## Graphs

## A Social Network

## Chemical Bonds




PANFLUTE FLOWCHART


A graph is a mathematical structure for representing relationships.

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A graph consists of a set of nodes connected by edges.

## Some graphs are directed.



## Some graphs are undirected.



## How can we represent graphs in $\mathrm{C}++$ ?

## Representing Graphs

We can represent a graph as a map from nodes to the list of nodes each node is connected to.


Node Vector<Node>
Node Adjacent To



## Representing Graphs

- The approach we just saw is called an adjacency list in comes in a number of different forms:

Map<string, Vector<string>> HashMap<string, HashSet<string>> Map<string, Set<string>> Vector<Vector<int>>

- The core idea is that we have some kind of mapping associating each node with its outgoing edges.


## Representing Graphs

The approach we just saw is called an adjacency list in comes in a number of different forms:
Map<string, Vector<string>>

HashMap<string, HashSet<string>>


## Other Graph Representations



This representation is called an adjacency matrix.
For those of you in Math 51: if $A$ is an adjacency matrix for a graph $G$, what is the significance of the matrix $A^{2}$ ?

## Other Representations



Many problems work on an implicit graph.

## You'll find graphs just about everywhere you look.

They're an extremely versatile and powerful abstraction to be aware of.

Going forward, unless stated otherwise, assume we're using an adjacency list.


## Traversing Graphs

## Iterating over a Graph

- In a singly-linked list, there's pretty much one way to iterate over the list: start at the front and go forward!
- In a binary search tree, there are many traversal strategies:
- An inorder traversal that produces all the elements in sorted order.
- A postorder traversal used to delete all the nodes in the BST.
- There are many ways to iterate over a graph, each of which have different properties.

One Search Strategy




> Core idea: Find everything one hop away from the start, then two hops away, then three hops away, etc.

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## Implementing this Idea



Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.

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Queve: (D) (B)

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Load newly-discovered nodes into a queue.


Queue: (B)

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Queue: (B)

Visit nodes in ascending order of
Load newly-discovered distance from the start node $E$. nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


Queve: (B) (A) (G)

Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


Queve: (B) (A) (C)

Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


Queve: (B) (A) (C)

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Queve: (A) (G)

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## Queve: (C) (C)

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## Queue:



Load newly-discovered nodes into a queue.


## Queue: $C$

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Queve: (C) (H)

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Queve: $(H) \oplus$

Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


Queve: $(H) \oplus$

Load newly-discovered nodes into a queue.



Queve: $(H)(F$

Visit nodes in ascending order of distance from the start node $E$.

Load newly-discovered nodes into a queue.


Queue:
(F)

Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


Queve: $F$

Load newly-discovered nodes into a queue.


Visit nodes in ascending order of distance from the start node $E$.


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Queve: $F$

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Queve: $(F)$

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## Queue:



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Queve: (I)

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## Breadth-First Search

- The Queue-based search strategy we just saw is called breadth-first search (or just BFS for short).
- In pseudocode:

```
bfs-from(node v) {
    make a queue of nodes, initially seeded with v.
    while the queue isn't empty:
    dequeue a node curr.
    process the node curr.
    for each node adjacent to curr:
    if that node has never been enqueued:
                enqueue that node.
}
```


## BFS Efficiency

```
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\section*{BFS Efficiency}
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```


Whenever we process a node, we do
1. some fixed amount of work to process the node, then
2. some amount of work proportional to the number of edges touching that node.

\section*{BFS Efficiency}
- Suppose our graph has \(\boldsymbol{n}\) nodes and \(\boldsymbol{m}\) edges.
- These letters are the standard conventions for talking about graphs.
- Average work done per node: \(\mathrm{O}(1)\) baseline work, plus \(\mathrm{O}(m / n)\) work processing edges.
- Number of nodes: \(n\).
- Total work done: \(n \cdot(\mathrm{O}(\mathrm{m} / n)+\mathrm{O}(1))=\mathbf{O}(\boldsymbol{m}+\boldsymbol{n})\).


Whenever we process a node, we do
1. some fixed amount of work to process the node, then
2. some amount of work proportional to the number of edges touching that node.

\section*{BFS Efficiency}
- The amount of work done to run breadthfirst search is \(\mathrm{O}(m+n)\), assuming the graph is represented as an adjacency list.
- Great question to ponder: how fast is breadthfirst search if we use an adjacency matrix?
- The work done is proportional to the number of objects (nodes and edges) that make up the graph.
- We say that BFS is a linear-time graph algorithm.

\section*{A Nifty BFS Trick}


Run BFS, but have each node store a pointer back to the

Queue: node that first discovered it.


Run BFS, but have each node store a pointer back to the node that first discovered it.

Queue:


Run BFS, but have each node store a pointer back to the node that first discovered it.

Queue: \(E\)


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Queve: (A) (C)


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\section*{Queve: (H)}


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\section*{Queve: \(F\)}


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Queve: \(F\)


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Start at any node and follow pointers until you reach \(E\). What path are you
tracing out?

\section*{Shortest Path Routing}
- Breadth-first search can be used to find a shortest path from each node back to the start node.
- The tree you get when you do this is called a breadth-first search tree and has lots of fun properties and cool applications.
- Want to learn more? Take CS161!

\section*{Next Time}
- Minimum Spanning Trees
- Linking Things Cheaply
- Applications of MSTs
- Oh, there are many!```

