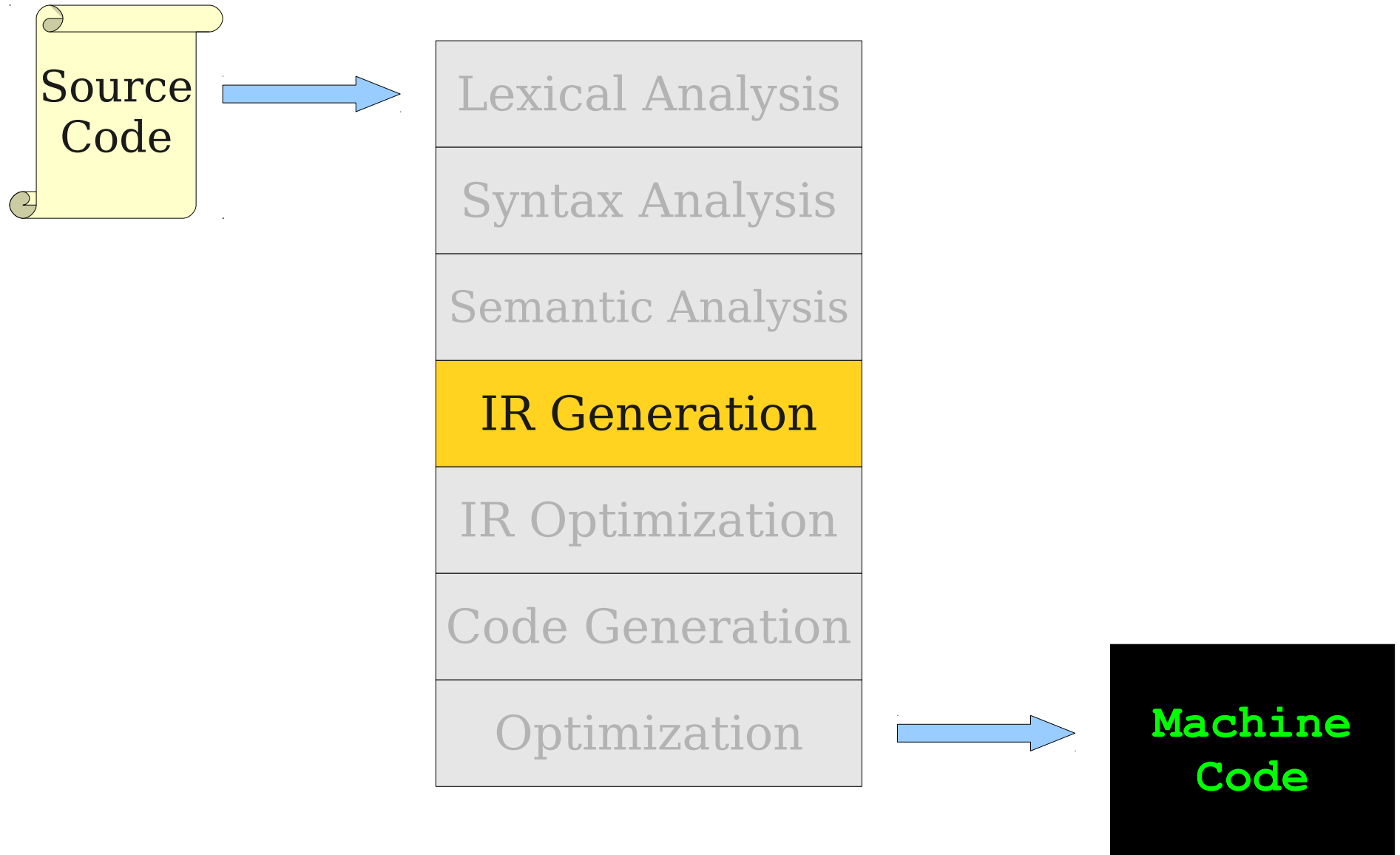


IR Optimization

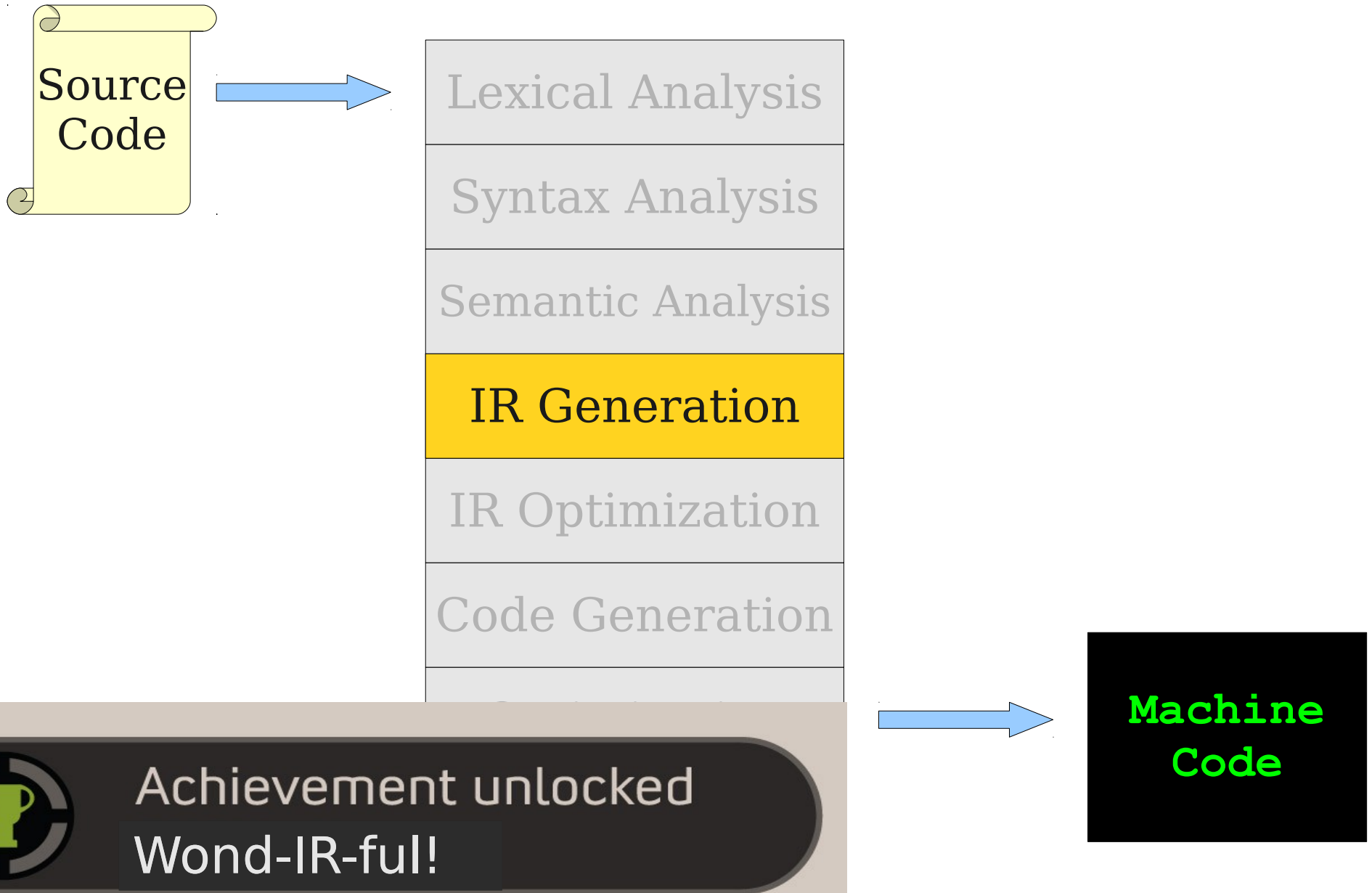
Announcements

- Programming Project 3 due Monday at 11:59PM.
- Programming Project 3 checkpoint graded; feedback emailed out.
 - Please let us know ASAP if you haven't heard back yet!

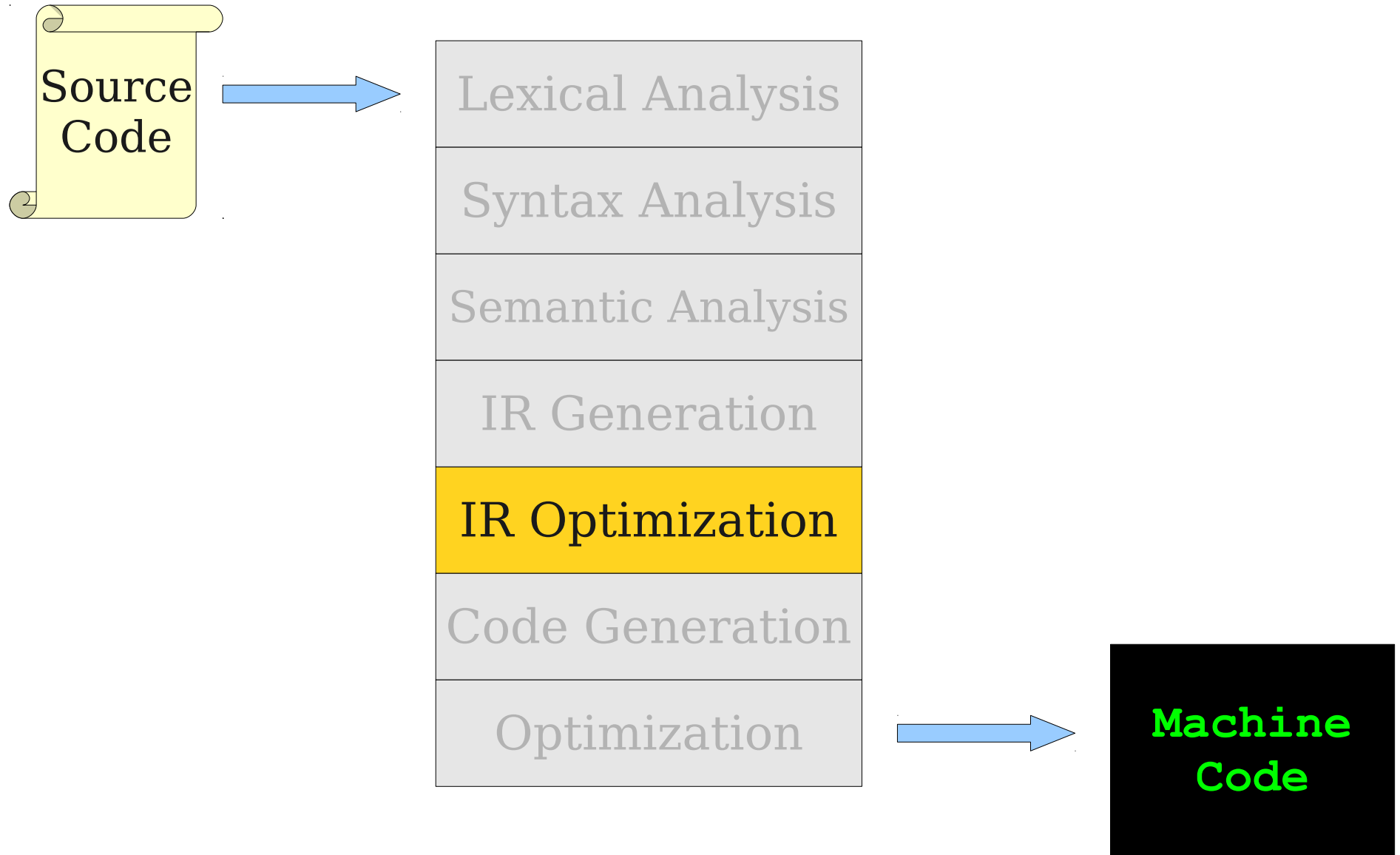
Where We Are



Where We Are



Where We Are



IR Optimization

- **Goal:** Improve the IR generated by the previous step to take better advantage of resources.
- One of the most important and complex parts of any modern compiler.
- A very active area of research.
- There is a whole class (CS243) dedicated to this material.

