# Programming Turing Machines 

## Turing Machines are Hard

|  |  | 1 |  |  |  |  |  | B |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{q}_{5}$ | B | R | $\mathrm{q}_{1}$ | B |  | B | $\mathrm{R} \mathrm{q}_{=}$ | B |  |
| $\mathrm{q}_{1}$ | 1 | R | $\mathrm{q}_{1}$ | 1 | $R q_{x}$ | 1 | $R \mathrm{q}_{=}$ |  | L $\mathrm{q}_{\mathrm{B}}$ |
| $\mathrm{q}_{\mathrm{x}}$ | $\times$ | R | $\mathrm{q}_{1}$ | $\times$ | $R \mathrm{q}_{x}$ | $\times$ | $R \mathrm{q}_{=}$ |  |  |
| $\mathrm{q}_{\text {- }}$ | $=$ | R | $\mathrm{q}_{1}$ | $=$ | $\mathrm{R} \mathrm{q}_{x}$ |  | $R \mathrm{q}_{=}$ |  | L |
| $\mathrm{q}_{\mathrm{R}}$ | 1 | L | $\mathrm{q}_{\mathrm{R}}$ | $\times$ | L $\mathrm{q}_{\mathrm{R}}$ |  | L $\mathrm{q}_{\mathrm{R}}$ | B | R $\mathrm{q}^{\text {d }}$ |
| $\mathrm{q}_{52}$ | 1 |  |  | $\times$ | $\mathrm{R} \mathrm{q}_{\times 2}$ |  | Leject |  | Reject |
| $\mathrm{q}_{\mathrm{x} 2}$ | 1 |  | $\mathrm{q}_{\times 2}$ |  | Reject |  | $\mathrm{R} \mathrm{q}_{=2}$ |  | eject |
| $\mathrm{q}_{\text {- }}$ | 1 |  | $\mathrm{q}_{=2}$ |  | eject |  | Reject |  | L |
| $\mathrm{a}_{12}$ | 1 |  | $\mathrm{q}_{1}$ |  |  |  |  | B |  |

## Outline for Today

- A programming language for Turing machines.
- Design a simple programming language that "compiles" down to Turing machines.
- Keep extending our language to see just how powerful the Turing machine is.


## Our Initial Language: WB

- Programming language WB ("Wang B-machine") controls a tape head over a singly-infinite tape, as in a normal Turing machine.
- Language has six commands:
- Move direction
- Moves the tape head the specified direction (either left or right)
- Write S
- Writes symbol $\boldsymbol{s}$ to the tape.
- Go to $\mathbf{N}$
- Jumps to instruction number $N$ (all instructions are numbered)
- If reading $\boldsymbol{s}$, go to $\mathbf{N}$
- If the current tape symbol is $\boldsymbol{s}$, jump to the instruction numbered $\boldsymbol{N}$.
- Accept and Reject
- Ends the program.
- Statements in WB are executed in the order in which they appear, unless control flow changes.


## A Simple Program in WB

0: If reading $B$, go to 4.
1: If reading 1 , go to 5.
2: Move right.
3: Go to 0.
4: Accept.
5: Reject.

## A Simple Program in WB

## 0000

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5: Reject.

## A WB Program for Even Palindromes

- Suppose we want to test if a string is an even-length palindrome.
- Idea: Cross off the first symbol and match it with the symbol on the far side of the tape.
- If it matches, great! Repeat.
- Otherwise, we should reject.


## A WB Program for Even Palindromes



## A WB Program for Even Palindromes


// Start // M1
0 : If reading 0, go to MO. 10: Write B.
1: If reading 1, go to M1. 11: Move right.
2: Accept
// мо
3: Write B.
4: Move right.
5: If reading 0, go to 4.
6: If reading 1 , go to 4.
7: Move left.
// Next
17: Write B.
8: If reading 0, go to Next.
9: Reject.
12: If reading 0 , go to 11.
13: If reading 1 , go to 11.
14: Move left.
15: If reading 1, go to Next.
16: Reject.
18: Move left.
19: If reading 0 , go to 18
20: If reading 1 , go to 18
21: Move right
22: Go to Start.

## A WB Program for Even Palindromes




## A WB Program for Even Palindromes




## A WB Program for Even Palindromes



## 10010



## A WB Program for Even Palindromes




## A WB Program for Even Palindromes


// Start
0: If reading 0, go to M0.
1: If reading 1, go to M1.
2: Accept
// m0
3: Write B.
4: Move right.
5: If reading 0, go to 4 .
6: If reading 1, go to 4 .
7: Move left.
8: If reading 0 , go to Next.
9: Reject.

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2: Accept
// мо
3: Write B.
4: Move right.
5: If reading 0, go to 4.
6: If reading 1 , go to 4.
7: Move left.
8: If reading 0, go to Next. 9: Reject.
// Next
17: Write B.
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// M1

12: If reading 0 , go to 11.
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6: If reading 1 , go to 4.
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3: Write B.
4: Move right.
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15: If reading 1, go to Next. 16: Reject.
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17: Write B.
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## A WB Program for Even Palindromes



// M1

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// Next
17: Write B.
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19: If reading 0, go to 18
20: If reading 1 , go to 18
21: Move right
22: Go to Start.

