

$P(n)$ : Every planar graph is 5-colorable.

Base case...

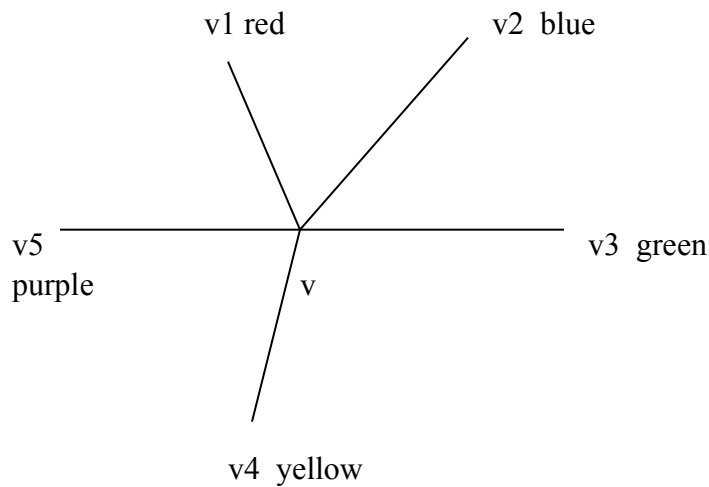
Inductive hypothesis...

Proof:

Let  $G$  be a planar graph with  $n+1$  vertices. We know that  $G$  contains a vertex  $v$  whose degree is at most 5. The graph  $G - v$  is a planar graph with  $n$  vertices and is covered by the inductive hypothesis, so can be colored with 5 colors.

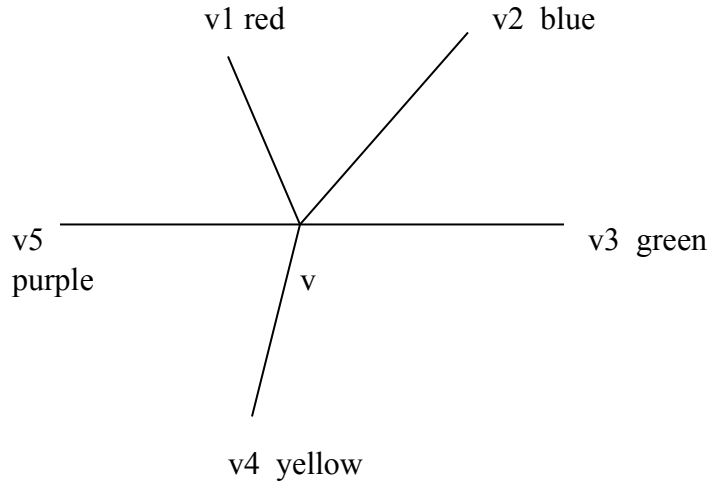
We may assume that  $v$  has exactly 5 neighbors and that they are differently colored, since otherwise there would be at most 4 vertices adjacent to  $v$ . This leaves a spare color for coloring  $v$ .

Here is the situation:



Define  $H(x,y)$  to be the two-colored subgraph of  $G$  induced by all vertices colored  $x$  or  $y$ . Consider  $H(\text{red}, \text{green})$ , the colors of  $v1$  and  $v3$ :

1) If  $v1$  and  $v3$  lie in different components of  $H$ , then we can interchange the colors red and green of all the vertices in  $H$  containing  $v1$ . The result of this recoloring is  $v1$  and  $v3$  will both be green so  $v$  can be colored red.



2) If  $v_1$  and  $v_3$  lie in the same component of  $H$  then there is a circuit of the form  $v \rightarrow v_1 \rightarrow \dots \rightarrow v_3 \rightarrow v$  with alternating colors red and green. Notice that  $v_2$  lies entirely in the circuit and  $v_4$  lies outside, so there cannot be any two-color path between  $v_2$  and  $v_4$  lying entirely in  $H(\text{blue, yellow})$  – the graph is planar and such a path would cross an edge of the circuit. We can therefore interchange the colors of the vertices in the component of  $H(\text{blue, yellow})$  containing  $v_2$ . Then, vertex  $v_2$  and  $v_4$  would both be colored yellow enabling  $v$  to be colored blue.

This completes the proof.