Sources of Optimization

- In order to optimize our IR, we need to understand why it can be improved in the first place.
- **Reason one:** IR generation introduces redundancy.
 - A naïve translation of high-level language features into IR often introduces subcomputations.
 - Those subcomputations can often be sped up, shared, or eliminated.
- **Reason two:** Programmers are lazy.
 - Code executed inside of a loop can often be factored out of the loop.
 - Language features with side effects often used for purposes other than those side effects.

Optimizations from IR Generation

```
int x;  
int y;  
bool b1;  
bool b2;  
bool b3;
```

```
b1 = x + x < y  
b2 = x + x == y  
b3 = x + x > y
```


Optimizations from IR Generation

```
int x;  
int y;  
bool b1;  
bool b2;  
bool b3;
```

```
b1 = x + x < y  
b2 = x + x == y  
b3 = x + x > y
```

```
_t0 = x + x;  
_t1 = y;  
b1 = _t0 < _t1;
```

```
_t2 = x + x;  
_t3 = y;  
b2 = _t2 == _t3;
```

```
_t4 = x + x;  
_t5 = y;  
b3 = _t5 < _t4;
```

Optimizations from IR Generation

```
int x;  
int y;  
bool b1;  
bool b2;  
bool b3;  
  
b1 = x + x < y  
b2 = x + x == y  
b3 = x + x > y
```

```
_t0 = x + x;  
_t1 = y;  
b1 = _t0 < _t1;  
  
_t2 = x + x;  
_t3 = y;  
b2 = _t2 == _t3;  
  
_t4 = x + x;  
_t5 = y;  
b3 = _t5 < _t4;
```

Optimizations from IR Generation

```
int x;  
int y;  
bool b1;  
bool b2;  
bool b3;  
  
b1 = x + x < y  
b2 = x + x == y  
b3 = x + x > y
```

```
_t0 = x + x;  
_t1 = y;  
b1 = _t0 < _t1;  
  
b2 = _t0 == _t1;  
  
b3 = _t0 > _t1;
```

```
while (x < y + z) {  
    x = x - y;  
}
```



Optimizations from Lazy Coders

```
while (x < y + z) {  
    x = x - y;  
}
```

```
_L0:  
    _t0 = y + z;  
    _t1 = x < _t0;  
    IfZ _t1 Goto _L1;  
    x = x - y;  
    Goto _L0;  
_L1:
```

Optimizations from Lazy Coders

```
while (x < y + z) {  
    x = x - y;  
}
```

```
_L0:  
    _t0 = y + z;  
    _t1 = x < _t0;  
    IfZ _t1 Goto _L1;  
    x = x - y;  
    Goto _L0;  
_L1:
```

Optimizations from Lazy Coders

```
while (x < y + z) {  
    x = x - y;  
}
```

```
    _t0 = y + z;  
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    _t1 = x < _t0;  
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    Goto _L0;  
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```

Optimizations from Lazy Coders

```
while (x < y + z) {  
    x = x - y;  
}
```

```
    _t0 = y + z;  
_L0:  
    _t1 = x < _t0;  
    IfZ _t1 Goto _L1;  
    x = x - y;  
    Goto _L0;  
_L1:
```


A Note on Terminology

- The term “optimization” implies looking for an “optimal” piece of code for a program.
- This is, in general, undecidable.
 - e.g. create a program that can be simplified iff some other program halts.
- Our goal will be IR *improvement* rather than IR *optimization*.

The Challenge of Optimization

- A good optimizer
 - Should never change the observable behavior of a program.
 - Should produce IR that is as efficient as possible.
 - Should not take too long to process inputs.
- Unfortunately:
 - Even good optimizers sometimes introduce bugs into code.
 - Optimizers often miss “easy” optimizations due to limitations of their algorithms.
 - Almost all interesting optimizations are **NP**-hard or undecidable.

What are we Optimizing?

- Optimizers can try to improve code usage with respect to many observable properties.
- What are some quantities we might want to optimize?

What are we Optimizing?

- Optimizers can try to improve code usage with respect to many observable properties.
- What are some quantities we might want to optimize?
- **Runtime** (make the program as fast as possible at the expense of time and power)
- **Memory usage** (generate the smallest possible executable at the expense of time and power)
- **Power consumption** (choose simple instructions at the expense of speed and memory usage)
- Plus a lot more (minimize function calls, reduce use of floating-point hardware, etc.)

IR Optimization vs Code Optimization

- There is not always a clear distinction between what belongs to “IR optimization” versus “code optimization.”
- Typically:
 - IR optimizations try to perform simplifications that are valid across all machines.
 - Code optimizations try to improve performance based on the specifics of the machine.
- Some optimizations are somewhere in-between:
 - Replacing $x * 0.5$ with $x / 2$

Overview of IR Optimization

- **Formalisms and Terminology** (Today)
 - Control-flow graphs.
 - Basic blocks.
- **Local optimizations** (Today)
 - Speeding up small pieces of a function.
- **Global optimizations** (Monday)
 - Speeding up functions as a whole.
- **The dataflow framework** (Monday/Wednesday)
 - Defining and implementing a wide class of optimizations.

Formalisms and Terminology

Analyzing a Program

- In order to optimize a program, the compiler has to be able to reason about the properties of that program.
- An analysis is called **sound** if it never asserts an incorrect fact about a program.
- All the analyses we will discuss in this class are sound.
 - *(Why?)*

Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
  
Print(x);
```

Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
  
Print (x) ;
```

Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
Print (x) ;
```

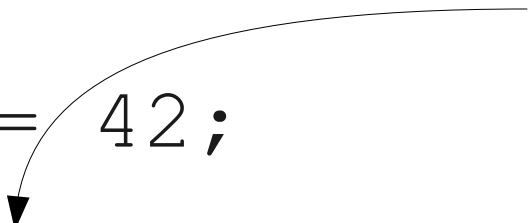
“At this point in the program, **x** holds some integer value.”



Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
Print (x) ;
```

“At this point in the program, **x** is either 137 or 42”



Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
Print (x) ;
```

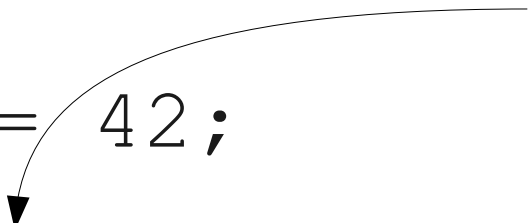
“At this point in the program, **x** is 137”



Soundness

```
int x;  
int y;  
  
if (y < 5)  
    x = 137;  
else  
    x = 42;  
Print (x) ;
```

“At this point in the program, **x** is either 137, 42, or 271”



Semantics-Preserving Optimizations

- An optimization is **semantics-preserving** if it does not alter the semantics of the original program.
- Examples:
 - Eliminating unnecessary temporary variables.
 - Computing values that are known statically at compile-time instead of runtime.
 - Evaluating constant expressions outside of a loop instead of inside.
- Non-examples:
 - Replacing bubble sort with quicksort.
- The optimizations we will consider in this class are all semantics-preserving.

A Formalism for IR Optimization

- Every phase of the compiler uses some new abstraction:
 - Scanning uses regular expressions.
 - Parsing uses CFGs.
 - Semantic analysis uses proof systems and symbol tables.
 - IR generation uses ASTs.
- In optimization, we need a formalism that captures the structure of a program in a way amenable to optimization.

Visualizing IR

```
main:
    BeginFunc 40;
    _tmp0 = LCall _ReadInteger;
    a = _tmp0;
    _tmp1 = LCall _ReadInteger;
    b = _tmp1;
_L0:
    _tmp2 = 0;
    _tmp3 = b == _tmp2;
    _tmp4 = 0;
    _tmp5 = _tmp3 == _tmp4;
    IfZ _tmp5 Goto _L1;
    c = a;
    a = b;
    _tmp6 = c % a;
    b = _tmp6;
    Goto _L0;
_L1:
    PushParam a;
    LCall _PrintInt;
    PopParams 4;
    EndFunc;
```

Visualizing IR

```
main:
  BeginFunc 40;
  _tmp0 = LCall _ReadInteger;
  a = _tmp0;
  _tmp1 = LCall _ReadInteger;
  b = _tmp1;
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  _tmp2 = 0;
  _tmp3 = b == _tmp2;
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  _tmp5 = _tmp3 == _tmp4;
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  c = a;
  a = b;
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  PushParam a;
  LCall _PrintInt;
  PopParams 4;
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```

Visualizing IR

```
main:
  BeginFunc 40;
  _tmp0 = LCall _ReadInteger;
  a = _tmp0;
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  b = _tmp1;
_L0:
  _tmp2 = 0;
  _tmp3 = b == _tmp2;
  _tmp4 = 0;
  _tmp5 = _tmp3 == _tmp4;
  IfZ _tmp5 Goto _L1;
  c = a;
  a = b;
  _tmp6 = c % a;
  b = _tmp6;
  Goto _L0;
_L1:
  PushParam a;
  LCall _PrintInt;
  PopParams 4;
  EndFunc;
```