## A WB Program for Even Palindromes




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## A WB Program for Even Palindromes




## WB and Turing Machines

- Recall: A language $L$ is recursively enumerable iff there is a TM for it.
- Theorem: A language $L$ is recursively enumerable iff there is a WB program for it.
- Need to show the following:
- Any TM can be converted into an equivalent WB program.
- Any WB program can be converted into an equivalent TM.


## From Turing Machines to WB

- Basic idea: Construct a small WB program for each state that simulates that state.
- Combine all programs together to get an overall WB program that simulates the Turing machine.


## A State in a Turing Machine

- There are three kinds of states in a Turing machine:
- Accepting states,
- Rejecting states, and
- "Working" states.
- We can easily build WB programs for the first two:
$/ / q_{\mathrm{acc}}$
$0:$ Accept
// $\mathrm{q}_{\mathrm{rej}}$
0: Reject


## Working States

- At a given working state in a Turing machine, we will do exactly the following, in this order:
- Read the current symbol.
- Write back a new symbol based on this choice of symbol.
- Transition to some destination state.
- Could we build a WB program for this?


## Working States



```
// \(q_{0}\)
    \(0:\) If reading 0 , go to \(0 q_{0}\).
    1: If reading 1 , go to \(1 q_{0}\).
    2: If reading B , go to \(\mathrm{Bq}_{0}\).
// \(0 \mathrm{q}_{0}\)
    3: Write B.
    4: Move right.
    5: Go to \(\mathrm{q}_{1}\)
// \(1 q_{0}\)
    6: Write 0
    7: Move left.
    8: Go to \(\mathrm{q}_{0}\)
```

// $0 \mathrm{q}_{0}$
3: Write B.
4: Move right.
5: Go to $\mathrm{q}_{1}$
// $1 q_{0}$
6: Write 0
7: Move left.
8: Go to $\mathrm{q}_{0}$
// Bq ${ }_{0}$
9: Write B
10: Move right.
11: Go to $\mathrm{q}_{\text {acc }}$

## A Complete Construction


// $\mathrm{q}_{0}$
0 : If reading 0 , go to 3 .
1: If reading 1 , go to 6.
2: If reading $B$, go to 9.
3: Write 0.
4: Move right.
5: Go to $\mathrm{q}_{1}$.
6: Write 1.
7: Move right.
8: Go to $\mathrm{q}_{\text {re }}$.
9: Write 1.
10: Move right.
11: Go to $\mathrm{q}_{\mathrm{acc}}$.
// $\mathrm{q}_{\mathrm{acc}}$
12: Accept.

13: If reading 0 , go to 16.
14: If reading 1 , go to 19.
15: If reading $B$, go to 22.
16: Write 0.
17: Move right.
18: Go to $\mathrm{q}_{\mathrm{rej}}$.
19: Write 1.
20: Move right.
21: Go to $\mathrm{q}_{0}$.
22: Write 1.
23: Move right.
24: Go to $q_{\text {acc }}$.
// $\mathrm{q}_{\mathrm{rej}}$
25: Reject.

## From WB to Turing Machines

- We now need a way to convert a WB program into a Turing machine.
- Construction sketch:
- Create a state in the TM for each line of the WB program.
- Introduce extra "helper" states to implement some of the trickier instructions.
- Connect the states by transitions that simulate the WB program.
- We will show how to translate each WB command into a collection of states plus transitions.


## Refresher: Turing Machine Notation

$$
B \rightarrow B, R
$$

start

## Refresher: Turing Machine Notation

- The accept and reject states are denoted

- A transition of the form

means "on seeing $x$, write $y$ and move direction $D . "$


## Accept and Reject

- The Accept and Reject commands are the easiest to translate.
- To translate N : Accept into TM states, construct the following:

- To translate N : Reject into TM states, construct the following:



## Move left and Move right

- We can translate N : Move left and N : Move right by having the TM do the following:
- Write back the same symbol that was already on the tape (ensuring that we don't change the tape).
- Move in the indicated direction.
- Transition into the state representing line $\mathrm{N}+1$.



## Go to $L$

- The line N : Go to m needs to change into the state for line m without moving the tape head.
- All TM transitions move the tape head; how might we address this?
- Move right and change into a new state that then moves back to the left.



## Write S

- The line $\mathbf{N}$ : Write $\boldsymbol{s}$ needs to
- Write the symbol $\boldsymbol{s}$,
- Leave the tape head where it is, and
- Move to line $\mathbf{N}+1$.
- We use a similar trick as before:



## If reading $s$, go to $M$

- The line $\mathrm{N}: ~ I f ~ r e a d i n g ~ s, ~ g o ~ t o ~ m e i t h e r ~$
- Executes a "go to m" step as before if reading $\mathbf{s}$, or
- Does nothing and transitions to state $\mathbf{N}+1$.



## A Complete Conversion



## The Story So Far

- We have just built a simple programming language that is equivalent in power to a Turing machine.
- This language, however, makes for some very complicated programs.
- Let's add some new features to our programming language to make it a bit easier to work with.


## Revisiting Even Palindromes

// m0
3: Write B.
4: Move right.
5: If reading 0 , go to 4.
6: If reading 1 , go to 4.
7: Move left.
8: If reading 0 , go to Next.
9: Reject.
// Next
17: Write B.
18: Move left.
19: If reading 0 , go to 18
20: If reading 1 , go to 18
21: Move right.
22: Go to Start.

- Steps 4-6 essentially say "move right, then move right until you read a blank."
- Steps 18 - 20 essentially say "move left, then move left until you read a blank."
- Is it really necessary to write this out each time?


