# Thinking Recursively Part V 

## Recap from Last Time

## Recursive Backtracking

- In a recursive enumeration problem, we list all solutions to a problem.
- In a recursive optimization problem, we find the best solution to a problem.
- In a recursive backtracking problem, we see whether there even is a solution.


## A Little Word Puzzle

"What nine-letter word can be reduced to a single-letter word one letter at a time by removing letters, leaving it a legal word at each step?"

## One Solution

## S T A R T L I NG

## One Solution

## STARTING

## One Solution

## S TARING

## One Solution

## S T R I NG

## One Solution

## S T I NG

## One Solution

## S I N G

## One Solution

## S I N

## One Solution



## One Solution

I

New Stuff!

## Our Solution, In Action

## The Incredible Shrinking Word



```
bool isShrinkable(const string& word, const Lexicon& english) {
    if (!english.contains(word)) return false;
    if (word.length() == 1) return true;
    for (int i = 0; i < word.length(); i++) {
        string shrunken = word.substr(0, i) + word.substr(i + 1);
        if (isShrinkable(shrunken, english)) {
                        return true;
        }
    }
    return false;
```

\}

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        string shrunken = word.substr(0, i) + word.substr(i + 1);
        return isShrinkable(shrunken, english); // « Bad Idea «
    }
    return false;
```

$\}$

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bool isShrinkable(const string& word, const Lexicon& english) {
    if (!english.contains(word)) return false;
    if (word.length() == 1) return true;
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        string shrunken = word.substr(0, i) + word.substr(i + 1);
        return isShrinkable(shrunken, english); // « Bad Idea «
    }
    return false;
```

\}

## Tenacity is a Virtue



# When backtracking recursively, don't give up if your first try fails! 

Hold out hope that something else will work out. It very well might!

## Recursive Backtracking

if (problem is sufficiently simple) \{
return whether the problem is solvable
\} else \{
for (each choice) \{ try out that choice if (that choice leads to success) \{ return success;
\}
\}
return failure;
\}

Note that if the recursive call succeeds, then we return success. If it doesn't succeed, that doesn't mean we 're failed - it just means we need to try out the next option.

## How do we know we're correct?

## Output Parameters

- An output parameter (or outparam) is a parameter to a function that stores the result of that function.
- Caller passes the parameter by reference, function overwrites the value.
- Often used with recursive backtracking:
- The return value says whether a solution exists.
- If one does, it's loaded into the outparameter.


## Generating the Answer



## Generating the Answer



## Generating the Answer

## CART



## Generating the Answer



## Generating the Answer



## Generating the Answer



## Generating the Answer



## Generating the Answer



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## Generating the Answer



## Generating the Answer



## Dense Crosswords



New York Times Mini Crossword, February 1, 2019


New York Times Mini Crossword, February 1, 2019

Can we design a crossword puzzle where every square must be filled in?






Idea: Fill this in using recursive backtracking.

$8,636^{5}=48,035,594,312,821,554,176$
At one billion grids per second, this will take about three hundred years to complete.

## Speeding Things Up

## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
|  |  |  |  |  |
|  |  |  |  |  |

## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |
|  |  |  |  |  |

## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |

## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | L | I | I |

## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | R | G | H |

## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | H | E | D |
| A | A | R | G | H |$\underbrace{}_{\text {These columns are silly. No }}$| Nords start witht three A's, <br> or three H's, etc. |
| :--- |

## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



The Lexicon has a fast function containsPrefix that's perfect for this.

## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :--- | :--- | :--- | :--- | :--- |
| A | B | A | C | A |
| A | A | H | E | D |
|  |  |  |  |  |
|  |  |  |  |  |

## Generating Dense Crosswords

| A | A | H | E | D |
| :---: | :---: | :---: | :---: | :---: |
| A | B | A | C | A |
| A | A | H | E | D |
|  |  |  |  |  |
|  |  |  |  |  |

## Generating Dense Crosswords



## Generating Dense Crosswords

| A | A | H | E | D |
| :---: | :---: | :---: | :---: | :---: |
| A | B | A | C | A |
| A | A | L | I | I |
|  |  |  |  |  |
|  |  |  |  |  |

## Generating Dense Crosswords

| A | A | H | E | D |
| :---: | :---: | :---: | :---: | :---: |
| A | B | A | C | A |
| A | A | I | I | I |
|  |  |  |  |  |
|  |  |  |  |  |

## Generating Dense Crosswords



## Let's Code it Up!




This word's length is the number of rows.


# bool canMakeCrosswordRec(Grid<char>\& crossword, int nextRow, const Lexicon\& rowWords, const Lexicon\& colWords); 


bool canMakeCrosswordRec(Grid<char>\& crossword, int nextRow, const Lexicon\& rowWords, const Lexicon\& colWords);

## Recursive Backtracking

if (problem is sufficiently simple) \{
return whether the problem is solvable
\} else \{
for (each choice) \{ try out that choice if (that choice leads to success) \{ return success;
\}
\}
return failure;

## Going Deeper

- You can speed this up even more if you're more clever. Here are some thoughts to get you started:
- Once you've placed a few rows down, the columns will be very constrained. Consider switching to going one column at a time versus one row at a time at that point.
- Figure out which row or column is most constrained at each point, and only focus on that row/column.
- Completely optional challenge: Make this program run faster, and find a cool dense crossword. If you find something interesting (and PG-13), we'll share it with the rest of the class!


## Closing Thoughts on Recursion

You now know how to use recursion to view problems from a different perspective that can lead to short and elegant solutions.

You've seen how to use recursion to enumerate all objects of some type, which you can use to find the optimal solution to a problem.

You've seen how to use recursive backtracking to determine whether something is possible and, if so to find some way to do it.

# You've seen that optimizing code is more about changing strategy than writing less code. 

Congratulations on making it this far!

## Your Action Items

- Finish Chapter 9 of the textbook.
- It's all about backtracking, and there are some great examples in there!
- Keep working on Assignment 3.
- You should be done with the Sierpinski Triangle and Human Pyramids, and be making good progress on Shift Scheduling.
- Aim to complete Shift Scheduling and to have started Riding Circuit by Monday.


## Next Time

- Algorithmic Analysis
- How do we formally analyze the complexity of a piece of code?
- Big-O Notation
- Quantifying efficiency!