```
_tmp0 = LCall _ReadInteger;
a = _tmp0;
_tmp1 = LCall _ReadInteger;
b = _tmp1;
```

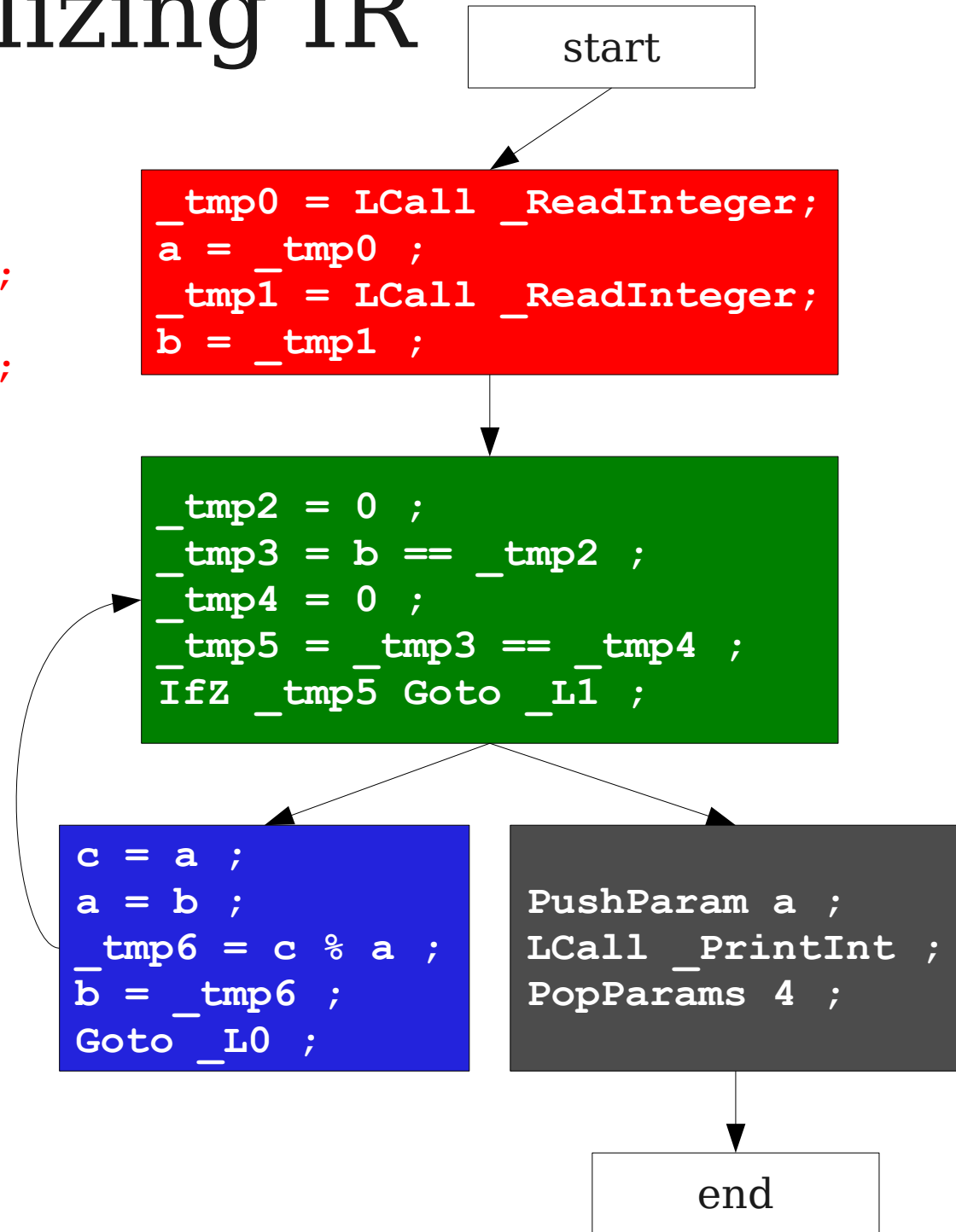
```
_tmp2 = 0;
_tmp3 = b == _tmp2;
_tmp4 = 0;
_tmp5 = _tmp3 == _tmp4;
IfZ _tmp5 Goto _L1;
```

```
c = a;
a = b;
_tmp6 = c % a;
b = _tmp6;
Goto _L0;
```

```
PushParam a;
LCall _PrintInt;
PopParams 4;
```

Visualizing IR

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  IfZ _tmp5 Goto _L1;
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  a = b;
  _tmp6 = c % a;
  b = _tmp6;
  Goto _L0;
_L1:
  PushParam a;
  LCall _PrintInt;
  PopParams 4;
  EndFunc;
```



Basic Blocks

- A **basic block** is a sequence of IR instructions where
 - There is exactly one spot where control enters the sequence, which must be at the start of the sequence.
 - There is exactly one spot where control leaves the sequence, which must be at the end of the sequence.
- Informally, a sequence of instructions that always execute as a group.

Control-Flow Graphs

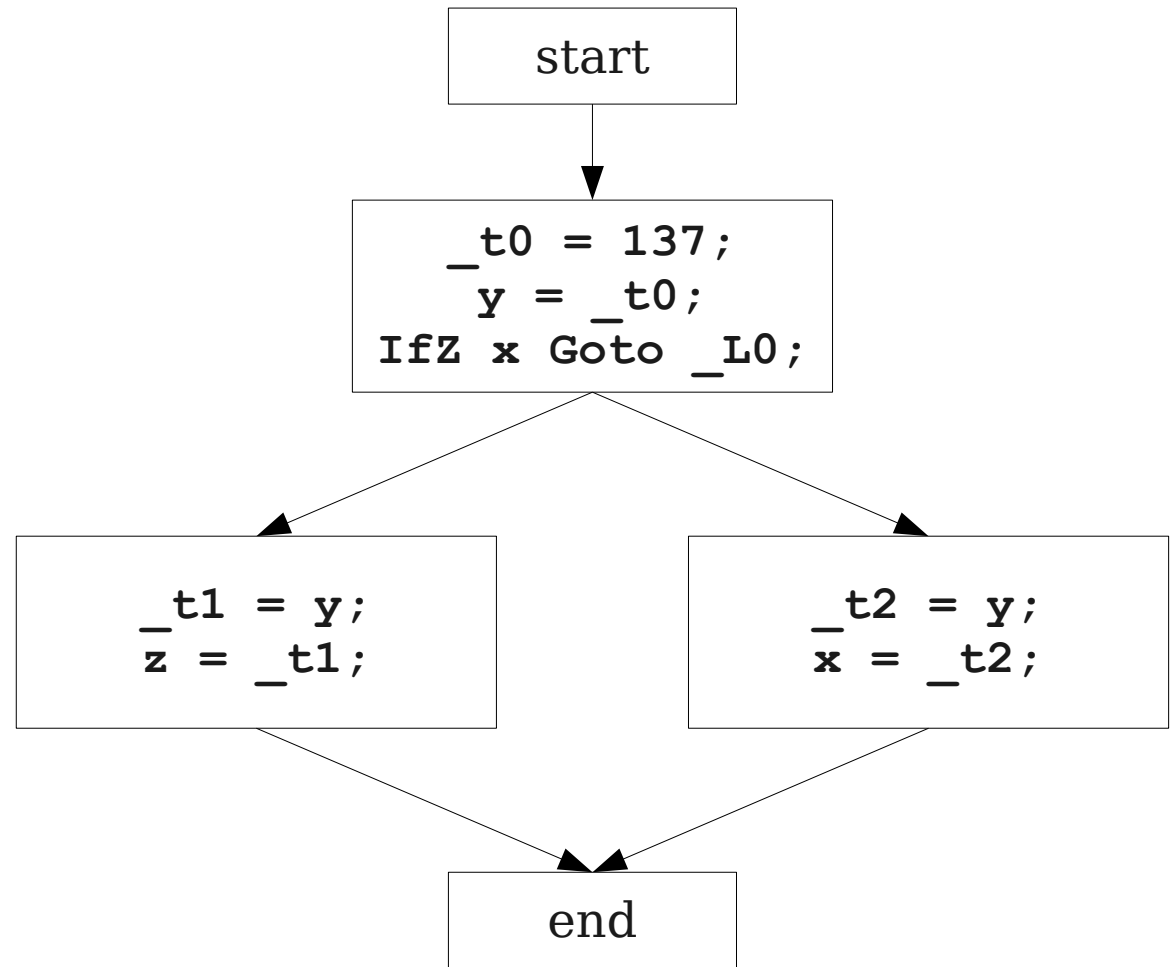
- A **control-flow graph** (CFG) is a graph of the basic blocks in a function.
 - The term CFG is overloaded – from here on out, we'll mean “control-flow graph” and not “context-free grammar.”
- Each edge from one basic block to another indicates that control can flow from the end of the first block to the start of the second block.
- There is a dedicated node for the start and end of a function.

Types of Optimizations

- An optimization is **local** if it works on just a single basic block.
- An optimization is **global** if it works on an entire control-flow graph.
- An optimization is **interprocedural** if it works across the control-flow graphs of multiple functions.
 - We won't talk about this in this course.

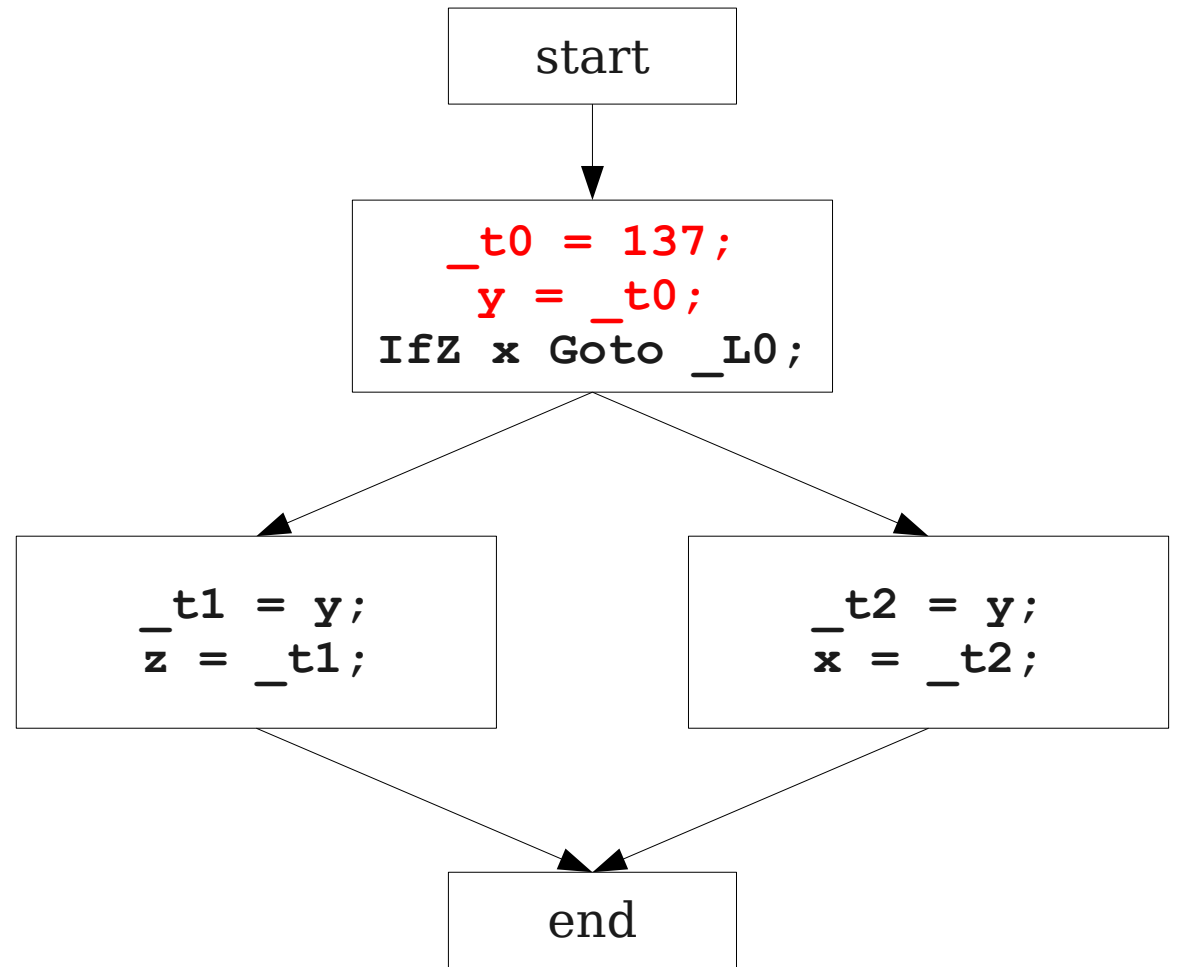
Local Optimizations

```
int main() {  
    int x;  
    int y;  
    int z;  
  
    y = 137;  
    if (x == 0)  
        z = y;  
    else  
        x = y;  
}
```



