# o5: Independence

Jerry Cain April 10<sup>th</sup>, 2024

Lecture Discussion on Ed

# Independence I

#### Independence

Two events *E* and *F* are defined as independent if: P(EF) = P(E)P(F)

Otherwise *E* and *F* are called <u>dependent</u> events.

If *E* and *F* are independent, then:

$$P(E|F) = P(E)$$

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#### Intuition through proof

Independent events *E* and *F* P(EF) = P(E)P(F)

Statement:

If E and F are independent, then P(E|F) = P(E).

Proof:

$$P(E|F) = \frac{P(EF)}{P(F)}$$
$$= \frac{P(E)P(F)}{P(F)}$$
$$= P(E)$$

Definition of conditional probability

Independence of E and F

Taking the bus to cancellation city

Knowing that F happened does not change our belief that E happened.

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#### Dice, our misunderstood friends events *E* and *F* P(E|F) = P(E)Roll two 6-sided dice, yielding values $D_1$ and $D_2$ . Let event *E*: $D_1 = 1$ event F: $D_2 = 6$ event *G*: $D_1 + D_2 = 5$ $G = \{(1,4), (2,3), (3,2), (4,1)\}$ |G| = 4**1.** Are *E* and *F* independent? 2. Are *E* and *G* independent? EF > EAF = { (1, b)} $EG = \{(1,4)\}$ P(E) = 1/6 $P(E) P(F) = \frac{1}{6} \cdot \frac{1}{6} \cdot \frac{1}{36}$ P(E) = 1/6 $P(E_6) = \frac{1}{36}$ P(F) = 1/6 P(EF) = 1/36 $P(EF) = \frac{1}{31}$ P(G) = 4/36 = 1/9 $P(EG) = 1/36 \neq P(E)P(G)$ $\times \underline{dependent} \qquad \boxed{V_b} \qquad \boxed{V_q} \Rightarrow \sqrt{54}$ ✓ independent no!

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Independent  $\square P(EF) = P(E)P(F)$ 

#### Generalizing independence

Three events *E*, *F*, and *G* are independent if:

$$P(EFG) = P(E)P(F)P(G), \text{ and}$$

$$P(EF) = P(E)P(F), \text{ and}$$

$$P(EG) = P(E)P(G), \text{ and}$$

$$P(FG) = P(F)P(G)$$

$$P(FG) = P(F)P(G)$$

$$P(FG) = P(F)P(G)$$

n events 
$$E_1, E_2, \ldots, E_n$$
 are  
independent if:  
$$for r = 1, \ldots, n:$$
for every subset  $E_1, E_2, \ldots, E_r$ :  
$$P(E_1E_2 \ldots E_r) = P(E_1)P(E_2) \cdots P(E_r)$$
informelly; - we need pairwise indeputere for all paire.  
- we need frib - wise indeputere for all prifter.  
- we need gradet - wise indeputere for all grates.  
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#### Dice, increasingly misunderstood (still our friends)

- Each roll of a 6-sided die is an independent trial.
- Two rolls:  $D_1$  and  $D_2$ . •
  - Let event *E*:  $D_1 = 1$ event *F*:  $D_2 = 6$ event *G*:  $D_1 + D_2 = 7$

$$G = \{(1,6), (2,5), (3,4), (4,3), (5,2), (6,1)\}$$

**1.** Are E and F **2.** Are E and G **3.** Are F and G **4.** Are E, F, GEF 10 still { (1,6)}

P(EF) = 1/36



#### Dice, increasingly misunderstood (still our friends)

- Each roll of a 6-sided die is an independent trial.
- Two rolls:  $D_1$  and  $D_2$ .



 $G = \{(1,6), (2,5), (3,4), (4,3), (5,2), (6,1)\}$ 

1. Are *E* and *F* independent?  $\begin{array}{c}
\text{2. Are$ *E*and*G* independent? $<math display="block">
\begin{array}{c}
\text{5. Are$ *E*and*G* independent? $}$   $\begin{array}{c}
\text{6. Are$ *E*and*G* independent? $}$   $\begin{array}{c}
\text{6. Are$ *E*and*G* independent? $}$   $\begin{array}{c}
\text{6. Are$ *E*and*G*  $}$   $\begin{array}{c}
\text{6. Are$ *E*and*G*  $}$   $\begin{array}{c}
\text{6. Are$ *E*and*G*  $}$   $\begin{array}{c}
\text{6. Are$ *E* 

Pairwise independence is not sufficient to prove independence of 3 or more events!

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# Independence II

#### Independent trials

We often are interested in experiments consisting of *n* independent trials.

- *n* trials, each with the same set of possible outcomes
- *n*-way independence: an event in one subset of trials is independent of events in other subsets of trials

Examples:

- Flip a coin *n* times
- Roll a die *n* times
- Send a multiple-choice survey to *n* people
- Send *n* web requests to *k* different servers

#### Network reliability

Consider the following parallel network:

- *n* independent routers, each with probability  $p_i$  of functioning (where  $1 \le i \le n$ )
- E = functional path from A to B exists.

What is P(E)?





#### Network reliability



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### Exercises

#### Independence?

Independent events E and F P(EF) = P(E)P(F)P(E|F) = P(E)

- **1.** True or False? Two events *E* and *F* are independent if:
- A. Knowing that F happens means that E can't happen.
- B. Knowing that F happens doesn't change probability that E happened.
- 2. Are *E* and *F* independent in the following pictures?