## Introducing WB2

- The programming language WB2 is the language WB with two new commands:
- Move left until $\left\{\boldsymbol{S}_{1}, \boldsymbol{S}_{2}, \ldots, \boldsymbol{S}_{n}\right\}$.
- Moves the tape head left until we read one of $\boldsymbol{s}_{1}, \boldsymbol{s}_{2}, \boldsymbol{s}_{3}, \ldots, \boldsymbol{s}_{\boldsymbol{n}}$.
- Move right until $\left\{\boldsymbol{S}_{1}, \boldsymbol{S}_{2}, \ldots, \boldsymbol{S}_{\boldsymbol{n}}\right\}$.
- Moves the tape head right until we read one of $\boldsymbol{s}_{1}, \boldsymbol{s}_{2}, \boldsymbol{s}_{3}, \ldots, \boldsymbol{s}_{\boldsymbol{n}}$.
- Both commands are no-ops if we're already reading one of the specified symbols.
- We can write programs in WB2 that are much easier to read than in WB.


## A WB Program for Even Palindromes



## A WB2 Program for Even Palindromes

| // Start | // M1 |
| :---: | :---: |
| 0: If reading 0, go to M0. | 9: Write B. |
| 1: If reading 1, go to M1. | 10: Move right. |
| 2: Accept | 11: Move right until \{B\} |
|  | 12: Move left. |
| // M0 | 13: If reading 1, go to Next. |
| 3: Write B. | 14: Reject. |
| 4: Move right. |  |
| 5: Move right until \{B\}. | // Next |
| 6: Move left. | 15: Write B. |
| 7: If reading 0, go to Next. | 16: Move left. |
| 8: Reject. | 17: Move left until \{B\}. |
|  | 18: Move right. |

## A WB2 Program for BALANCE

- Let $\Sigma=\{0,1\}$ and consider the language BALANCE:

$\left\{w \in \Sigma^{*} \mid w\right.$ has the same number of 0s and 1s. \}

- Let's write a WB2 program for BALANCE.


## A WB2 Program for BALANCE

// Start 0 : Move right until $\{0,1, B\}$.

1: If reading 0 , go to Match0.
2: If reading 1 , go to Match1.
3: Accept.
// Match1
9: Write B.
10: Move right.
11: Move right until $\{0, B\}$.
12: If reading 0 , go to Found.
13: Reject.
// Match0
4: Write B.
5: Move right.
6: Move right until $\{1, \mathrm{~B}\}$.
7: If reading 1, go to Found.
8: Reject.
// Found
14: Write x.
15: Move left until \{B\}.
16: Move right.
17: Go to Start.

## WB2 and Turing Machines

- Theorem: A language is recursively enumerable iff there is a WB2 program for it.
- We could directly prove this again by showing equivalence with Turing machines.
- Instead, we'll connect it to WB:



## From WB2 to WB

- We will show how to turn any WB2 program into an equivalent WB program.
- All old instructions are still valid.
- We need to show how to implement the new Move ... until commands using just WB.


## Implementing Move ... until

- Replace N: Move dir until $\left\{\boldsymbol{s}_{1}, \ldots, \boldsymbol{s}_{n}\right\}$ as follows:

| $\mathrm{N}+0:$ | If reading $s_{1}$, go to $\mathrm{N}+\mathrm{n}+2$. |
| :--- | :--- |
| $\mathrm{N}+1:$ | If reading $s_{2}$, |
| $\mathrm{N}+2:$ | If reading $s_{3}$, go to $\mathrm{N}+\mathrm{n}+2$. |

$\mathrm{N}+(\mathrm{n}-1)$ : If reading $\mathrm{s}_{\mathrm{n}}$, go to $\mathrm{N}+\mathrm{n}+2$.
$\mathrm{N}+\mathrm{n}$ : Move dir.
$\mathrm{N}+\mathrm{n}+1$ : Go to N

- Renumber other lines as appropriate.


## Why This Matters

- We are starting to move more and more away from the Turing machine with from we started.
- The structure of our approach is
- Find some simple programming language that can be directly translated into a Turing machine (and viceversa).
- Add new features to the language, and show how to implement those new features using the old language.
- Add new features to that language, and show how to implement those features using the previous language.
- (etc.)
- Conclude that the final language is equivalent to a Turing machine.


## A Repeating Pattern

// Match0
4: Write B.
5: Move right.
6: Move right until $\{1, B\}$.
7: If reading 1, go to Found.
8: Reject.
// Match1
9: Write B.
10: Move right.
11: Move right until $\{0, B\}$.
12: If reading 0 , go to Found.
13: Reject.

## A Simple Memory

- Right now, our programming language WB2 has no variables in it.
- To solve larger classes of problems, let's invent a new language WB3 that has support for variables.
- We will severely limit the scope of our variables:
- Only finitely many total variables throughout the program.
- Each variable can only hold a single tape symbol.
- Each variable initially holds the blank symbol.


## Our New Commands

- We will define WB3 as WB2 with the following extra commands:
- Load $s$ into $v$.
- Sets the variable $\boldsymbol{v}$ equal to tape symbol $\boldsymbol{s}$.
- Load current into $v$.
- Sets the variable $\boldsymbol{v}$ equal to the currently-scanned tape symbol.
- If $v_{1}=v_{2}$, go to $L$.
- If $v_{1}$ and $v_{2}$ have the same value, go to instruction $\boldsymbol{L}$.
- These may be constants or variables.
- Additionally, any command that referenced a tape symbol (for example, write, if reading, move ... until) can refer to variables in addition to constants.