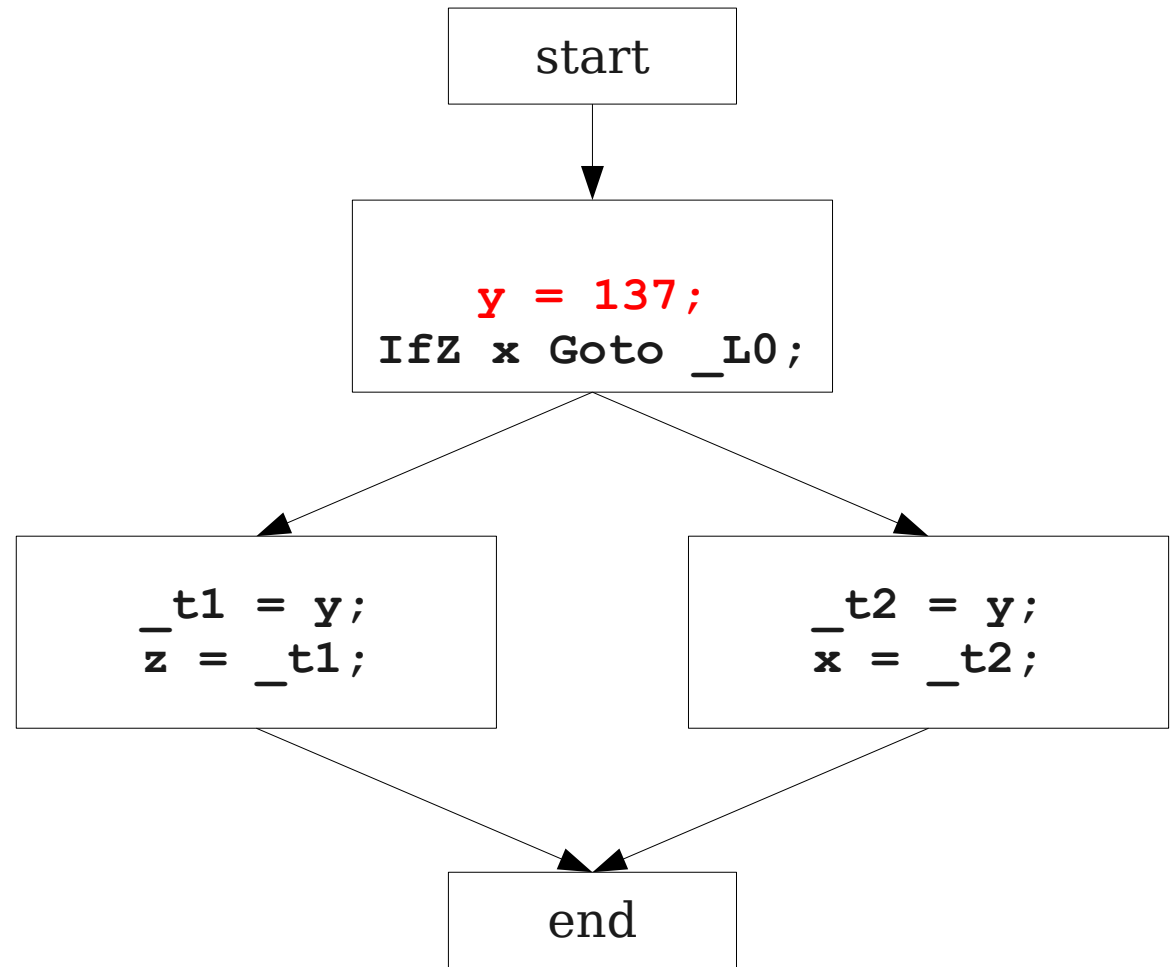
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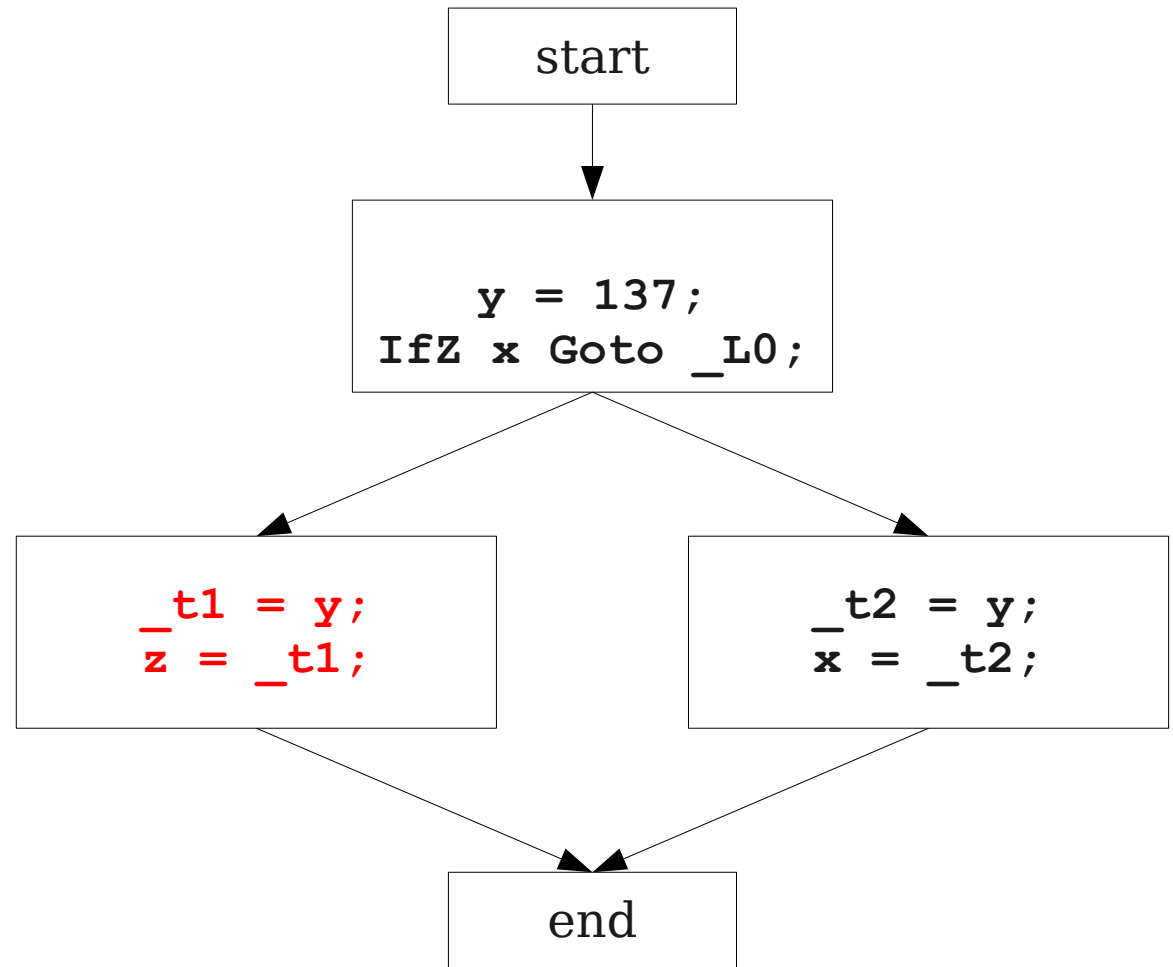
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}
```



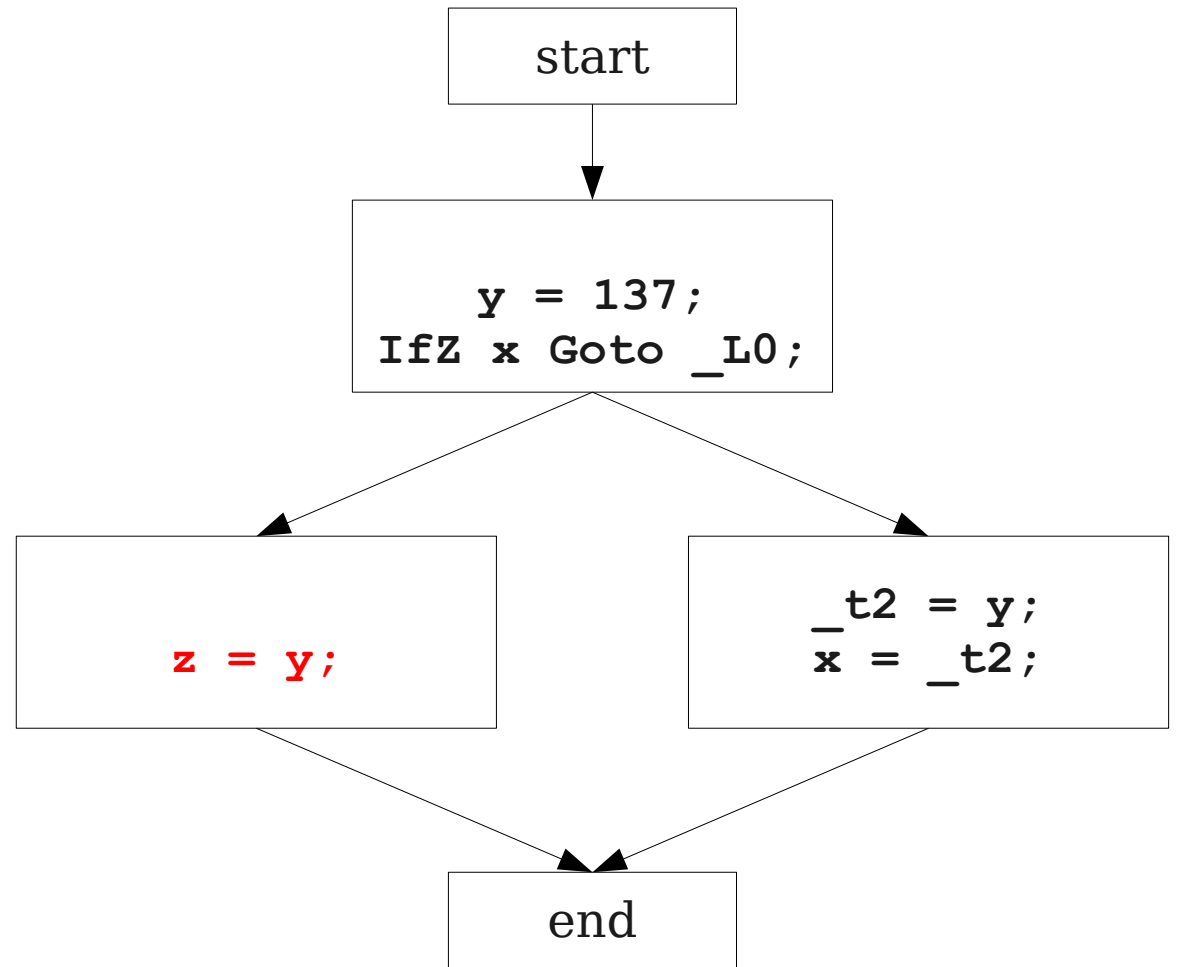
Local Optimizations

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}
```



