# Linked Lists <br> Part Two 

## Recap from Last Time

## Linked Lists at a Glance

- A linked list is a data structure for storing a sequence of elements.
- Each element is stored separately from the rest.
- The elements are then chained together into a sequence.



## Representing a Cell

- For simplicity, let's assume we're building a linked list of strings.
- We can represent a cell in the linked list as a structure:

> struct Cell $\{$ string value; Cell* next;
\};

- The structure is defined recursively!


## A Linked List is Either...

... an empty list, represented by nullptr, or...

a single linked list ... at another linked cell that points...

## Traversing a Linked List

- Once we have a linked list, we can traverse it by following the links one at a time. for (Cell* ptr = list; ptr != nullptr; ptr = ptr->next) \{ /* ... use ptr ... */
\}
list



## Traversing a Linked List

- Once we have a linked list, we can traverse it by following the links one at a time.
for (Cell* ptr = list; ptr != nullptr; ptr = ptr->next) \{
/* ... use ptr ... */
\}



## A Linked List is Either...

... an empty list, represented by nullptr, or...

a single linked list ... at another linked cell that points...

## New Stuff!

## Cleaning Up Our Messes

## Freeing a Linked List

- All good things must come to an end, and we eventually need to reclaim the memory for a linked list.
- The following code triggers undefined behavior. Don't do this!
for (Cell* ptr = list; ptr != nullptr; ptr = ptr->next) \{ delete ptr;
\}


## Freeing a Linked List

- All good things must come to an end, and we eventually need to reclaim the memory for a linked list.
- The following code triggers undefined behavior. Don't do this!
for (Cell* ptr = list; ptr != nullptr; ptr = ptr->next) delete ptr;

$$
\text { ptr } ? ? ?
$$

## Freeing a Linked List Properly

- To properly free a linked list, we have to be able to
- Destroy a cell, and
- Advance to the cell after it.
- How might we accomplish this?


## while (list != nullptr) \{

Cell* next = list->next; delete list; list = next;

## A Linked List is Either...

... an empty list, represented by nullptr, or...

a single linked list ... at another linked cell that points...

## Linked Lists: The Tricky Parts

- Suppose that we want to write a function that will add an element to the front of a linked list.
- What might this function look like?

What went wrong?

```
int main() {
Cell* list = nullptr;
    listInsert(list, "A"):
    listInsert(list, "B");
    listInsert(list, "C");
    return 0;
}
    list
    ~
```



```
int main() {
    Cell* list = nullptr;
    listInsert(list, "A"):
    listInsert(list, "B");
    listInsert(list, "C");
```

    return 0;
    \}
list


## Why does

 nobody love me?
## Pointers by Reference

- To resolve this problem, we can pass the linked list pointer by reference.
- Our new function:
void listInsert(Cell*\& list, const string\& value) \{ Cell* newCell = new Cell; newCell->value = value; newCell->next = list;
list = newCell;
\}


## Pointers by Reference

- To resolve this problem, we can pass the linked list pointer by reference.
- Our new function:
void listInsert(Cell*\& list, const string\& value) \{ Cell* newCell = new Cell; newCell->value = value; newCell->next = list;
list = newCell;
This is a reference to a pointer to a Cell. If we change where list points in this function, the changes will stick:

```
int main() {
Cell* list = nullptr;
    listInsert(list, "A"):
    listInsert(list, "B");
    listInsert(list, "C");
    return 0;
}
    list
    ~
```




## Pointers by Reference

- If you pass a pointer into a function by value, you can change the contents at the object you point at, but not which object you point at.
- If you pass a pointer into a function by reference, you can also change which object is pointed at.


## Time-Out for Announcements!

## Assignment 5

- Assignment 5 (Priority Queue) goes out today. It's due next Friday at the start of class.
- It's a four-parter, and we've included a timetable on the front of the assignment.
- Start this assignment as soon as you get it! You'll have plenty of time to finish everything, but not if you put it off to the last minute.
- Working in pairs is permitted - and encouraged! on this assignment.
- Anton will be holding YEAH hours tomorrow evening. We'll announce the time and location on Piazza and over email.


## Stanford Women in Computer Science

## CASUAL CS DINNER

## [w]

Monday, February 27 from 6-7 PM at the WCC RSVP link here!

Come have dinner with CS students and faculty.
Everyone is welcome, especially students

## Midterm Timetable

- You're done with the midterm exam! Woohoo!
- We'll be grading it over the weekend and returning graded exams on Monday along with stats and solutions.
- Have any questions in the meantime? Just ask!


## Back to Linked Lists!

## Tail Pointers

- A tail pointer is a pointer to the last element of a linked list.
- Tail pointers make it easy and efficient to add new elements to the back of a linked list.
- We can use tail pointers to implement an efficient Queue using a linked list.


## Tail Pointers

- A tail pointer is a pointer to the last element of a linked list.
- Tail pointers make it easy and efficient to add new elements to the back of a linked list.
- We can use tail pointers to implement an efficient Queue using a linked list.



## Enqueuing Things

- Case 1: The queue is empty.

- Case 2: The queue is not empty.


## head

## Dequeuing Things

- Case 1: Dequeuing when there are $2+$ elements.
head

- Case 2: Dequeuing the last element.


## Analyzing Efficiency

- What is the big-O complexity of a dequeue?
- Answer: O(1).
- What is the big-O complexity of an enqueue?
- Answer: O(1).