## A WB2 Program for Even Palindromes

| // Start | // M1 |
| :---: | :---: |
| 0: If reading 0, go to M0. | 9: Write B. |
| 1: If reading 1, go to M1. | 10: Move right. |
| 2: Accept | 11: Move right until \{B\} |
|  | 12: Move left. |
| // M0 | 13: If reading 1, go to Next. |
| 3: Write B. | 14: Reject. |
| 4: Move right. |  |
| 5: Move right until \{B\}. | // Next |
| 6: Move left. | 15: Write B. |
| 7: If reading 0, go to Next. | 16: Move left. |
| 8: Reject. | 17: Move left until \{B\}. |
|  | 18: Move right. |

## A WB3 Program for Even Palindromes

// Start
0 : Read current into X .
1: If $\mathrm{X}=\mathrm{B}$, go to Acc.
2: Write B.
3: Move right.
4: Move right until $\{B\}$.
5: Move left.
6: If reading X , go to Match.
7: Reject.
// Match 8: Write B.

9: Move left.
10: Move left until B.
11: Move right.
12: Go to Start.
// Acc:
13: Accept.

## A WB2 Program for BALANCE

// Start 0 : Move right until $\{0,1, B\}$.

1: If reading 0 , go to Match0.
2: If reading 1 , go to Match1.
3: Accept.
// Match1
9: Write B.
10: Move right.
11: Move right until $\{0, B\}$.
12: If reading 0 , go to Found.
13: Reject.
// Match0
4: Write B.
5: Move right.
6: Move right until $\{1, \mathrm{~B}\}$.
7: If reading 1, go to Found.
8: Reject.
// Found
14: Write x.
15: Move left until \{B\}.
16: Move right.
17: Go to Start.

## A WB3 Program for BALANCE

// Start
0: Move right until $\{0,1, B\}$. 8: Write B.
1: If reading $B$, go to Acc.
2: If reading 0 , go to 5.
3: Load O into Y.
4: Go to Scan.
5: Load 1 into Y.
6: Go to Scan.
// Scan

9: Move right.
10: Move right until \{Y, B\}
11: If reading $Y$, go to 13.
12: Reject.
13: Write x.
14: Move left until B.
15: Move right.
16: Go to Start.
// Acc:
17: Accept.

## Equivalence of WB2 and WB3

- Theorem: A language is recursively enumerable iff there is a WB3 program for it.
- Adding in these sorts of variables adds no power to our model of computation!
- To prove the theorem, we will show
- Any WB2 program can be converted to a WB3 program, and
- Any WB3 program can be converted to a WB2 program.



## The Proof: An Intuition

- Our programs allow only finitely many variables holding only one of finitely many different values (tape symbols).
- We could just replicate the program for each possible assignment to the variables, then hardcode in the behavior in each of these cases.
- Could make the program staggeringly huge, but it will still be finite!


## The Transformation, Part I

- Let $\mathrm{V}_{1}, \mathrm{~V}_{2}, \ldots, \mathrm{~V}_{\mathrm{n}}$ be the variables referenced in the program.
- We can just look at the source code to determine this.
- Make $\mid \Gamma{ }^{\mid n}$ copies of the initial program, one for each possible assignment of tape symbols to the variables $\mathrm{V}_{\mathrm{i}}$.
- Order the copies arbitrarily, but make the version where all variables hold B come first.


## The Transformation, Part II

- We now have a whole bunch of copies of WB3 programs.
- We need to convert them into legal WB2 programs.
- This works in two steps:
- Removing variables from older WB2 commands like Write, If reading ..., and Move ... while.
- For example: "Write $\boldsymbol{x}$," where $\boldsymbol{x}$ is a variable.
- Rewriting all new WB3 commands that reference variables to use only WB2 commands.
- For example: "Load current into x."


## Eliminating Variables from WB2

- Removing variables from purely WB2 statements is easy because we've copied the program so many times.
- For each copy, replace all variables in WB2 statements with the value that the variable has in that copy.

$$
\begin{aligned}
& 0: \text { Load } 0 \text { into } Y . \\
& 1: \text { Write } Y . \\
& 2: \text { Accept }
\end{aligned}
$$

```
0: Load O into Y.
1: Write Y.
2: Accept
```

$$
Y=B
$$


$\mathrm{Y}=0$


## Eliminating Variables from WB2

- Removing variables from purely WB2 statements is easy because we've copied the program so many times.
- For each copy, replace all variables in WB2 statements with the value that the variable has in that copy.

$$
\begin{aligned}
& 0: \text { Load } 0 \text { into } Y . \\
& 1: \text { Write } Y . \\
& 2: \text { Accept }
\end{aligned}
$$

```
0 : Load 0 into \(Y\).
1: Write B.
2: Accept
```

$$
Y=B
$$


$\mathrm{Y}=0$


## Eliminating Variables from WB3

- We can eliminate commands that manipulate variables by replacing them with Go mos.
- There are three commands to eliminate:
- Load s into V.
- Load current into V.
- If $\boldsymbol{v}_{1}=\boldsymbol{v}_{2}$, go to $L$.