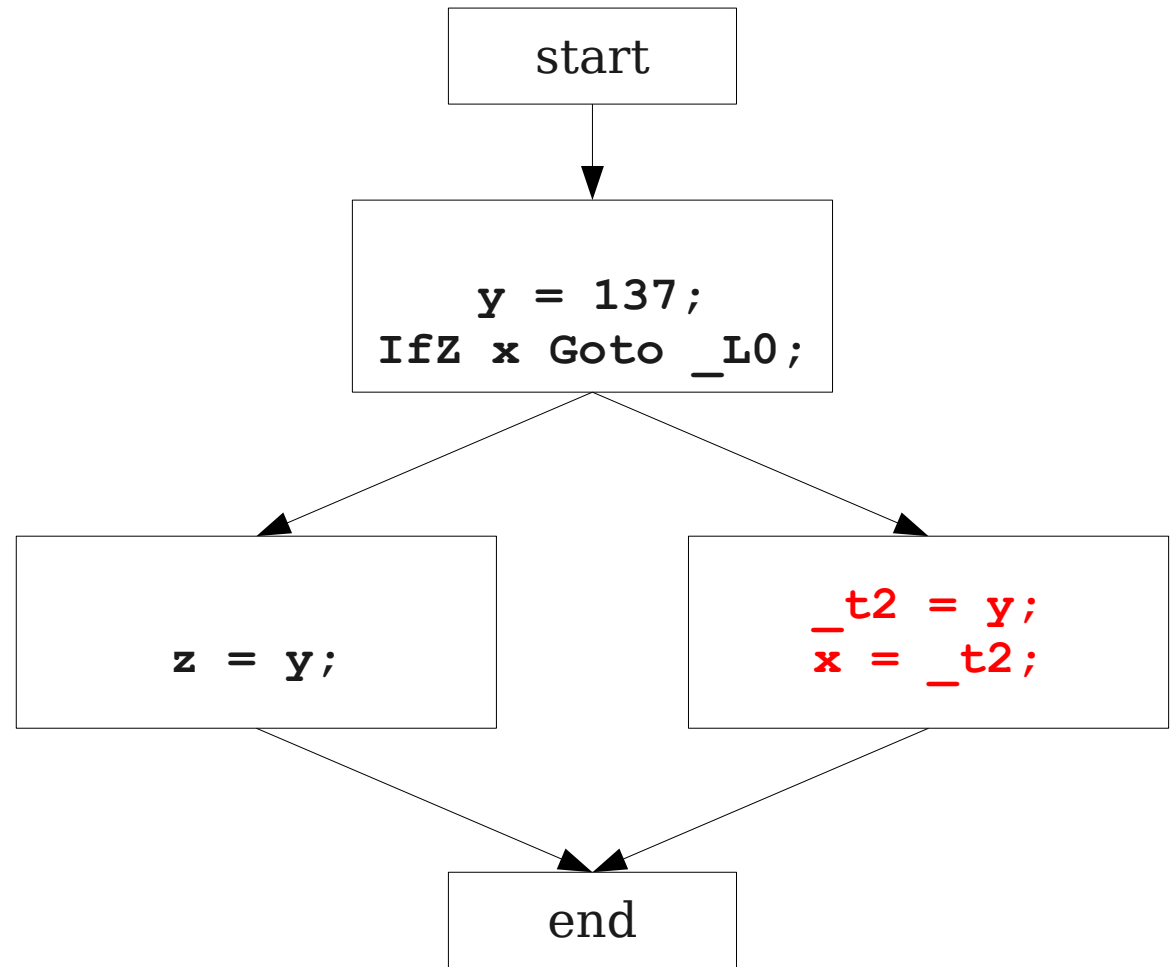
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```