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### Coin Flips

Suppose we flip a coin n times. Each coin flip is an **independent trial** with probability p of coming up heads. Write an expression for the following:

- **1.** P(n heads on n coin flips)
- 2. P(n tails on n coin flips)
- 3. P(first k heads, then n k tails)
- **4.** *P*(exactly *k* heads on *n* coin flips)



#### **Coin Flips**

Suppose we flip a coin n times. Each coin flip is an **independent trial** with probability p of coming up heads. Write an expression for the following: S) n consecutive heads -> HHHHH... H => p<sup>n</sup> S) n consecutive tails -> TTT... T => (1-p)<sup>n</sup> = q where q P(n heads on n coin flips)P(n tails on n coin flips)2.  $P(\text{first } k \text{ heads, then } n - k \text{ tails}) \longrightarrow \underbrace{\text{HH} \cdot \text{H}}_{k} \underbrace{\text{TTT} \cdot \text{T}}_{n-k} \Rightarrow P^{k}g^{n-k}$ P(exactly k heads on n coin flips) (any particular segurate of n flips with exactly k heads somewhere within => pk (1-p) and there are  $\binom{n}{k}$  such sequences 's total probability is  $\binom{n}{k}P^{k}(1-p)^{n-k}$  $\binom{n}{k} p^k (1-p)^{n-k}$ # of mutually *P*(a particular outcome's exclusive k heads on n coin flips) outcomes Make sure you understand #4! It will come up again. Stanford University 17

#### Probability of events



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#### Probability of events





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#### De Morgan's Laws

De Morgan's lets you switch between AND and OR.

$$\begin{array}{c} \textbf{S} \\ \textbf{E} \\ \textbf{E} \\ \textbf{F} \\ \textbf$$



#### Hash table fun

- *m* strings are hashed (not uniformly) into a hash table with *n* buckets.  $\sum_{i=1}^{n} p_i = 1$
- Each string hashed is an independent trial w.p.  $p_i$  of getting hashed into bucket i.

What is P(E) if

**1.** E = bucket 1 has  $\geq$  1 string hashed into it?



**2.** E = at least 1 of buckets 1 to k has  $\geq 1$  string hashed into it?



#### Hash table fun

- *m* strings are hashed (not uniformly) into a hash table with *n* buckets.  $\sum_{n=1}^{\infty} \mathbb{P}_{n} = \mathbb{V}$
- Each string hashed is an independent trial w.p.  $p_i$  of getting hashed into bucket i.

What is P(E) if

**1.** E =bucket 1 has  $\ge 1$  string hashed into it?



#### Hash table fun

- *m* strings are hashed (not uniformly) into a hash table with *n* buckets.
- Each string hashed is an independent trial w.p.  $p_i$  of getting hashed into bucket i.

What is P(E) if **1.** E = bucket 1 has  $\geq 1$  string hashed into it? Define  $S_i$  = string *i* is hashed into bucket 1 <u>WTF</u> (not-real acronym for Want To Find):  $S_i^C$  = string *i* is <u>not</u> hashed into bucket 1  $P(E) = P(S_1 \cup S_2 \cup \cdots \cup S_m)$  $= 1 - P((S_1 \cup S_2 \cup \dots \cup S_m)^C)$ Complement  $P(S_i) = p_1$  $= 1 - P(S_1^C S_2^C \cdots S_m^C)$ De Morgan's Law  $P(S_{i}^{C}) = 1 - p_{1}$  $= 1 - P(S_1^{C})P(S_2^{C}) \cdots P(S_m^{C}) = 1 - (P(S_1^{C}))^m$  $S_i$  independent trials  $= 1 - (1 - p_1)^m$ Stanford University 24 Lisa Yan, Chris Piech, Mehran Sahami, and Jerry Cain, CS109, Spring 2024

#### More hash table **fun**: Possible approach?

- m strings are hashed (not uniformly) into a hash table with n buckets.
- Each string hashed is an independent trial w.p.  $p_i$  of getting hashed into bucket *i*.

What is P(E) if

- 1. E = bucket 1 has  $\geq$  1 string hashed into it?
- 2. E = at least 1 of buckets 1 to k has  $\geq 1$  string hashed into it?

$$P(E) = P(F_1 \cup F_2 \cup \cdots \cup F_k)$$
  
=  $1 - P((F_1 \cup F_2 \cup \cdots \cup F_k)^C)$   
=  $1 - P(F_1^C F_2^C \cdots F_k^C)$   
? =  $1 - P(F_1^C) P(F_2^C) \cdots P(F_k^C)$ 

Define  $F_i$  = bucket *i* has at least one string in it

 $F_i$  bucket events are dependent!

So we cannot approach with complement.

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#### More hash table fun

- *m* strings are hashed (not uniformly) into a hash table with *n* buckets.
- Each string hashed is an independent trial w.p.  $p_i$  of getting hashed into bucket i.

What is P(E) if

- 1. E = bucket 1 has  $\geq$  1 string hashed into it?
- 2. E = at least 1 of buckets 1 to k has  $\geq 1$  string hashed into it?

$$P(E) = P(F_1 \cup F_2 \cup \dots \cup F_k)$$
  
=  $1 - P((F_1 \cup F_2 \cup \dots \cup F_k)^C)$   
=  $1 - P(F_1^C F_2^C \cdots F_k^C)$   
=  $P(buckets 1 to k all denied strings)$   
=  $(P(each string hashes to k + 1 or higher))^m$   
=  $(1 - p_1 - p_2 \dots - p_k)^m$ 

$$= 1 - (1 - p_1 - p_2 \dots - p_k)^m$$

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