## If $v_{1}=v_{2}$, go to $L$

- We can eliminate this statement by just hardcoding the jump in place.
- If in the current copy of the program $\boldsymbol{v}_{1}$ and $\boldsymbol{v}_{2}$ have the same values, replace with

Go to $L$
where $L$ is the corresponding version of $L$ in this copy.

- Otherwise, replace with

Go to $N$
where $\boldsymbol{N}$ is the number of the next line in the program.

## Load $s$ into $v$

- To simulate the effect of loading $\boldsymbol{s}$ into $\boldsymbol{v}$, we can jump out of the current copy of the program into the copy where $v$ has value $s$.
0: Load 0 into $Y$.
1: Write $Y$.
2: Accept
$0:$ Load 0 into $Y$.
1: Write Y.
2: Accept

$\mathrm{Y}=0$

$$
\mathrm{Y}=0
$$

6: Load 0 into $Y$.
7: Write Y.
8: Accept

$$
Y=B
$$

$$
\mathrm{Y}=1
$$

## Load $s$ into $v$

- To simulate the effect of loading $\boldsymbol{s}$ into $\boldsymbol{v}$, we can jump out of the current copy of the program into the copy where $v$ has value $s$.
0: Load 0 into $Y$.
1: Write $Y$.
2: Accept
$0:$ Load 0 into $Y$.
1: Write B.
2: Accept

$\mathrm{Y}=0$

$$
\mathrm{Y}=0
$$

6: Load 0 into $Y$.
7: Write 1.
8: Accept

$$
Y=B
$$

$$
Y=1
$$

## Load $s$ into $v$

- To simulate the effect of loading $\boldsymbol{s}$ into $\boldsymbol{v}$, we can jump out of the current copy of the program into the copy where $v$ has value $s$.

```
O: Load O into Y.
1: Write Y.
2: Accept
```

```
0: Go to 4.
1: Write B.
2: Accept
```

3: Go to 4.
4: Write 0.
5 : Accept

$$
\begin{aligned}
& 6: \text { Go to } 4 . \\
& 7: \text { Write } 1 . \\
& 8: \text { Accept }
\end{aligned}
$$

$\mathrm{Y}=0$

$$
Y=1
$$

## Load current into $V$

- We can simulate this instruction using a similar trick to before.
- Replace this instruction as follows:

```
If reading s}\mp@subsup{s}{1}{}\mathrm{ , go to LoadS 
If reading }\mp@subsup{s}{2}{},\mathrm{ go to LoadS 
If reading }\mp@subsup{s}{n}{}\mathrm{ , go to LoadS 
// LoadS 
Load s}\mp@subsup{s}{1}{}\mathrm{ into v.
Go to Done.
// LoadS
Load }\mp@subsup{s}{n}{}\mathrm{ into v.
Go to Done.
// Done:
```


## Souping up our Tape

- Up to this point, we've been improving our WB programming language by adding in new ways of scanning over the tape.
- What if we made changes to the tape itself?


## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5 .
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5 .
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5 .
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into x .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
$X$
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into x .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



$\mathbf{O}$
$\mathbf{X}$
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



$\mathbf{0}$
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



0
$\mathbf{X}$
// Start
0 : Read track 1 into x .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



$\mathbf{0}$
$\mathbf{x}$
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5 .
4: Go to 0
5: /* ... */

## A Multitrack Tape



$\mathbf{0}$
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



$\mathbf{0}$
X
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape



1
$X$
// Start
0: Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0: Read track 1 into X.
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5 .
4: Go to 0
5: /* ... */

## A Multitrack Tape


// Start
0 : Read track 1 into X .
1: Move right.
2: Write X into track 2
3: If reading $B$ on track 1 , go to 5.
4: Go to 0
5: /* ... */

## Introducing WB4

- Let's define WB4 to be WB3 with the introduction of finitely many tracks on the tape.
- The tape head still moves as a unit to the left or right, but we can now issue read and write commands to any cell in the current track.
- All previous commands updated to specify which track is to be read or written.


## A Surprising Theorem

- Theorem: A language is recursively enumerable iff there is a WB4 program for it.
- This is not obvious... it seems like adding in more tracks should increase the power of our programming language!
- As with before, will prove that all WB4 programs are equivalent to WB3 programs.



## The Intuition

- Treat a single tape as a "fat tape" where each tape symbol encodes the contents of the cells of all four tracks.
- Each read or write to a specific location replaces the entire tape cell with a new symbol representing the change.



## A Sketch of the Construction

- Replace each instruction that reads or writes a track with a huge cascading "if" that checks for every possible tape symbol and reacts accordingly.
- Can make the program enormously bigger, but it still ends up finite.
- I'm not even going to attempt to fit something like that onto these slides.


## Where We Are Now

- Starting with WB, we have added
- Loops to search for a value. (WB2)
- Variables with finite storage. (WB3)
- Multiple tracks. (WB4)
- Yet we still accept exactly the same set of languages.
- Every WBn program can be converted back to a TM.



## Making Things Crazier

- What do you get when you combine a PDA and a WB4 program?
- A program with an infinite tape, plus multiple stacks!



## Introducing WB5

- The programming language WB5 is the programming language WB4 with the addition of a finite number of stacks.
- We add three extra commands:
- Push s onto stack $\boldsymbol{V}$.
- Pushes the symbol $\boldsymbol{s}$ onto the stack named $\boldsymbol{v}$.
- If stack $V$ is empty, go to $L$.
- If stack $v$ is empty, go to instruction $L$.
- Pop stack $v$ into w.
- If stack $\boldsymbol{v}$ is nonempty, pops $\boldsymbol{v}$ and puts the top into $\boldsymbol{w}$.