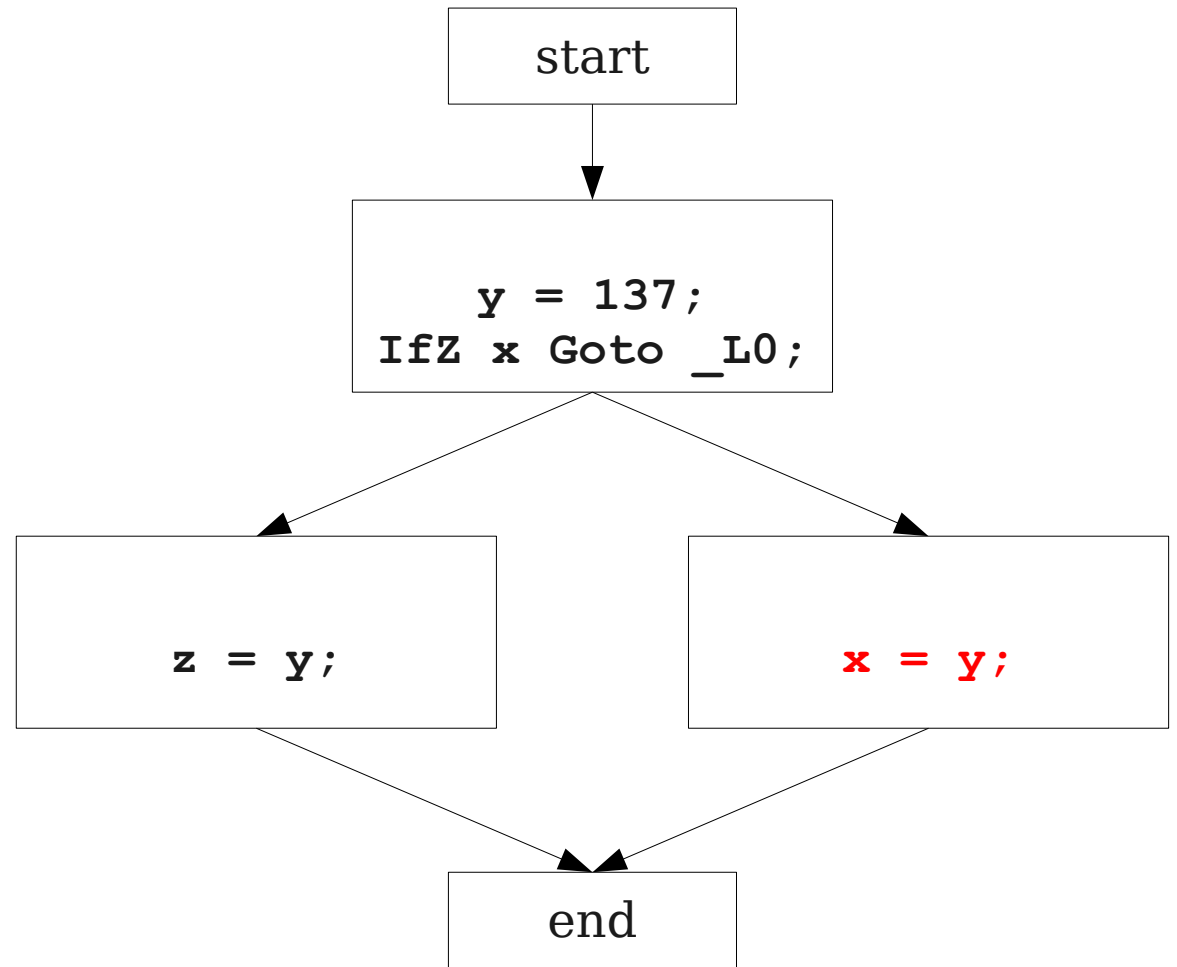
Local Optimizations

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    if (x == 0)  
        z = y;  
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}
```



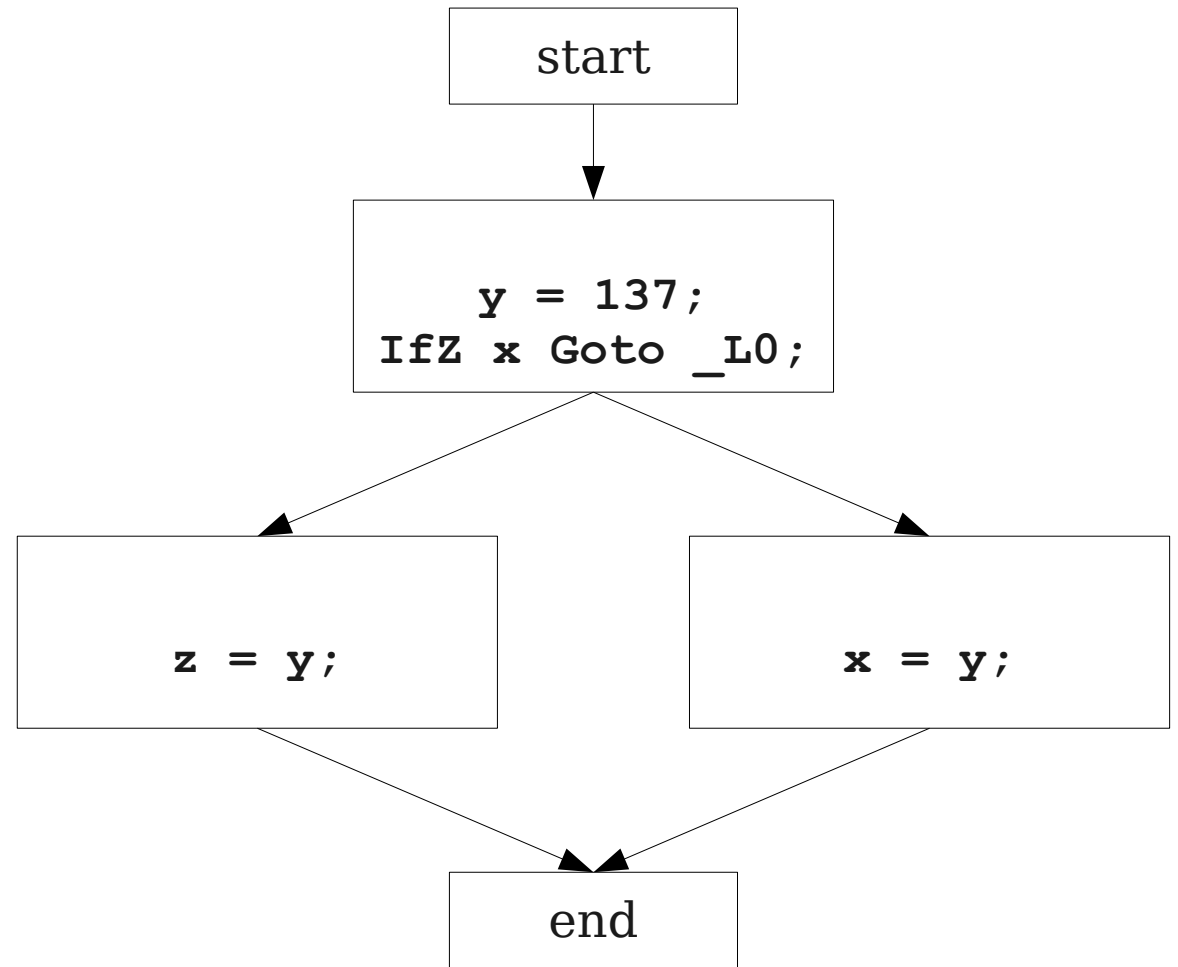
Local Optimizations

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    int z;  
  
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```



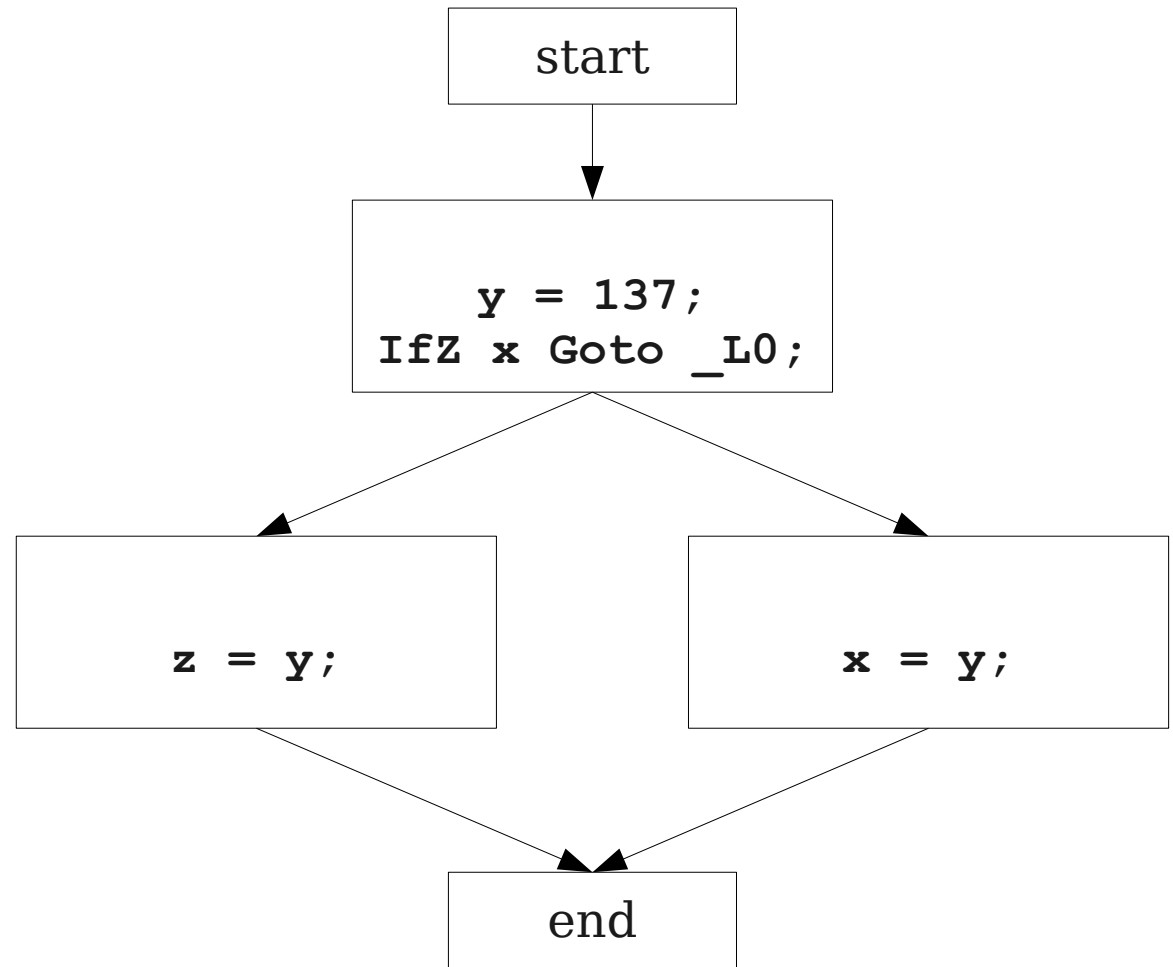
Local Optimizations

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```



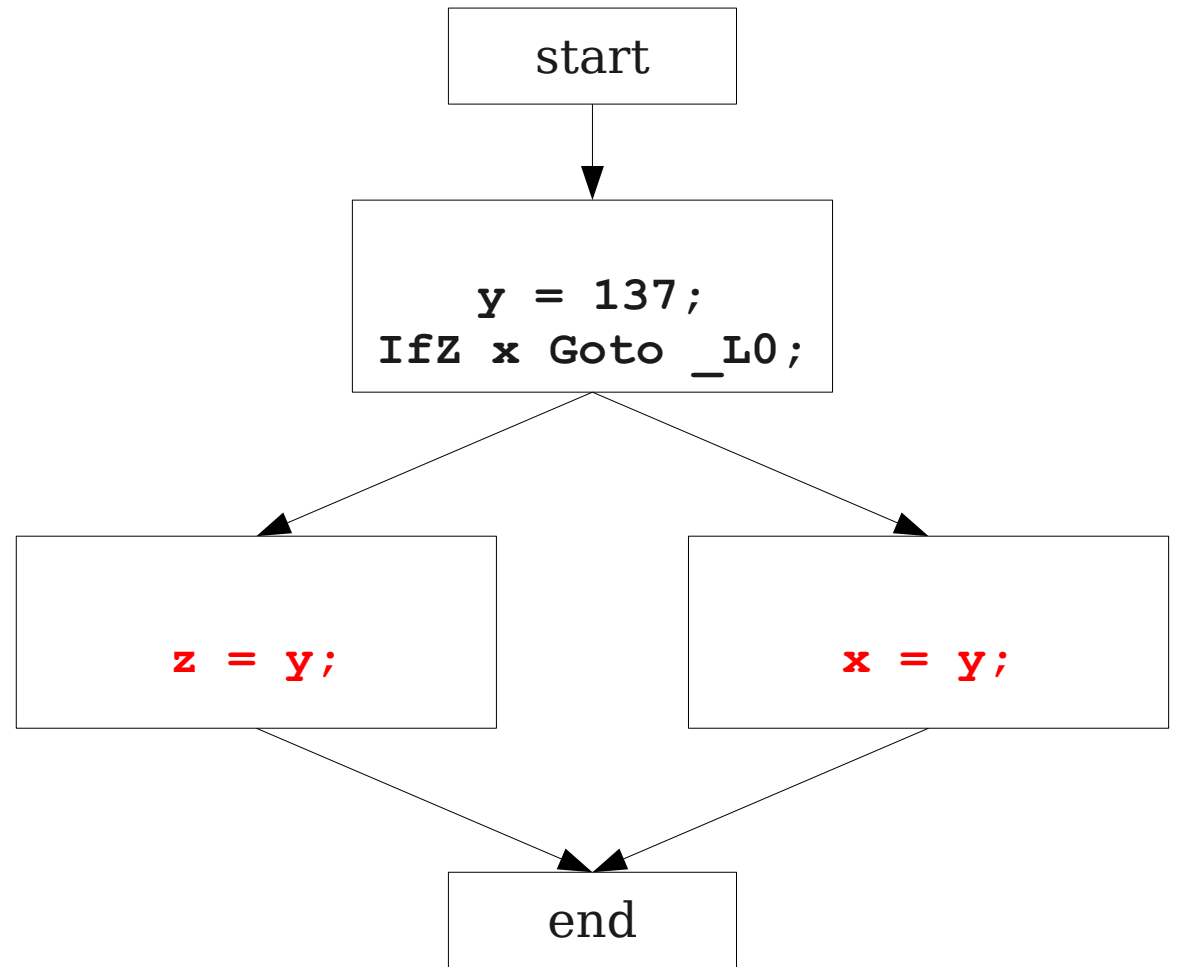
Global Optimizations

```
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    int z;  
  
