g_\sigma$

terry : $Terry : h_\sigma$

greeted : $\lambda X. \lambda Y. \text{greet}(X, Y) : g_\sigma \multimap (h_\sigma \multimap f_\sigma)$

greeted-terry : $\lambda X. \text{greet}(X, Terry) : g_\sigma \multimap f_\sigma$

sam \otimes **terry** \otimes **greeted** (*Premises.*)

\vdash **sam** \otimes **greeted-terry**

\vdash $\text{greet}(Sam, Terry) : f_\sigma$

Choice of glue language:

Klein and Sag (1985, page 172):

Translation rules in Montague semantics have the property that the translation of each component of a complex expression occurs exactly once in the translation of the whole. . . . That is to say, we do not want the set S [of semantic representations of a phrase] to contain *all* meaningful expressions of IL which can be built up from the elements of S , but only those which use each element exactly once.

Linear logic

multiplicative conjunction: \otimes

linear implication: \multimap

INCORRECT: $A \vdash (A \otimes A)$

INCORRECT: $(A \otimes B) \vdash A$

CORRECT: $(A \otimes (A \multimap B)) \vdash B$

INCORRECT: $(A \otimes (A \multimap B)) \vdash (A \otimes B)$

INCORRECT: $(A \otimes (A \multimap B)) \vdash (A \multimap B) \otimes B$

Completeness and coherence:

*Terry yawns the sink. [**incoherent**]

*Terry devoured. [**incomplete**]

Resources in syntax and semantics: F-descriptions

$(ip \text{ SUBJ}) = np$
 $np = n$
 $(n \text{ PRED}) = \text{'Terry'}$
 $(n \text{ NUM}) = \text{sg}$
 $(n \text{ NUM}) = \text{sg}$
 $(n \text{ NUM}) = \text{sg}$
 $ip = i'$
 $i' = vp$
 $vp = v$
 $(v \text{ PRED}) = \text{'yawn'}$
 $(v \text{ SUBJ NUM}) = \text{sg}$

Minimal solution to f-description:

$$ip, i', vp, v : \left[\begin{array}{cc} \text{PRED} & \text{'yawn'} \\ \text{SUBJ} & np, n : \left[\begin{array}{cc} \text{PRED} & \text{'Terry'} \\ \text{NUM} & \text{sg} \end{array} \right] \end{array} \right]$$

Resources in syntax and semantics:
Glue language premises

$$f : \left[\begin{array}{ll} \text{PRED} & \text{'yawn'} \\ \text{SUBJ} & g : \left[\begin{array}{ll} \text{PRED} & \text{'Terry'} \\ \text{NUM} & \text{sg} \end{array} \right] \end{array} \right]$$

terry : $Terry : g_\sigma$

terry : $Terry : g_\sigma$

terry : $Terry : g_\sigma$

yawns : $\lambda X.yawn(X) : g_\sigma \multimap f_\sigma$

terry \otimes **terry** \otimes **terry** \otimes **yawns**

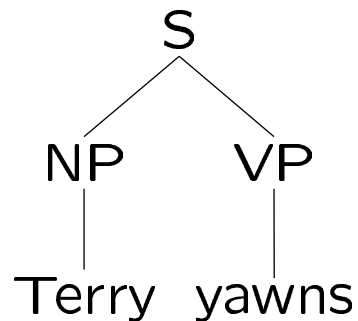
\vdash $yawn(Terry) : f_\sigma \otimes \mathbf{terry} \otimes \mathbf{terry}$ Incoherent!

Resource accounting in syntax: R-LFG (Johnson 1999)

Terry NP $Terry : \text{NOM} \multimap e$
 yawns VP $\lambda x.yawn(x) : \text{SUBJ } e \multimap t$

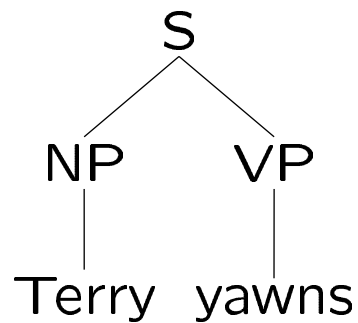
S \longrightarrow NP VP
 SUBJ(NOM, \downarrow) \downarrow

C-structure and f-term for *Terry yawns*:



$$\left[\text{SUBJ} \left[\begin{array}{l} Terry : \text{NOM} \multimap e \\ \text{NOM} \end{array} \right] \right]$$

$$\lambda x.yawn(x) : \text{SUBJ } e \multimap t$$



$$\left[\text{SUBJ} \left[\begin{array}{l} \text{Terry} : \text{NOM} \rightarrow e \\ \text{NOM} \end{array} \right] \right]$$

$$\left[\lambda x. \text{yawn}(x) : \text{SUBJ } e \rightarrow t \right]$$

$$\frac{\frac{\text{Terry} : \text{SUBJ}(\text{NOM} \rightarrow e)}{\text{Terry} : \text{SUBJ NOM} \rightarrow \text{SUBJ } e} \quad \text{SUBJ NOM}}{\text{Terry} : \text{SUBJ } e} \quad \frac{\lambda x. \text{yawn}(x) : \text{SUBJ } e \rightarrow t}{\text{yawn}(\text{Terry}) : t}$$

Johnson (1999): All NPs must receive exactly one case.