## The Multiplication Language

- Let $\Sigma=\{0,1,2\}$ and consider the language 01MULT defined as
$\left\{w \in \Sigma^{*} \mid\right.$ the number of $2^{\prime} \mathrm{s}$ in $w$ is the product of the number of 1 's and the number of 0's. \}
- For example:
- 00112222 E 01MULT
- 22001122122 G 01MULT
- This language is neither context-free nor regular.
- How could we write a WB5 program for it?


## One Approach



## One Approach



## One Approach



## One Approach



## One Approach



## One Approach



## One Approach



## One Approach

## 0



## One Approach



## One Approach



## One Approach



## One Approach



## One Approach

## One Approach



## One Approach

## 0



## One Approach



## One Approach



## One Approach



## One Approach

## One Approach



## One Approach



## One Approach



## One Approach



## WB5 Program for 01MULTI

| // Start | // Load1 |
| :--- | :--- |
| 0: If reading 0, go to Load0. | $7:$ Push 1 onto Stack 1. |
| 1: If reading 1, go to Load1. | 8: Move right. |
| 2: If reading 2, go to Load2. | 9: Go to Start. |
| 3: Go to Check. |  |
|  | // Load2 |
| // Load0 | $10:$ Push 2 onto Stack 2. |
| 4: Push 0 onto Stack 0. | $11:$ Move right. |
| 5: Move right. | $12:$ Go to Start. |
| 6: Go to Start. |  |

## WB5 Program for 01MULTI

## // Check:

13: If Stack 0 is empty, go to Ver.
14: Pop Stack 0.
15: If Stack 1 is empty, go to Fix. 24: Push 1 onto Stack 1.
16: Pop Stack 1.
17: Push 1 onto Stack 1T.
18: If Stack 2 is empty, go to Rej. // Ver:
19: Pop Stack 2.
20: Go to 15.
// Rej:
21: Reject.
// Fix:
22: If $S t 1 T$ is empty, go to Check.
23: Pop Stack 1T.

25: Go to Fix.

26: If Stack 2 is empty, go to Acc. 27: Reject.
// Acc:
28: Accept.

## A Pretty Ridiculous Theorem

- Theorem: A language is recursively enumerable iff there is a WB5 program for it.
- So adding in finitely many infinite stacks doesn't give us any more expressive power!
- As with before, will prove that all WB5 programs are equivalent to WB4 programs.



## From Stacks to Tracks

- The key idea behind the construction for converting WB5 programs into WB4 programs is to represent each stack with its own track.
- If there are $n$ stacks in the program, we will add $n+1$ tracks:
- One track for each of the $n$ stacks, and
- One track for bookkeeping.
- If the WB5 program was using any tracks, we'll keep them as well and add these new ones in separately.



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |  |  |  |



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\gg$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |


| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\boldsymbol{>}$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |



| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\gg$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\boldsymbol{>}$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |



0 : Write $\times$ on track 5.


| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\boldsymbol{>}$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\boldsymbol{>}$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |



0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\gg$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4
6: Move left until $\{>\}$ on track 4.

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\gg$ | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4
6: Move left until $\{>\}$ on track 4.

| 1 | 1 | 0 | 0 |  | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  | ... |

0: Push 1 onto Stack 3.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4
6: Move left until $\{>\}$ on track 4.
7: Move right until $\{x\}$ on track 5.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |  | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  | Pu | ush | 1 | on | to | S | ac | 3 |  |  |  |
| 11 |  |  |  |  |  |  | 0 |  |  |  | le | on | unt | il | $\begin{aligned} & k \\ & \{> \end{aligned}$ |  | n | ra |  | 4. |
| 01 | 1 |  |  |  |  |  |  |  | i | ve <br> ite | $\begin{array}{r} \text { ri } \\ 1 \end{array}$ | ght on | $\begin{aligned} & \mathrm{ur} \\ & \mathrm{tr} \end{aligned}$ | $\begin{aligned} & \text { ti } \\ & \text { cac } \end{aligned}$ |  | $<\}$ | on | tr |  | $\text { k } 4$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 11 |  |  |  |  |  |  |  |  |  | ite | < | on | $t$ | ac | k |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  | ve | 1 | ft | un | il | \{> | \} | n | ra |  |  |
|  |  |  |  |  |  |  |  |  |  | ve | ri | ght | un | ti |  |  |  |  |  |  |


| 1 | 1 | 0 | 0 | 1 |  | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  |  |


| 1 | 1 |  |
| :--- | :--- | :--- |
| 0 | 1 | 1 |
| 1 | 1 |  |

0 : Write $x$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4
6: Move left until $\{>\}$ on track 4.
7: Move right until $\{x\}$ on track 5.
8: Write B on track 5.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |


| 1 | 1 |  |
| :--- | :--- | :--- |
| 0 | 1 | 1 |
| 1 | 1 |  |

0 : Write $x$ on track 5.
1: Move left until $\{>\}$ on track 4.
2: Move right until $\{<\}$ on track 4.
3: Write 1 on track 4.
4: Move right.
5: Write < on track 4
6: Move left until $\{>\}$ on track 4.
7: Move right until $\{x\}$ on track 5.
8: Write B on track 5.

| 1 | 1 | 0 |  | 0 | 1 | 1 | 0 | 1 | 0 | 0 |  | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 |  | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |

1: If Stack 1 is empty, go to L


| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

1: If Stack 1 is empty, go to L 0 : Write $\times$ on track 5.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\gg$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

1: If Stack 1 is empty, go to L 0 : Write $\times$ on track 5.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\gg$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\gg$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
011

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  | ... |

1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
011

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  | ... |

1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\boldsymbol{>}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |

1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $\gg$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |

1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into V


1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into V
$V=1$


1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
$v=1$

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.