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}
```



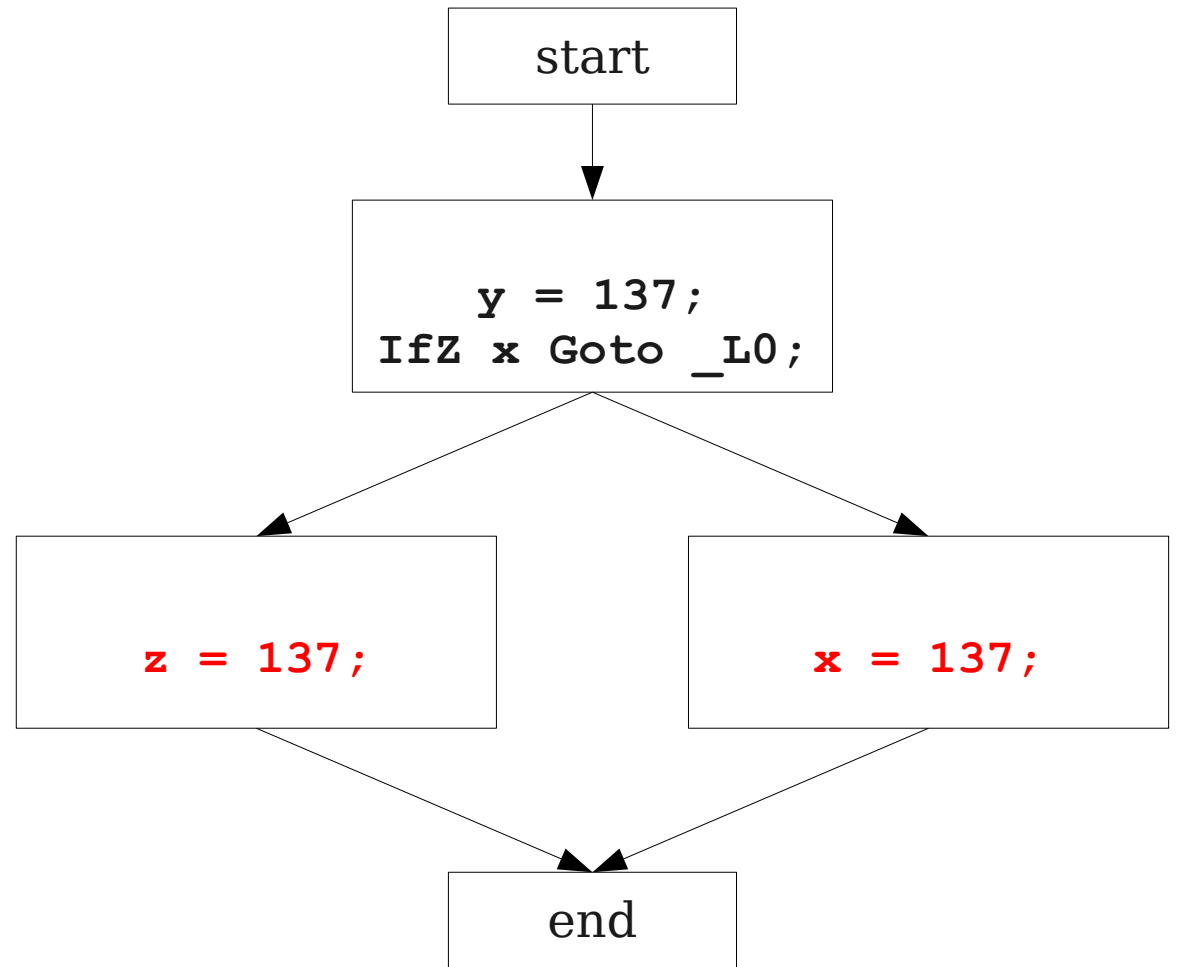
Global Optimizations

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```



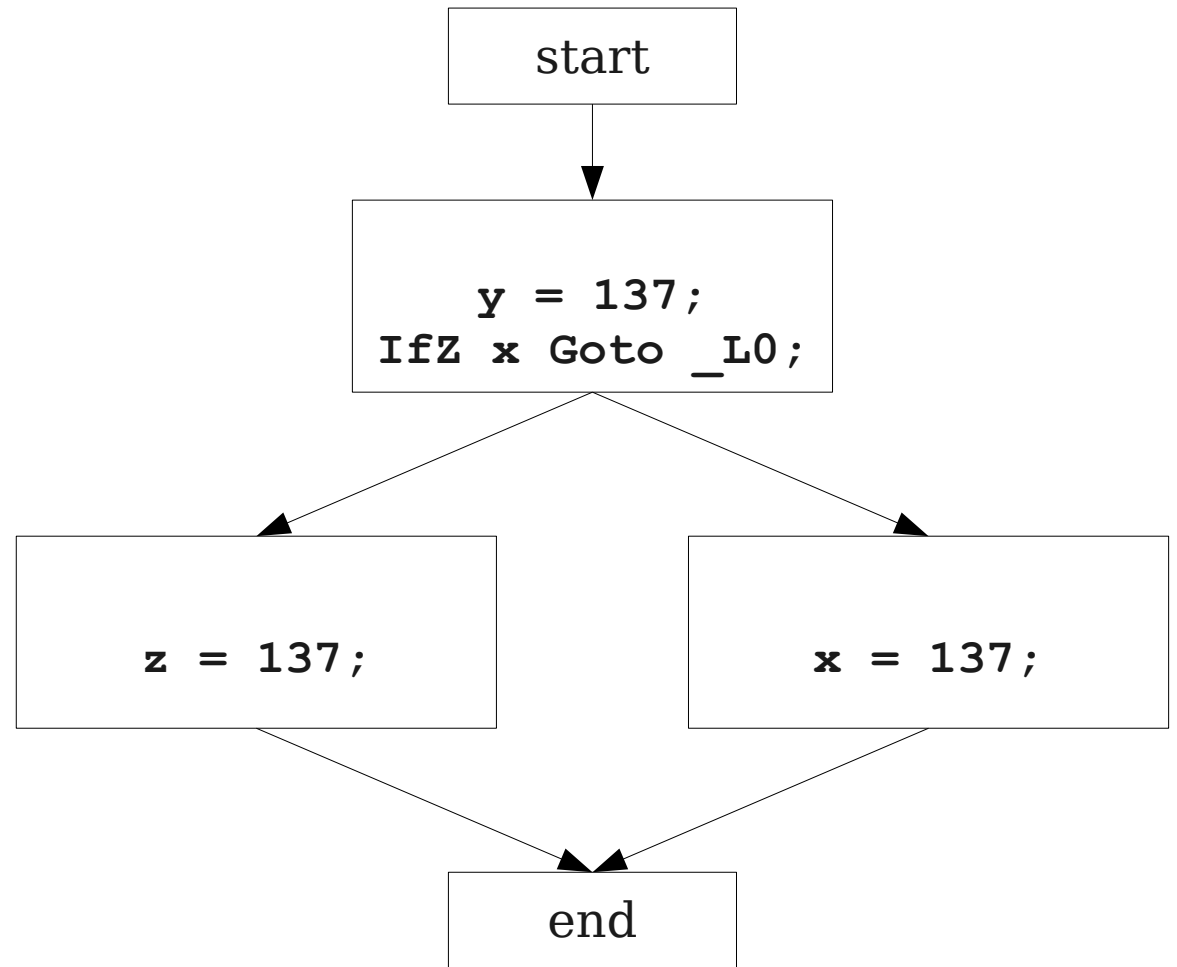
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    else  
        x = y;  
}
```



Local Optimizations

Common Subexpression Elimination

```
Object x;
```

```
int a;
```

```
int b;
```

```
int c;
```

```
x = new Object;
```

```
a = 4;
```

```
c = a + b;
```

```
x.fn(a + b);
```

Common Subexpression Elimination

```
Object x;  
int a;  
int b;  
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c = a + b;  
x.fn(a + b);
```

```
_tmp0 = 4 ;  
PushParam _tmp0 ;  
_tmp1 = LCall _Alloc ;  
PopParams 4 ;  
_tmp2 = Object ;  
*(_tmp1) = _tmp2 ;  
x = _tmp1 ;  
_tmp3 = 4 ;  
a = _tmp3 ;  
_tmp4 = a + b ;  
c = _tmp4 ;  
_tmp5 = a + b ;  
_tmp6 = *(x) ;  
_tmp7 = *(_tmp6) ;  
PushParam _tmp5 ;  
PushParam x ;  
ACall _tmp7 ;  
PopParams 8 ;
```

Common Subexpression Elimination

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Object x;  
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int b;  
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_tmp5 = _tmp4 ;  
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```

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x = _tmp1 ;  
_tmp3 = _tmp0 ;  
a = _tmp3 ;  
_tmp4 = a + b ;  
c = _tmp4 ;  
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int b;  
int c;  
  
x = new Object;  
a = 4;  
c = a + b;  
x.fn(a + b);
```

```
_tmp0 = 4 ;  
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PopParams 4 ;  
_tmp2 = Object ;  
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_tmp4 = a + b ;  
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_tmp6 = *(x) ;  
_tmp7 = *(_tmp6) ;  
PushParam _tmp5 ;  
PushParam x ;  
ACall _tmp7 ;  
PopParams 8 ;
```

Common Subexpression Elimination

```
Object x;  
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_tmp6 = *(x) ;  
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PushParam _tmp5 ;  
PushParam x ;  
ACall _tmp7 ;  
PopParams 8 ;
```

Common Subexpression Elimination

- If we have two variable assignments

$$\mathbf{v}_1 = \mathbf{a} \text{ op } \mathbf{b}$$

...

$$\mathbf{v}_2 = \mathbf{a} \text{ op } \mathbf{b}$$

and the values of \mathbf{v}_1 , \mathbf{a} , and \mathbf{b} have not changed between the assignments, rewrite the code as

$$\mathbf{v}_1 = \mathbf{a} \text{ op } \mathbf{b}$$

...

$$\mathbf{v}_2 = \mathbf{v}_1$$

- Eliminates useless recalculation.
- Paves the way for later optimizations.

Copy Propagation

```
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PushParam _tmp4 ;  
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PopParams 8 ;
```


Copy Propagation

- If we have a variable assignment

$$v_1 = v_2$$

then as long as v_1 and v_2 are not reassigned, we can rewrite expressions of the form

$$a = \dots v_1 \dots$$

as

$$a = \dots v_2 \dots$$

provided that such a rewrite is legal.

- This will help immensely later on, as you'll see.

Dead Code Elimination

```
Object x;  
int a;  
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a = 4;  
c = a + b;  
x.fn(a + b);
```

```
_tmp0 = 4 ;  
PushParam _tmp0 ;  
_tmp1 = LCall _Alloc ;  
PopParams 4 ;  
_tmp2 = Object ;  
*(_tmp1) = _tmp2 ;
```

```
_tmp4 = _tmp0 + b ;
```

```
_tmp7 = *(_tmp2) ;  
PushParam _tmp4 ;  
PushParam _tmp1 ;  
ACall _tmp7 ;  
PopParams 8 ;
```

Dead Code Elimination

- An assignment to a variable v is called **dead** if the value of that assignment is never read anywhere.
- **Dead code elimination** removes dead assignments from IR.
- Determining whether an assignment is dead depends on what variable is being assigned to and when it's being assigned.

For Comparison