## $V=1$

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | $\mathbf{1}$ | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
5: Move right until $\{x\}$ on track 5.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |  | ... |

1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
5: Move right until $\{x\}$ on track 5.
$V=1$


1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
5: Move right until $\{x\}$ on track 5.
6: Write $B$ on track 5.
$v=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |

1: If Stack 1 is empty, go to L
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
5: Move right until $\{x\}$ on track 5.
6: Write $B$ on track 5.
$v=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\gg$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

1: If Stack 1 is empty, go to L
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 2.
2: Move right.
3: Load current on track 2 into $V$
4: Move left.
5: Move right until $\{x\}$ on track 5.
6: Write B on track 5.
7: If V = <, go to L.

| 1 | 1 | 0 |  | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 0 | 1 |  | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 1 |  |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |



| 1 | 1 | 0 | 0 |  | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 1 | 1 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |



0: Write $x$ on track 5 .

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
|  |  |  |  |  |  |  |  |  |  |  | * |  |  |  |  |  |  | ... |



0 : Write $\times$ on track 5.

| 1 | 1 | 0 | 0 |  |  | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | ... |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | ... |
| $>$ | 0 | 1 | 1 |  | < |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  | ... |



0 : Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 3.

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\boldsymbol{>}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

2: Pop Stack 2 into X.
0: Write $\times$ on track 5.
1: Move left until \{>\} on track 3.
2: Move right until $\{<\}$ on track 3.


|  | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $>$ | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | $>$ | 0 | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | > | 1 | 1 | < |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  | $\times$ |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  | : P | Pop | St | tac | k | 2 i | nt | - x |  |  |  |
|  |  |  |  |  |  |  |  | Wri | ite | $\times$ | on | tr | rack |  |  |  |  |  |
| 1 | 1 |  |  |  |  |  |  | Mov |  |  | t | unt |  |  | ) | t | ack |  |
|  |  |  |  |  |  |  |  | Mov | ve | rig | ht | un | til | 1 | (<) |  | tra |  |
| 0 | 1 | 1 |  |  |  |  |  | Mov | ve |  |  |  |  |  |  |  |  |  |


| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X.
0: Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.



2: Pop Stack 2 into X .
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
$X=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X.
0: Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If $\mathrm{X}=>$, go to 7 .
$X=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X.
0: Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If X = >, go to 7.
6: Write < on track 3
$X=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | 1 | $<$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X .
0: Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If $\mathrm{X}=>$, go to 7.
6: Write < on track 3
$X=1$

| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $>$ | 1 | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 0 | $\mathbf{1}$ | $<$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |  |

2: Pop Stack 2 into X .
0: Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If $\mathrm{X}=>$, go to 7.
6: Write < on track 3
7: Move left until $\{>\}$ on track 3.
$x=1$

| 1 | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\ldots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{>}$ | $\mathbf{1}$ | $\mathbf{1}$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $>$ | $\mathbf{0}$ | $\mathbf{1}$ | $<$ | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
| $>$ | 1 | 1 | $<$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | $\ldots$ |

2: Pop Stack 2 into X.
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until \{<\} on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If X = >, go to 7.
6: Write < on track 3
7: Move left until $\{>\}$ on track 3.


2: Pop Stack 2 into X.
0 : Write $\times$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into X .
5: If $\mathrm{X}=>$, go to 7.
6: Write < on track 3
7: Move left until $\{>\}$ on track 3.
8: Move right until $\{x\}$ on track 5.


2: Pop Stack 2 into X .

0 : Write $x$ on track 5.
1: Move left until $\{>\}$ on track 3.
2: Move right until $\{<\}$ on track 3.
3: Move left.
4: Load current on track 3 into $X$.
5: If $\mathrm{X}=>$, go to 7 .
6: Write < on track 3
7: Move left until $\{>\}$ on track 3.
8: Move right until $\{x\}$ on track 5.



## Completing the Construction

- We've seen how to convert the new WB5 stack commands into WB4 code.
- For this to work, the extra tracks must be set up correctly.
- Add preamble code to the generated WB4 program to do this:

```
Write > to track 2.
Write > to track n.
Move right.
Write < to track 2.
Write < to track n.
Move left.
```


## But Why Stop There?

- Adding finitely many stacks to WB doesn't increase its expressive power.
- What if we added finitely many tapes to WB?
- We now have a programming language controlling
- Multiple tracks per tape,
- Finitely many stacks, and
- Finitely many tapes.



## Introducing WB6

- The programming language WB6 is WB5 with the addition of multiple tapes.
- All tape commands have been updated to specify which tape they apply to.
- If tape unspecified, it's assumed that it's tape 1.