```
_tmp0 = 4 ;  
PushParam _tmp0 ;  
_tmp1 = LCall _Alloc ;  
PopParams 4 ;  
_tmp2 = Object ;  
*(_tmp1) = _tmp2 ;  
x = _tmp1 ;  
_tmp3 = 4 ;  
a = _tmp3 ;  
_tmp4 = a + b ;  
c = _tmp4 ;  
_tmp5 = a + b ;  
_tmp6 = *(x) ;  
_tmp7 = *(_tmp6) ;  
PushParam _tmp5 ;  
PushParam x ;  
ACall _tmp7 ;  
PopParams 8 ;
```

```
_tmp0 = 4 ;  
PushParam _tmp0 ;  
_tmp1 = LCall _Alloc ;  
PopParams 4 ;  
_tmp2 = Object ;  
*(_tmp1) = _tmp2 ;  
_tmp4 = _tmp0 + b ;  
_tmp7 = *(_tmp2) ;  
PushParam _tmp4 ;  
PushParam _tmp1 ;  
ACall _tmp7 ;  
PopParams 8 ;
```


Applying Local Optimizations

- The different optimizations we've seen so far all take care of just a small piece of the optimization.
 - Common subexpression elimination eliminates unnecessary statements.
 - Copy propagation helps identify dead code.
 - Dead code elimination removes statements that are no longer needed.
- To get maximum effect, we may have to apply these optimizations numerous times.

Applying Local Optimizations

```
b = a * a;  
c = a * a;  
d = b + c;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = a * a;  
d = b + c;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = a * a;  
d = b + c;  
e = b + b;
```

Common Subexpression Elimination

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + c;  
e = b + b;
```

Common Subexpression Elimination

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + c;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + c;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + c;  
e = b + b;
```

Copy Propagation

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = b + b;
```

Copy Propagation

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = b + b;
```

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = b + b;
```

Common Subexpression Elimination (Again)

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = d;
```

Common Subexpression Elimination (Again)

Applying Local Optimizations

```
b = a * a;  
c = b;  
d = b + b;  
e = d;
```

Other Types of Local Optimization

- **Arithmetic Simplification**

- Replace “hard” operations with easier ones.
- e.g. rewrite $\mathbf{x} = 4 * \mathbf{a};$ as $\mathbf{x} = \mathbf{a} \ll 2;$

- **Constant Folding**

- Evaluate expressions at compile-time if they have a constant value.
- e.g. rewrite $\mathbf{x} = 4 * 5;$ as $\mathbf{x} = 20;.$

Implementing Local Optimization

Optimizations and Analyses

- Most optimizations are only possible given some analysis of the program's behavior.
- In order to implement an optimization, we will talk about the corresponding program analyses.

Available Expressions

- Both common subexpression elimination and copy propagation depend on an analysis of the **available expressions** in a program.
- An expression is called **available** if some variable in the program holds the value of that expression.
- In common subexpression elimination, we replace an available expression by the variable holding its value.
- In copy propagation, we replace the use of a variable by the available expression it holds.

Finding Available Expressions

- Initially, no expressions are available.
- Whenever we execute a statement $\mathbf{a = b + c}$:
 - Any expression holding \mathbf{a} is invalidated.
 - The expression $\mathbf{a = b + c}$ becomes available.
- **Idea:** Iterate across the basic block, beginning with the empty set of expressions and updating available expressions at each variable.

Available Expressions

$a = b;$

$c = b;$

$d = a + b;$

$e = a + b;$

$d = b;$

$f = a + b;$

Available Expressions

{ }

a = b;

c = b;

d = a + b;

e = a + b;

d = b;

f = a + b;

Available Expressions

```
    { }  
    a = b;  
  { a = b }  
    c = b;  
  
d = a + b;  
  
e = a + b;  
  
    d = b;  
  
f = a + b;
```

Available Expressions

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

e = a + b;

d = b;

f = a + b;

Available Expressions

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

d = b;

f = a + b;

Available Expressions

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

f = a + b;

Available Expressions

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

Available Expressions

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = b;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = a + b;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = d;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = d;

{ a = b, c = b, d = a + b, e = a + b }

d = b;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = d;

{ a = b, c = b, d = a + b, e = a + b }

d = a;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = d;

{ a = b, c = b, d = a + b, e = a + b }

d = a;

{ a = b, c = b, d = b, e = a + b }

f = a + b;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

{ }

a = b;

{ a = b }

c = a;

{ a = b, c = b }

d = a + b;

{ a = b, c = b, d = a + b }

e = d;

{ a = b, c = b, d = a + b, e = a + b }

d = a;

{ a = b, c = b, d = b, e = a + b }

f = e;

{ a = b, c = b, d = b, e = a + b, f = a + b }

Common Subexpression Elimination

a = b;

c = **a**;

d = a + b;

e = **d**;

d = **a**;

f = **e**;

Live Variables

- The analysis corresponding to dead code elimination is called **liveness analysis**.
- A variable is **live** at a point in a program if later in the program its value will be read before it is written to again.
- Dead code elimination works by computing liveness for each variable, then eliminating assignments to dead variables.

Computing Live Variables

- To know if a variable will be used at some point, we iterate across the statements in a basic block in reverse order.
- Initially, some small set of values are known to be live (which ones depends on the particular program).
- When we see the statement $\mathbf{a} = \mathbf{b} + \mathbf{c}$:
 - Just before the statement, \mathbf{a} is not alive, since its value is about to be overwritten.
 - Just before the statement, both \mathbf{b} and \mathbf{c} are alive, since we're about to read their values.
 - (*what if we have $\mathbf{a} = \mathbf{a} + \mathbf{b}$?*)

Liveness Analysis

a = b;

c = a;

d = a + b;

e = d;

d = a;

f = e;

Liveness Analysis

a = b;

c = a;

d = a + b;

e = d;

d = a;

f = e;

{ b, d }

Liveness Analysis

a = b;

c = a;

d = a + b;

e = d;

d = a;

{ b, d, e }

f = e;

{ b, d }

Liveness Analysis

a = b;

c = a;

d = a + b;

e = d;

{ a, b, e }

d = a;

{ b, d, e }

f = e;

{ b, d }

Liveness Analysis

a = b;

c = a;

d = a + b;

{ a, b, d }

e = d;

{ a, b, e }

d = a;

{ b, d, e }

f = e;

{ b, d }

Liveness Analysis

a = b;

c = a;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b, e }

d = a;

{ b, d, e }

f = e;

{ b, d }

Liveness Analysis

```
a = b;  
{ a, b }  
c = a;  
{ a, b }  
d = a + b;  
{ a, b, d }  
e = d;  
{ a, b, e }  
d = a;  
{ b, d, e }  
f = e;  
{ b, d }
```

Liveness Analysis

{ b }

a = b;

{ a, b }

c = a;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b, e }

d = a;

{ b, d, e }

f = e;

{ b, d }

Dead Code Elimination

{ b }

a = b;

{ a, b }

c = a;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b, e }

d = a;

{ b, d, e }

f = e;

{ b, d }

Dead Code Elimination

```
    { b }  
    a = b;  
    { a, b }  
    c = a;  
    { a, b }  
    d = a + b;  
    { a, b, d }  
    e = d;  
    { a, b, e }  
    d = a;  
    { b, d, e }  
    f = e;  
    { b, d }
```

Dead Code Elimination

{ b }

a = b;

{ a, b }

c = a;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b, e }

d = a;

{ b, d, e }

{ b, d }

Dead Code Elimination

```
    { b }  
    a = b;  
    { a, b }  
    c = a;  
    { a, b }  
    d = a + b;  
    { a, b, d }  
    e = d;  
    { a, b, e }  
    d = a;  
    { b, d, e }  
  
    { b, d }
```

Dead Code Elimination

```
    { b }  
    a = b;  
    { a, b }
```

```
    { a, b }  
    d = a + b;  
    { a, b, d }  
    e = d;  
    { a, b, e }  
    d = a;  
    { b, d, e }
```

```
    { b, d }
```

Dead Code Elimination

```
a = b;
```

```
d = a + b;
```

```
e = d;
```

```
d = a;
```

Liveness Analysis II

a = b;

d = a + b;

e = d;

d = a;

Liveness Analysis II

a = b;

d = a + b;

e = d;

d = a;
{ b, d }

Liveness Analysis II

a = b;

d = a + b;

e = d;

{ a, b }

d = a;

{ b, d }

Liveness Analysis II

```
a = b;
```

```
d = a + b;  
{ a, b, d }  
e = d;  
{ a, b }  
d = a;  
{ b, d }
```

Liveness Analysis II

a = b;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b }

d = a;

{ b, d }

Liveness Analysis II

{ b }

a = b;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b }

d = a;

{ b, d }

Dead Code Elimination

{ b }

a = b;

{ a, b }

d = a + b;

{ a, b, d }

e = d;

{ a, b }

d = a;

{ b, d }

Dead Code Elimination

```
{ b }  
a = b;
```

```
{ a, b }
```

```
d = a + b;  
{ a, b, d }  
e = d;  
{ a, b }  
d = a;  
{ b, d }
```

Dead Code Elimination

{ b }

a = b;

{ a, b }

d = a + b;

{ a, b, d }

{ a, b }

d = a;

{ b, d }

Dead Code Elimination

```
a = b;
```

```
d = a + b;
```

```
d = a;
```

Liveness Analysis III

`a = b;`

`d = a + b;`

`d = a;`

Liveness Analysis III

a = b;

d = a + b;

d = a;
{b, d}

Liveness Analysis III

a = b;

d = a + b;

{a, b}

d = a;

{b, d}

Liveness Analysis III

a = b;

{a, b}

d = a + b;

{a, b}

d = a;

{b, d}

Liveness Analysis III

{b}

a = b;

{a, b}

d = a + b;

{a, b}

d = a;

{b, d}

Dead Code Elimination

{b}
a = b;

{a, b}
d = a + b;

{a, b}
d = a;
{b, d}

Dead Code Elimination

{b}
a = b;

{a, b}

d = a + b;

{a, b}

d = a;
{b, d}

Dead Code Elimination

{b}
a = b;

{a, b}

{a, b}

d = a;
{b, d}

Dead Code Elimination

```
a = b;
```

```
d = a;
```


A Combined Algorithm

A Combined Algorithm

`a = b;`

`c = a;`

`d = a + b;`

`e = d;`

`d = a;`

`f = e;`

A Combined Algorithm

a = b;

c = a;

d = a + b;

e = d;

d = a;

f = e;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

e = d;

d = a;

f = e;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

e = d;

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

e = d;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

e = d;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

d = a + b;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

c = a;

{a, b}

d = a;

{b, d}

A Combined Algorithm

a = b;

{a, b}
d = a;

{b, d}

A Combined Algorithm

{b}
a = b;

{a, b}
d = a;

{b, d}

A Combined Algorithm

`a = b;`

`d = a;`

Next Time

- **Formalisms for Local Optimizations**
 - Transfer functions and semantics
- **Global optimization**
 - Optimizing across basic blocks.
 - Meet operators and the dataflow framework.