## A WB6 Program for $\operatorname{SEARCH}$

- Recall from Problem Sets 5 and 6 that the language $S E A R C H$ over $\Sigma=\{0,1$, ? $\}$ is the language

$$
\begin{gathered}
\left\{p ? t \mid p, t \in\{0,1\}^{*} \text { and } p\right. \\
\text { is a substring of } t\}
\end{gathered}
$$

- How would we write a WB6 program for SEARCH?
- (For simplicity, we'll assume that the input is properly formatted).


## A WB6 Program for Search


...

## A WB6 Program for Search


...

## A WB6 Program for Search


...

## A WB6 Program for Search


...

## A WB6 Program for Search


...

## A WB6 Program for Search


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## A WB6 Program for Search


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## A WB6 Program for Search


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## A WB6 Program for Search



-     - 


## A WB6 Program for Search


...

## A WB6 Program for Search


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## A WB6 Program for Search


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## A WB6 Program for Search



-     - 


## A WB6 Program for Search



- $\cdot$


## A WB6 Program for $S E A R C H$

// Start
0: Move tape 2 right.
1: If reading ? on tape 1.1 , go to Match.
2: Load curr on tape 1.1 into X .
3: Write X to tape 2.
4: Move tape 1 right.
5: Move tape 2 right.
6: Go to 1.

## A WB6 Program for SEARCH

// Match
7: Move tape 2 left until \{B\}
8: Move tape 2 right.
9: Move tape 1 right.
10: Write \$ to tape 1, track 2. // Acc
11: If B on tape 2, go to Acc. 22: Accept.
12: If $B$ on tape 1, go to Rej.
13: Load tape 1, track 1 into $\mathrm{x} . / / \mathrm{Rej}$
14: Load tape 2 into $Y$.
15: If $\mathrm{X}=\mathrm{Y}, \mathrm{go}$ to 17.
16: Go to Mismatch.
17: Move tape 1 right.
18: Move tape 2 right.
19: Go to 11.
// Mismatch
20: Move tape 1.2 left until \{\$\}
21: Go to Match.

23: Reject.

## Oh, Come On Already...

- Theorem: A language is recursively enumerable iff there is a WB6 program for it.
- We can really supercharge these languages without increasing our power!
- As with before, the construction will convert WB6 programs into WB5 programs.



## The Key Idea

- Represent an infinite tape with two stacks.



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| $\mathbf{E}$ |
| :---: |
| $\mathbf{F}$ |
| $\mathbf{G}$ |
| $\mathbf{H}$ |

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## A Sketch of the Construction

- At the start of the program, copy the contents of the initial tape into a pair of stacks that will henceforth represent the first tape.
- Convert all motion operations into stack manipulation operations to push and pop values from the appropriate stacks.
- Use variables to hold temporary values (for example, when moving the top of one stack to another).


## 0: Move tape 1 right.

0: If stack 1 R is empty, go to 2.
1: Go to 3.
2: Push B onto stack 1R.
3: Pop stack 1R into X .
4: Push X onto stack 1L.


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## What Else Can We Add?

- Function call and return.
- Have a stack to use as the call stack.
- Calling a function pushes the index of the instruction to which it should return.
- Returning pops the stack and jumps back.
- Named variables.
- Have a tape storing a sequence of values of the form name: value.
- Can read and write values from the tape.
- Pointers
- Have variables hold the names of other variables.
- Primitive types and arithmetic.
- Design subroutines for addition, subtraction, etc.
- Apply them to named variables.
- Pretty much any feature of any major programming language.


## The Conversion Back Down

- From WB6 to WB5:
- Add in two stacks per tape used.
- Replace all tape operations with appropriate stack manipulations.
- From WB5 to WB4:
- Add in one track per stack, plus one extra track.
- Replace all stack operations with appropriate manipulations of those tracks.


## The Conversion Back Down

- From WB4 to WB3:
- Expand the tape alphabet to include symbols for all track combinations.
- Replace all references to track symbols with cascading if's for each possible case.
- From WB3 to WB2:
- Replicate the code once for each possible assignment to variables.
- Hardcode in statements referencing variables.
- Replace variable manipulation code with code to jump to the appropriate copy.


## The Conversion Back Down

- From WB2 to WB:
- Expand out move ... until statements by replacing them with cascading if statements.
- From WB to Turing machines:
- Replace each statement with the appropriate Turing machine gadget.


## The Conversion Back Down

- The total conversion of a WB6 program using variables, multiple tracks, multiple stacks, and multiple tapes might produce an enormous Turing machine!
- But that said, the result is still a Turing machine.
- Turing machines are simple, yet have enormous computational power.


## Just how powerful are Turing machines?

## Effective Computation

- An effective method of computation is a form of computation with the following properties:
- The computation consists of a set of steps.
- There are fixed rules governing how one step leads to the next.
- Any computation that yields an answer does so in finitely many steps.
- Any computation that yields an answer always yields the correct answer.


## The Church-Turing Thesis states that

Every effective method of computation is either equivalent to or weaker than a Turing machine.

This statement cannot be proven or disproven, but is widely considered true.

## Problems

Solvable by Any Feasible Computing Machine

All Languages

## Next Time

- Encodings
- How do we do computations over arbitrary objects?
- The Universal Turing Machine
- A Turing machine for running other Turing machines.
- Nondeterministic Turing Machines
- What happens when we supercharge a TM? What does this even mean?
- $R$ and RE Languages
- A finer gradation within the RE languages.

